

## Pessimism Reduction in Static Timing Analysis using Timing and Logic Filtering

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Industry Support



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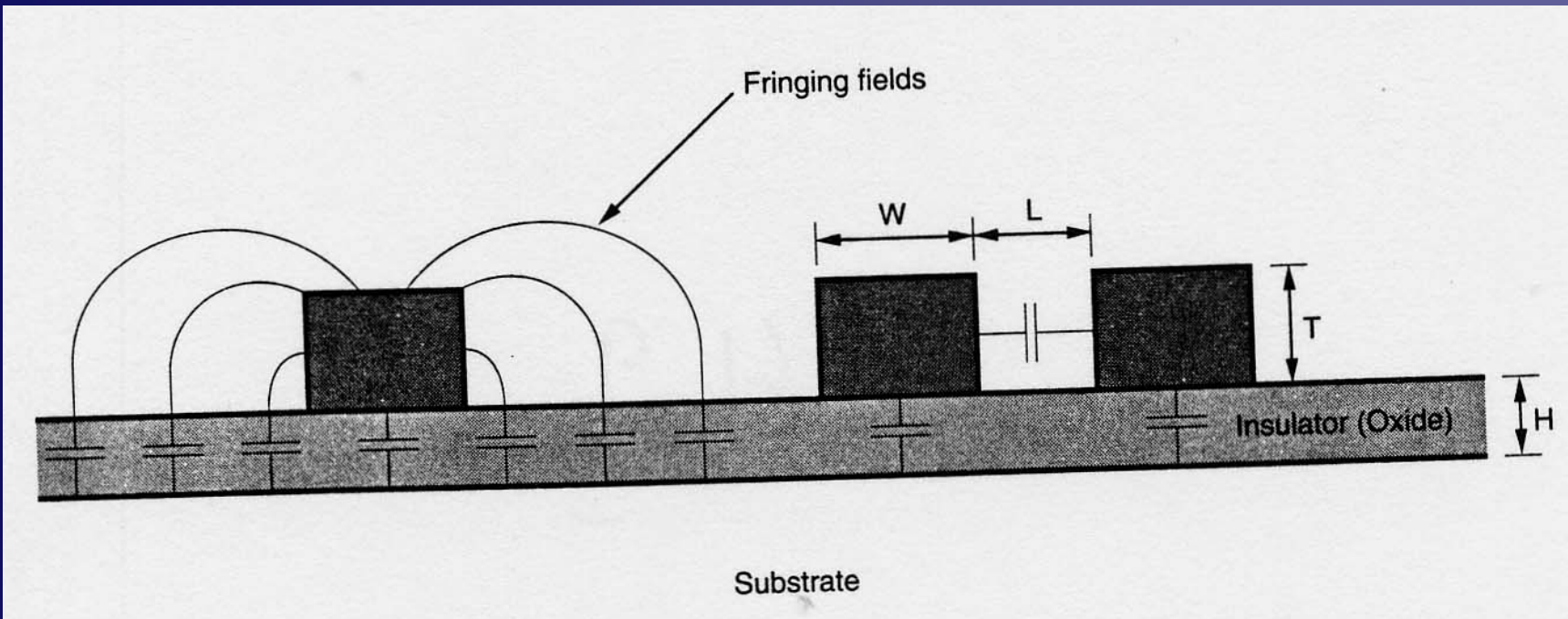
# Outline

- Coupling aware timing analysis
- Charge sharing based accurate coupling model
- Logic Filtering overview
- Filtering algorithm
- Experimental results
- Conclusions

## Origin of coupling : Increase in aspect ratio

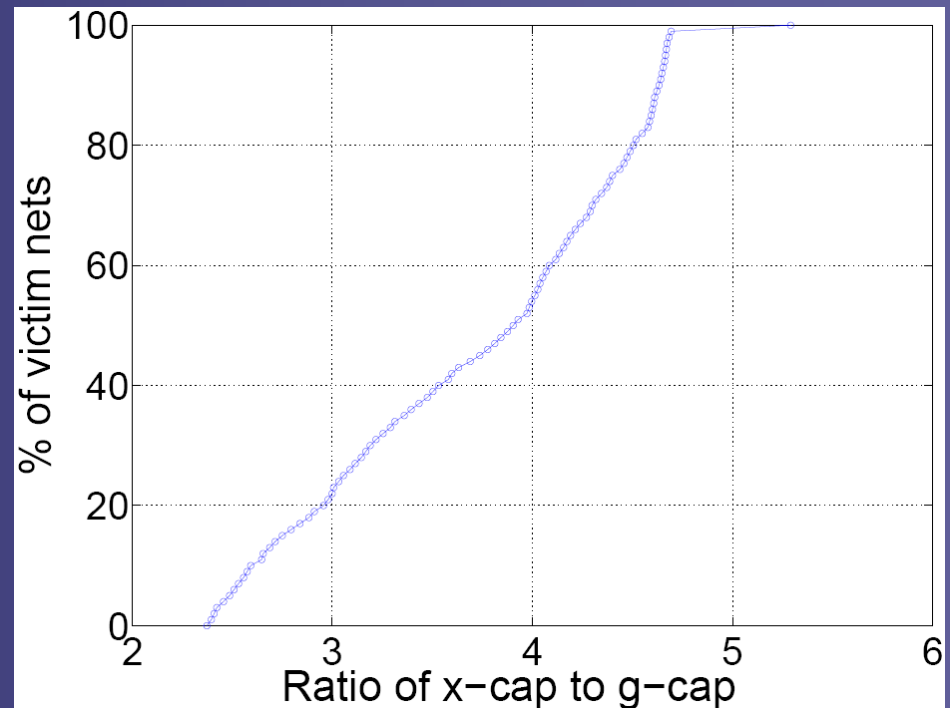
- Pitch =  $W + L$
- Aspect ratio:  $AR = T/W$ 
  - Old processes had  $AR \ll 1$   $T \ll W$
  - Modern processes have  $AR \approx 2$   $T \approx 2W$ 
    - Pack in many skinny wires

*Significant Lateral  
Cross Sectional Area*



# Coupling dominates

- Coupling cap dominates interconnect parasitics
- Graph shows ratio of coupling cap vs. ground cap of nets
  - Parasitics extracted from 65 nm logic block
  - Industrial Microprocessor design from Intel



# Coupling aware Timing Analysis

- **Timing Analysis with x-cap iterative**
- **Iterative analysis with continuous models: Chen et.al ICCAD 2000**
- **Iterative analysis with discrete models: Sapatnekar et.al ICCAD 2000, Chen et.al ICCAD 2000, Arunachalam et.al DAC 2000, Das et. al ICCD 2006**
- **Iterative coupling aware analysis without logic constraints is pessimistic**

# Pessimism reduction using Logic Filtering

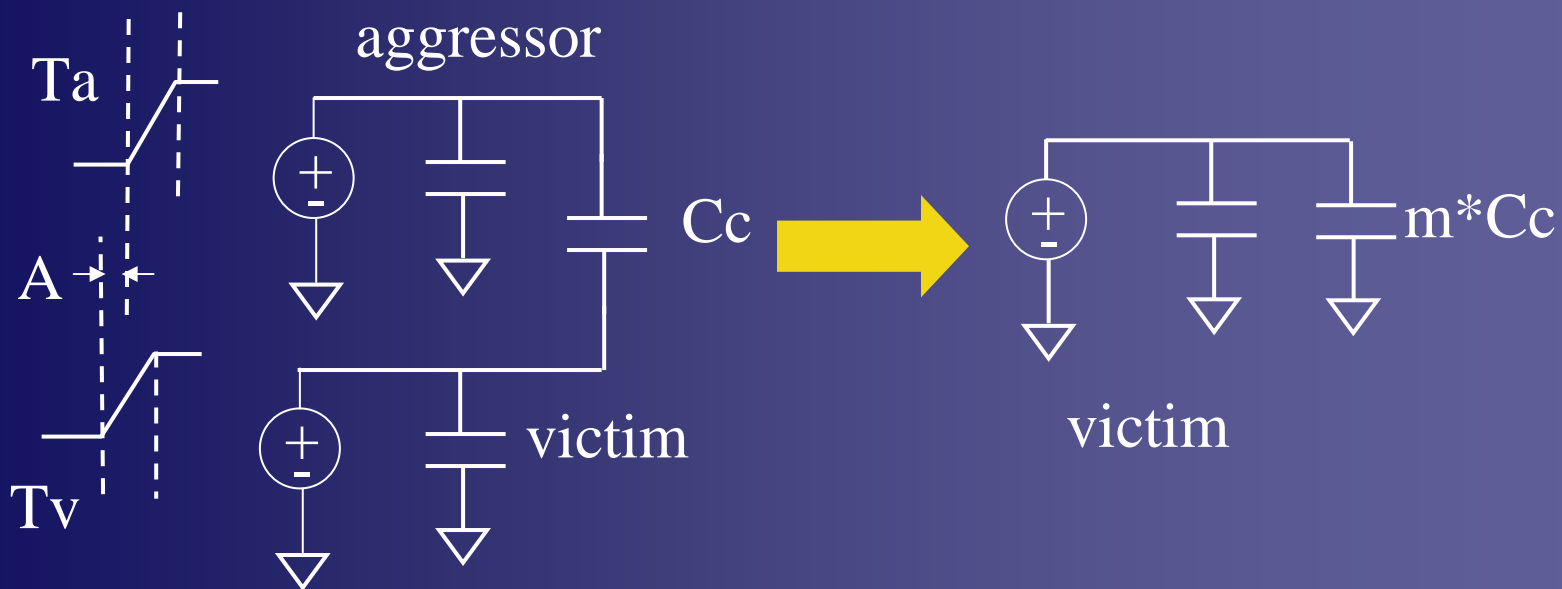
- **Generating Logic Conditions on a set of nets**
  - False path identification (numerous papers)
  - Finding false aggressors using logic implications (Glebov et. al.)
  - Finding false aggressors using propositional satisfiability (Sakallah et. al.)
  - Generating SAT formulas based on timing window convergence
    - False Coupling Exploration in Timing Analysis (Tseng et. al.)
    - Temporofunctional crosstalk noise analysis (Chai et. al.)
- **Previous research gives potential optimistic results**
  - No iteration on logic assumptions
  - Ignoring logic and timing interaction
  - Overly-simplistic coupling model

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# Coupling Model

- Arrival time and slope based coupling model (charge sharing)
  - $m = f(\text{victim arrival time and slope, aggressor arrival time and slope})$
- Circuit model for Noise cluster



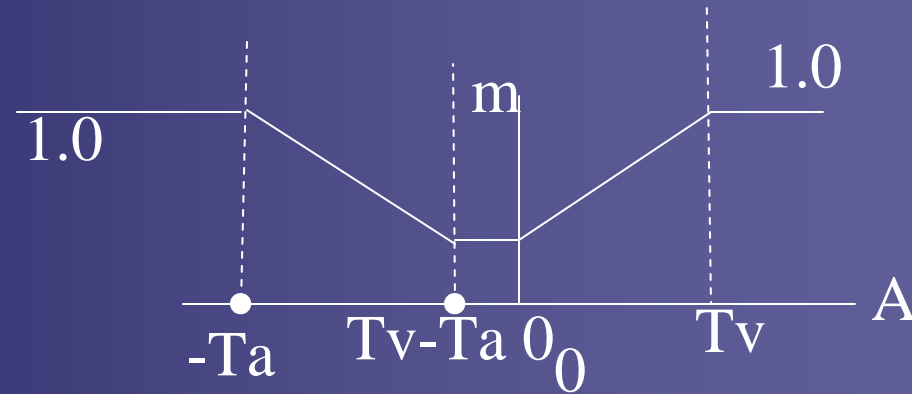
- Arrival Time : 0% , Slope : 0-100%, difference in arrival times =  $A$



# m vs A Plot (Similar Switching)

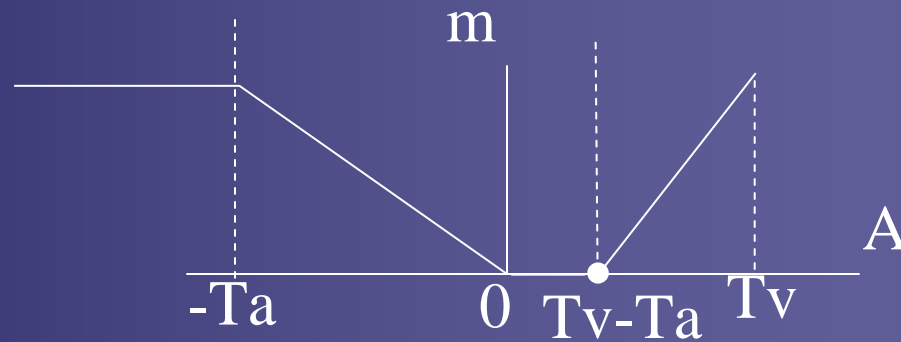
Case 1 :  $T_a > T_v$

$$m_{\min} = 1 - T_v / T_a$$

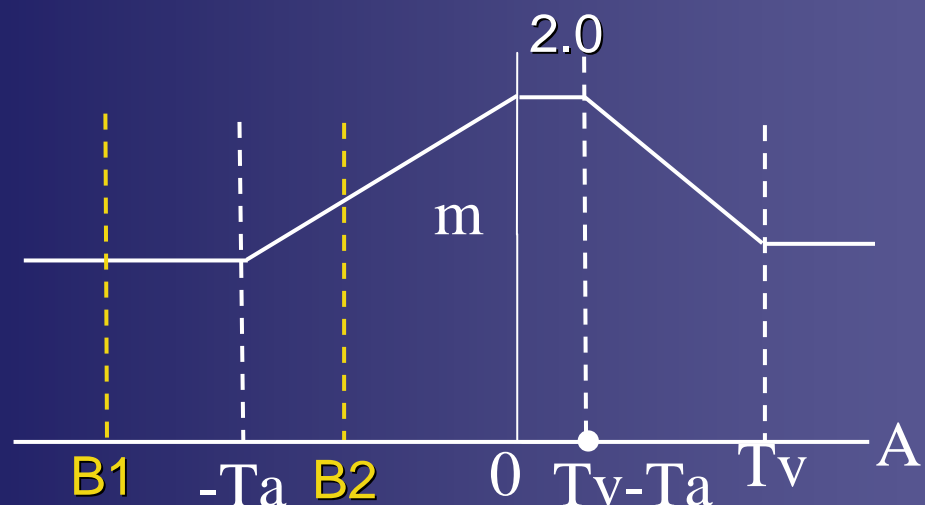


Case 2:  $T_a < T_v$

$$m_{\min} = 0$$



## Model Extension to Static Timing



- **Bounds  $B1, B2$  obtained by AT difference**
- **$T_a$  : Min TT associated with Aggressor**
- **$T_v$  : Max TT associated with Victim**
- **Worst  $m$  within  $B1$  and  $B2$**

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# Logic Filtering

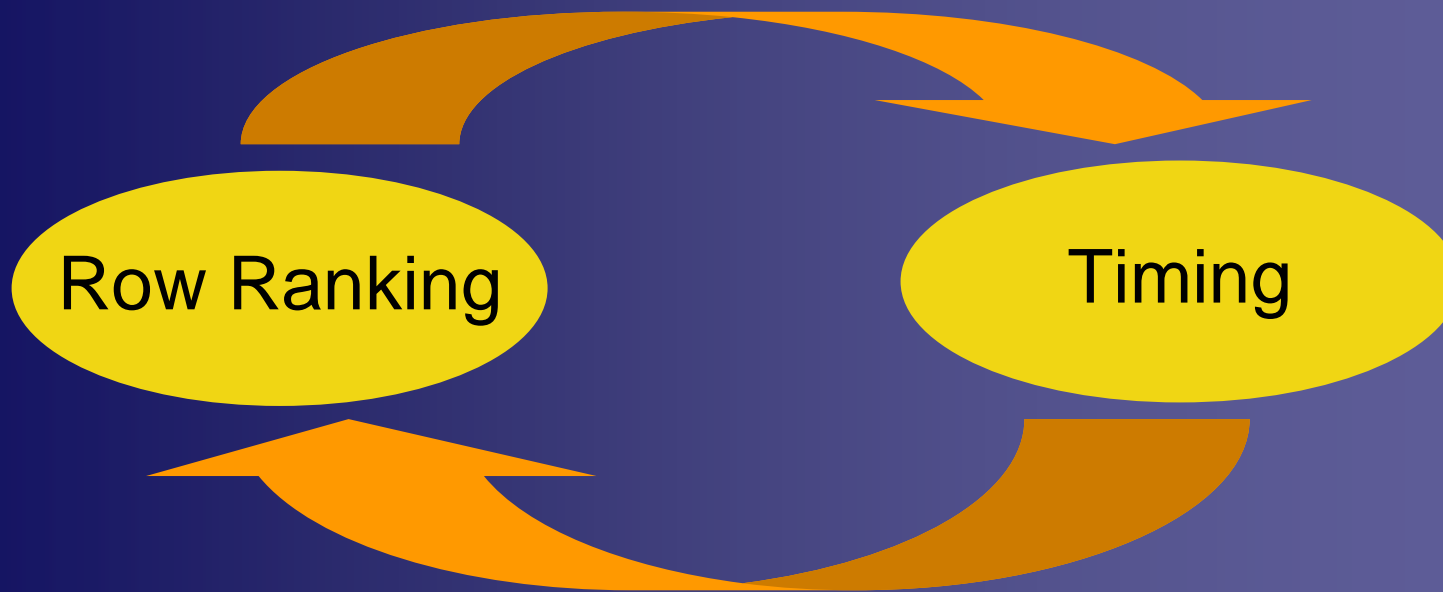
- Worst case might not happen (Due to Logic)

V	A1	A2	
R	R	R	X
R	S	R	✓
R	F	F	X
R	S	F	✓
.	.	.	
.	.	.	

- Idea is to find “worst” logically feasible row
  - Worst = Worst stage delay

# Logic Filtering: Chicken Egg Problem

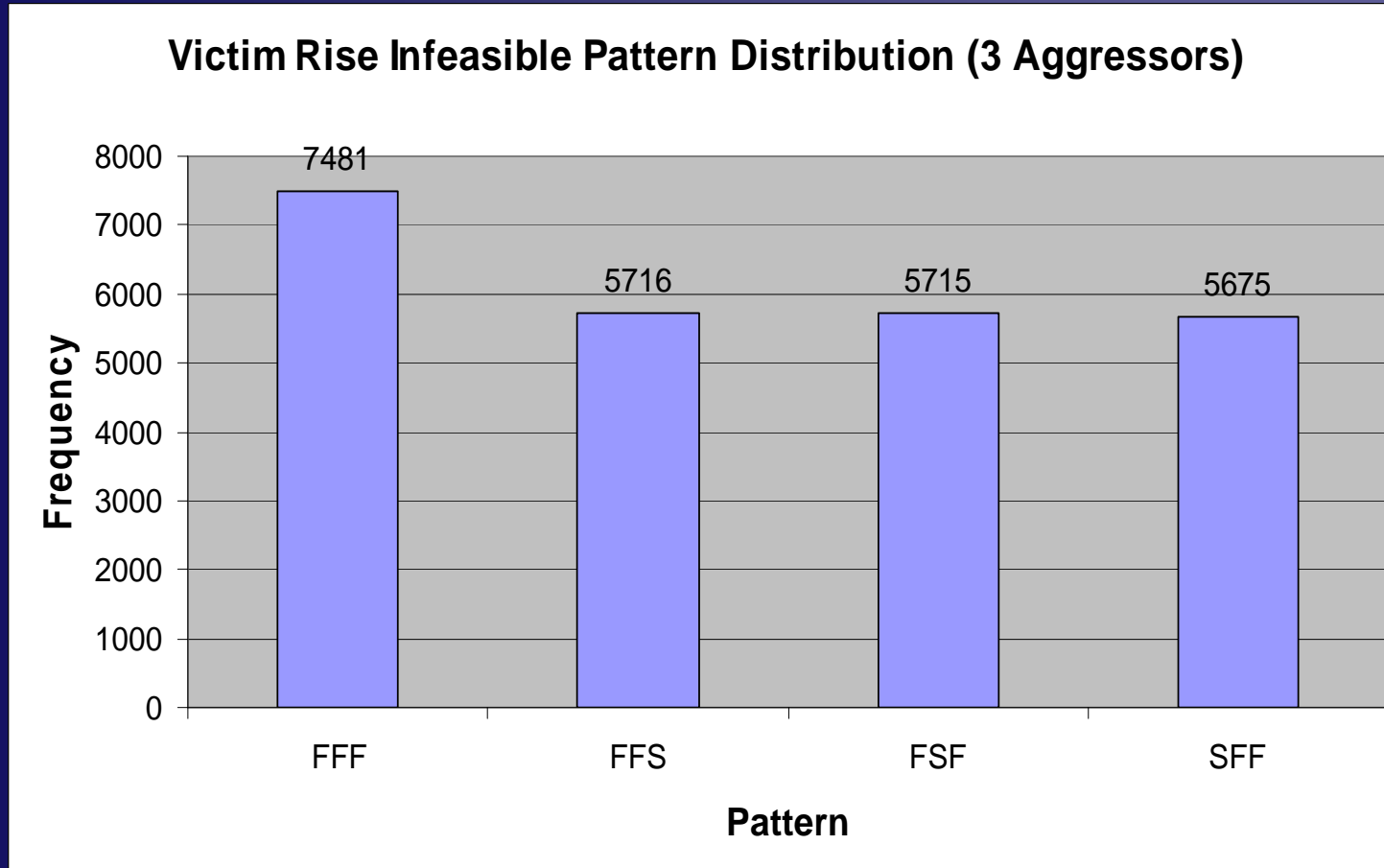
- Worst row generation (ranking) from Logic Tables
  - Row rank  $R = f(m, P)$  where  $P = \text{Parasitics}$
  - $m = g(A, Tv, Ta)$
  - $A = h1(m, P)$ ,  $Tv = h2(m, P)$ ,  $Ta = h3(m, P)$



# Logic Table Generation

- **We used**
  - Random circuit simulation with vectors
  - Logic implication based pruning
  - ATPG based techniques
  - SAT based techniques
- **Guaranteed conservative ignoring glitches**
  - 0-delay model
- **Complexity: expensive if SAT clause > 6 literals**
  - 20K Victim nets, 5 aggressors each :  $20K \times 2 \times 3^3 = 1,080,000 = \frac{1}{2}$  day
  - If Aggressors = Median of #aggs (25) :  $20K \times 2 \times 3^{25} = \text{EEK!}$
- **We used clause size 4: 1 victim, 3 aggressors**

# Logic Condition Generation

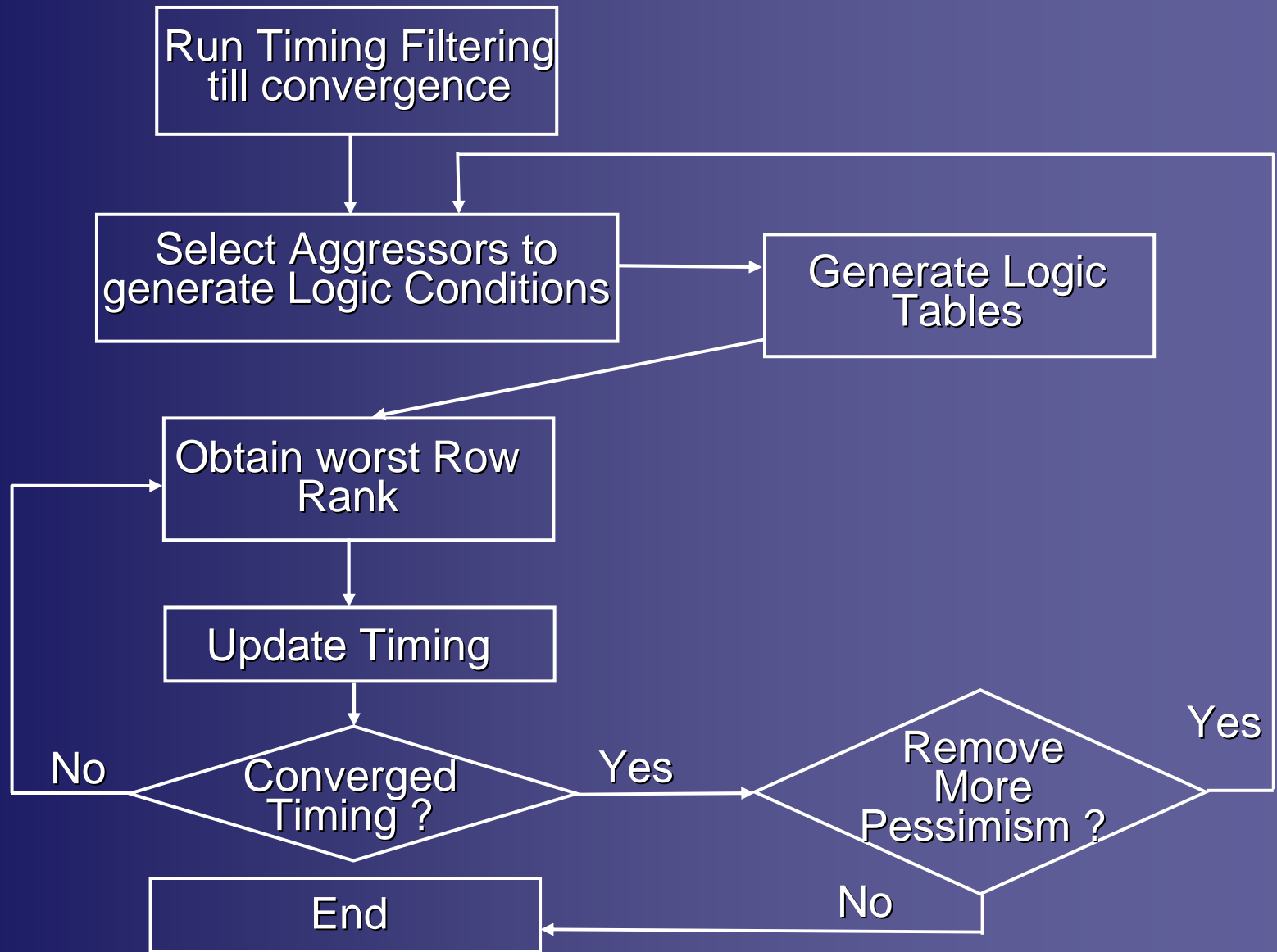


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# Proposed Flow



# Aggressor Selection

- Compute sensitivities of each aggressor
  - Victim V, 3 Aggressors A1, A2, A3
  - Converged Timing m's :  $m1_{nom}$ ,  $m2_{nom}$ ,  $m3_{nom}$
  - Receiver elmore delay =  $f(m1, m2, m3)$

- $$f(m1, m2, m3) = f(m1_{nom}, m2_{nom}, m3_{nom}) + (\partial f / \partial m1) x (m1 - m1_{nom}) + \dots$$

- $\partial f / \partial m1 = \{f(m1, m2_{nom}, m3_{nom}) - f(m1_{nom}, m2_{nom}, m3_{nom})\} / (m1 - m1_{nom})$
- $S_{m1} = (\partial f / \partial m1) / f(m1_{nom}, m2_{nom}, m3_{nom})$
- Change m1 to 1.0 for  $\partial f / \partial m1$  computation

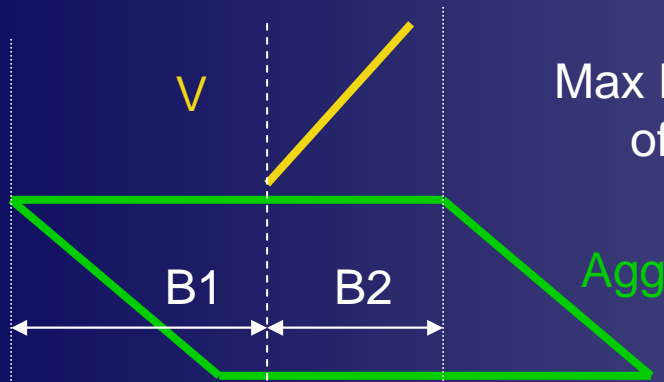
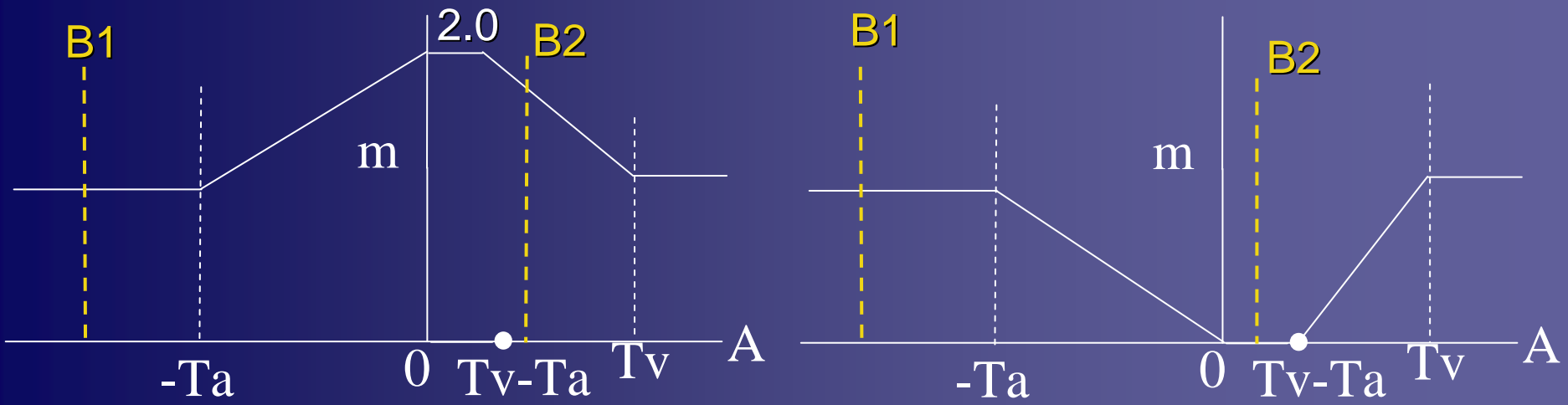
# Logic Filtering Algorithm

- Initialization :  $l_{mcf_0}, mcf_0 \leftarrow (0,2), Windows_0 \leftarrow f(mcf_0)$
- while(  $|Windows_i - Windows_{i-1}| > \epsilon$  )
  - if( LOGIC\_FILTERING)
    - Compute  $l_{mcf_i}$  for selected aggressors
    - $R \leftarrow$  Identify worse ranked row using  $l_{mcf_i}$
    - $mcf_i \leftarrow g(Windows_{i-1}, R)$
  - else if ( TIMING\_FILTERING)
    - $mcf_i \leftarrow h(Windows_{i-1})$
  - $Windows_i \leftarrow f(mcf_i)$

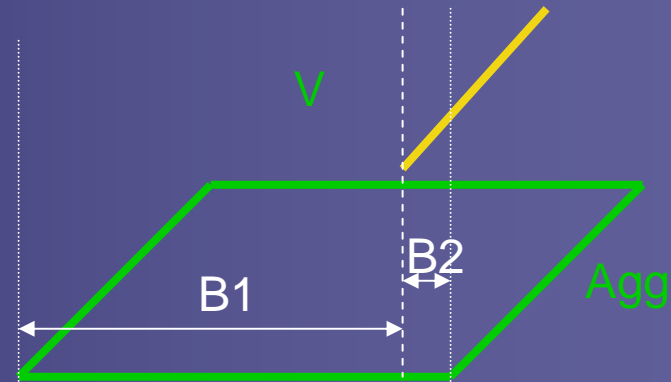
Logic MCF

Fix Point Computation

# Logic MCF Computation (Example)



Max Rise Event of Victim




Logic MCF Fall  
(Fall Twin of Aggressor)  
Max MCF = 2.0

Logic MCF  
Stable = 1.0

Logic MCF Rise  
(Rise Twin of Aggressor)  
Max MCF = 1.0

# Logic Filtering Algorithm: Worst Row Calculation

- Given Victim Cluster  $\langle V, A1, A2, A3 \rangle$
- x-caps:  $CC_{A1}, CC_{A2}, CC_{A3}$
- LMCFs:  $LMCF_{A1}, LMCF_{A2}, LMCF_{A3}$
- Logic Feasibility Table for cluster

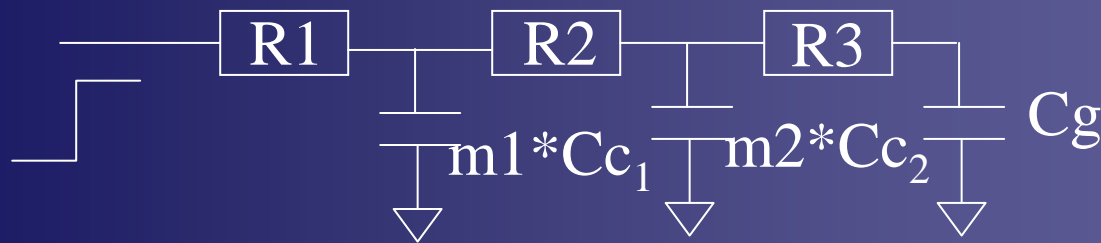


V	A1	A2	A3
R	F	S	S
R	S	R	S
F	S	S	R
F	S	R	S

$$\begin{aligned} \text{Rank} = & LMCF_{A1}[F] \times CC_{A1} + \\ & LMCF_{A2}[S] \times CC_{A2} + \\ & LMCF_{A3}[S] \times CC_{A3} \end{aligned}$$

# Accurate Ranking

- We should not ignore resistances and ground cap !
- Circuit model of victim cluster



$$E_d = m1Cc_1 * R1 + m2Cc_2 * (R1+R2) + C_g(R1+R2+R3)$$

- Ranking by Elmore delay computation
  - Choose the row with worst Elmore delay

# Conservative Row Ranking

- Consider MAX RISE event of victim
- $CC_{A1} = 3, CC_{A2} = 2.5, CC_{A3} = 2$
- $LMCF_{A1} = \{ F = 2.0, R = 1.0, S = 1 \}$ ,  $LMCF_{A2} = \{ F = 1.8, R = 1.0, S = 1 \}$

V	A1	A2	A3
R	F	S	S
R	S	R	S
F	S	S	R
F	S	R	S

Iteration i

$$\text{Rank}(i) = 2.0 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 10.5$$

$$\text{Rank}(i) = 1.8 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.9$$

# Conservative Row Ranking

- Consider MAX RISE event of victim
- $CC_{A1} = 3, CC_{A2} = 2.5, CC_{A3} = 2$
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V	A1	A2	A3
R	F	S	S
R	S	R	S
F	S	S	R
F	S	R	S

Iteration i

$$\text{Rank}(i) = 2.0 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 10.5$$

$$\text{Rank}(i) = 1.8 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.9$$



# Conservative Row Ranking

- Consider MAX RISE event of victim
- $CC_{A1} = 3, CC_{A2} = 2.5, CC_{A3} = 2$
- $LMCF_{A1} = \{F = 1.6, R = 1, S = 1\}, LMCF_{A2} = \{F = 1.7, R = 1, S = 1\}$

V	A1	A2	A3
R	F	S	S
R	S	R	S
F	S	S	R
F	S	R	S

Iteration i+1

$$\text{Rank}(i+1) = 1.6 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.3$$

$$\text{Rank}(i+1) = 1.7 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.6$$

# Conservative Row Ranking

- Consider MAX RISE event of victim
- $CC_{A1} = 3, CC_{A2} = 2.5, CC_{A3} = 2$
- $LMCF_{A1} = \{ F = 1.6, R = 1, S = 1 \}, LMCF_{A2} = \{ F = 1.7, R = 1, S = 1 \}$

	V	A1	A2	A3
→	R	F	S	S
→	R	S	R	S
	F	S	S	R
	F	S	R	S

Iteration i+1

$$\text{Rank}(i+1) = 1.6 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.3$$

$$\text{Rank}(i+1) = 1.7 \times 3 + 1.0 \times 2.5 + 1.0 \times 2 = 9.6$$

Must Consider same direction  
Even in MAX !

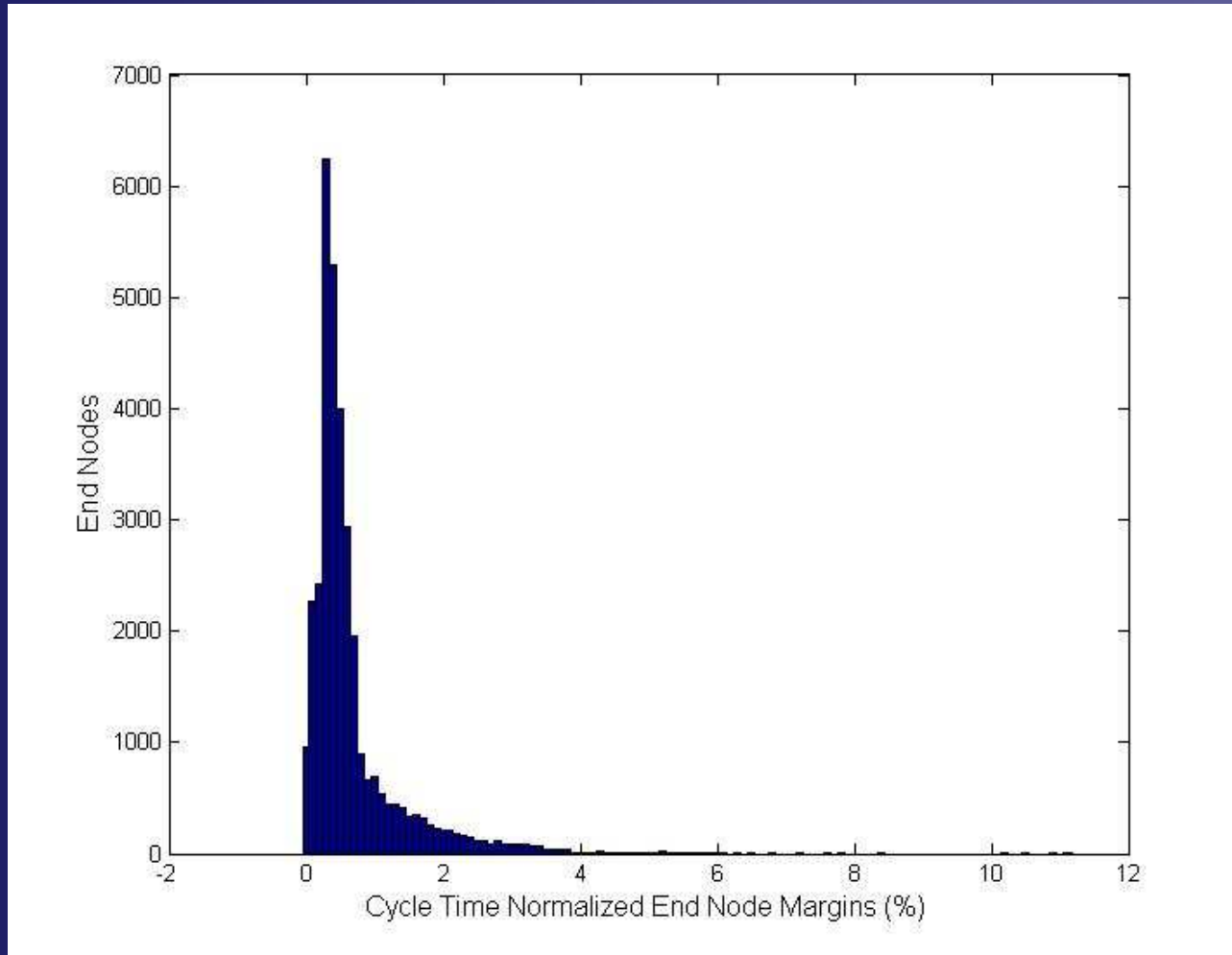
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# Results

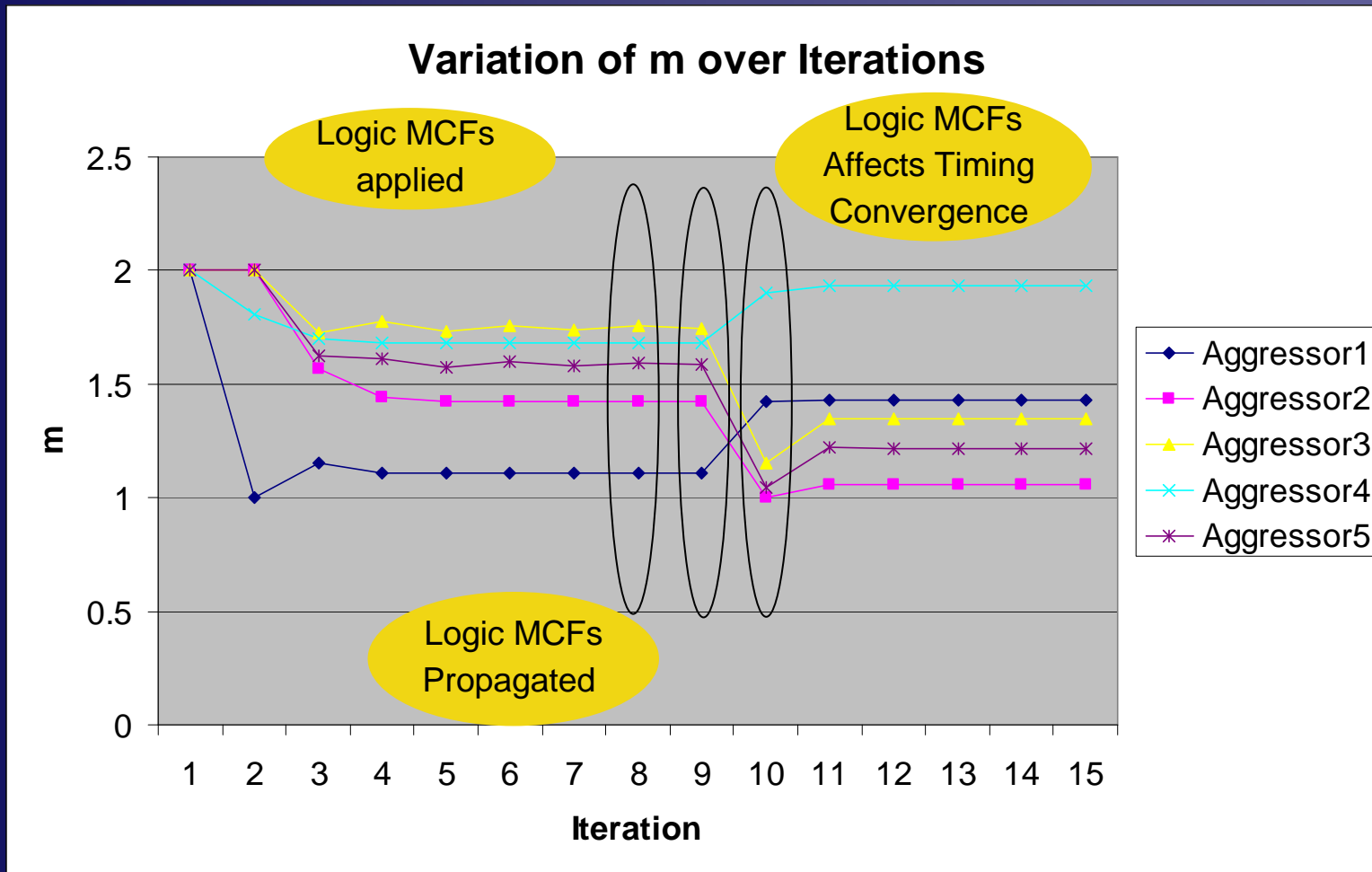
- Pessimism Reduction

- Logic Conditions applied after 8 iterations.
- Maximum : 11.18% Median : 0.44%



# Results

- Variation of  $m$  under logic + timing iterations



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## Conclusions

- Proposed novel algorithm to do coupling aware static timing analysis with logic constraints
- Dynamics of complex coupling model under timing and logic filtering is presented
- Iterations are needed to generate conservative timing data
- 11.18% maximum pessimism reduction observed on a functional block from 65 nm industrial microprocessor
- Future work includes extending such a model to statistical timing analysis and optimization

# Thank You