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Analyzing and Optimizing Energy Efficiency of Algorithms on DVS Systems:

--A First Step towards Algorithmic Energy Minimization--



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Optimizing Energy Consumption at Algorithmic Level

Optimizing energy consumption:

- Many HW mechanisms available (DVFS, DPM, ...)
- Important in system level design
- Used in earlier stages of SW development phase

However, fundamental concepts have not yet been completed.

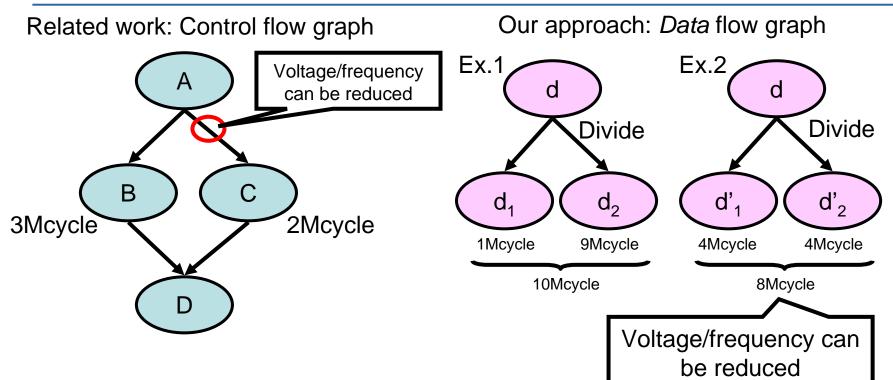
- Differences from performance optimization?
- Metrics?
- Programming logics and structure?
- Dataflow?

More precisely, we cannot answer: Which, either Quicksort or Heapsort, is more energy optimal?

Target and Objective

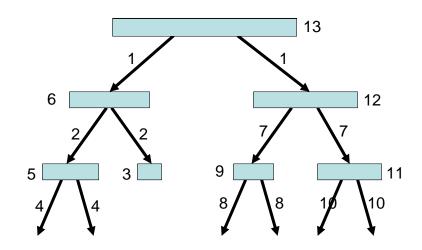
- Target
 - DVS systems
 - Deadline constraints
 - Algorithmic level
- Objective
 - Clarify the difference between energy optimization and performance optimization
 - Propose a measure for energy consumption
 - Study a case of algorithmic energy optimization
 - Answer "Quicksort vs Heapsort.
 Which is more energy optimal?"

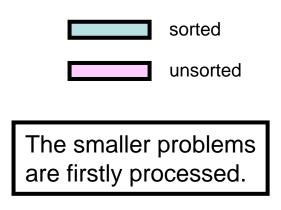
IntraDVS: Basic Concepts



- The selected branch decides the remaining cycles
 - Ex. Either block B or block C is executed
- The sizes of divided subproblems decide the remaining cycles

Review of Quicksort



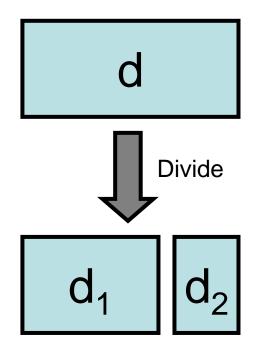


- Workload variation \rightarrow slack time
 - WCEC: $\delta_w(n) \propto n^2$
 - ACEC: $\delta_a(n) \propto n \log n$
- Problem: WCEC is too big!
 - \rightarrow Heapsort does not have much workload variance.

Remaining Predicted Execution Cycles

• WCEC

$$\delta_{\mathsf{W}}(d) = c(d) + \max_{d_1, d_2 \in \mathsf{div}(d)} (\delta_{\mathsf{W}}(d_1) + \delta_{\mathsf{W}}(d_2))$$



 $\delta_w(d)$: WCEC of processing d

c(d) + max(
$$\delta_w(d_1) + \delta_w(d_2)$$
)
Dividing time + Worst of sum of WCECs

Comparison of Remaining WCET of qsort

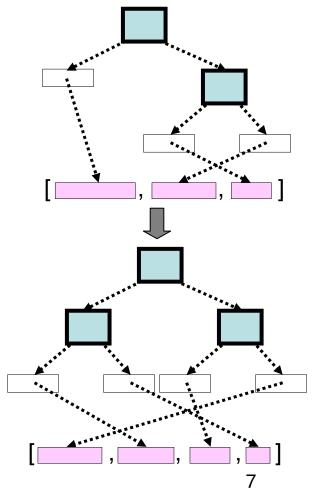
• Remaining WCET w/o concatenation:

 $\sum_{i=1}^{m} \delta_{\rm wc}(length({\tt xs}_i)) \qquad \text{(Approximation)}$

- In ordinary libraries, the smaller problem is firstly processed
 - PRO: Memory consumption is bound by O(log n)
 - CON: Remaining WCET does not rapidly decrease (early division technique)

	Decrease of RET	Memory Usage
Smallest	Slow	Low
Largest	Moderate	Middle
Larger	Rapid	High

Our approach: Processing larger problems



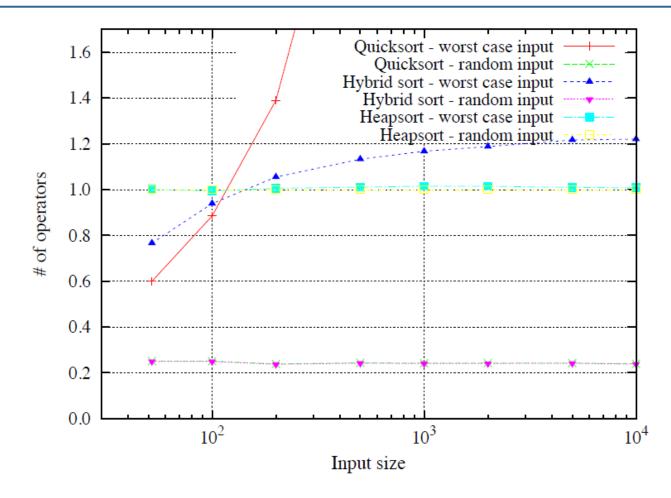
Energy-efficient Hybrid Sorting Algorithm

• Energy Efficiency = Original Algo. + Optimization

	Original Algo.	Energy Optimization
Quicksort	O(n ²)	Very good
Heapsort	Θ(n log n)	Bad
Hybrid	Θ(n log n)	Very good

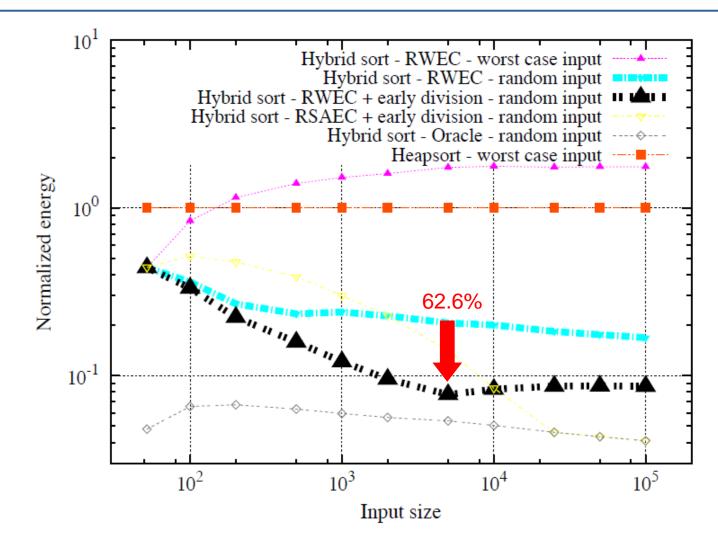
- Our solution: Hybrid sort = Quicksort + Heapsort
 - First, Quicksort: fast on average
 - Early division technique is used.
 - Performance is very high for almost all of the input.
 - At worst case, changed into Heapsort; WCEC is bounded.
 - ⇒ The energy efficiency of the sorting algorithm can be optimized on DVS systems.

Normalized Operator numbers of Sorting Programs

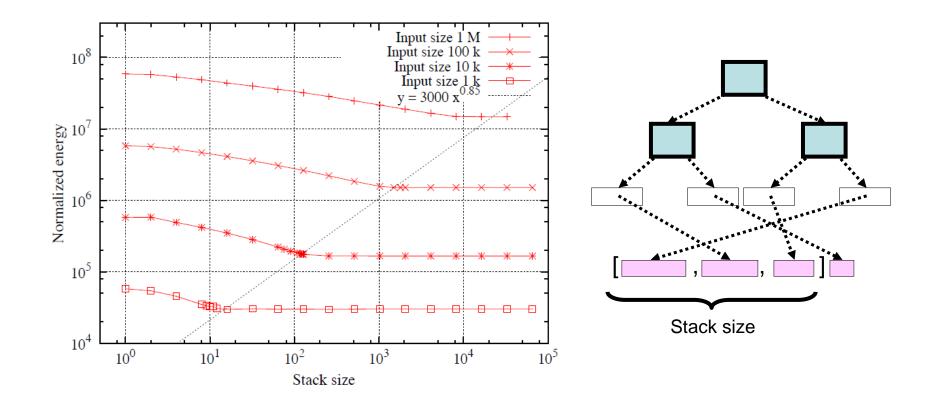


The data was obtained on a MIPS R5000 processor with 512 KB of secondary cache and 64 MB of main memory, using version 7.2.1 of the Silicon Graphics MIPSpro C++ compiler.

Normalized Energy of Sorting Programs Using DVS



Effects of the Stack Size on Energy Consumption Using SVS



- A tradeoff between energy and memory
- The greater stack size it has, the more energy is saved.
- Energy savings saturated to about 25%
- This saturation occurs at line $y = 3000 \times 0.85$

Measure for Evaluating Energy Consumption of Algorithm

- SVS: voltage scheduling before execution
- The optimal voltage of each task
 = The minimum voltage to finish the task execution exactly at the deadline

Lemma (Optimal SVS) For SVS without scaling bound, An algorithm of a given task is optimal on average iff $\delta_w^2 \delta_a$ is minimum.

Implications:

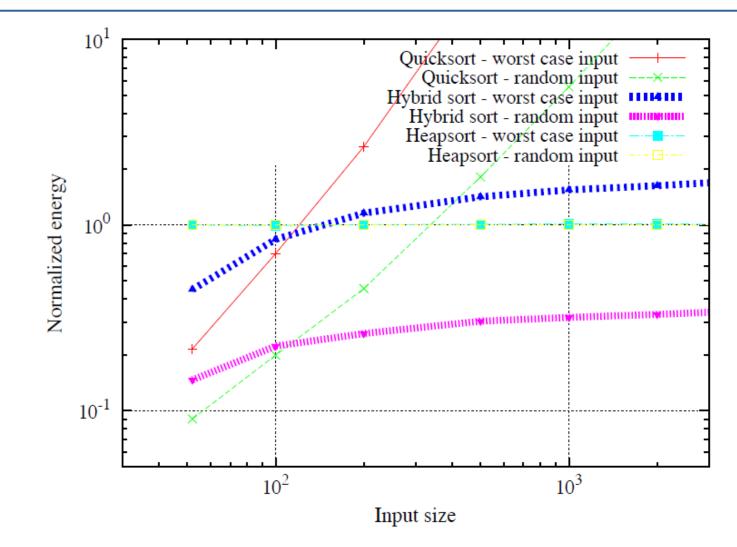
- Deadline does not affect the comparison of two algorithms
- Firstly, WCET δ_w should be reduced
- Secondly, ACET δ a

Proof

- 1. Let D be deadline.
- 2. The execution of frequency δ_w /D finishes exactly on deadline.
- 3. The corresponding power consumption is P(δ_w/D), which is the smallest in the case that f = δ_w/D because of the monotonicity of *P*.
- 4. The execution time is the number of cycles X divided by frequency, *i.e.*, X/(δ_w /D).
- 5. Thus, energy consumption becomes P(δ_w /D) · DX/ δ_w
- 6. Now, we are comparing the different algorithms under the same deadline, and therefore *D* is constant.
- 7. The average energy consumption is proportional to the average of $(\delta_w)^2 X$.

CAVEAT: If frequencies range over $[f_{min}, f_{max}]$, the objective function becomes $(\delta_w)^2 \cdot max(Df_{min}, min(Df_{max}, \delta_a))$

Normalized Energy of Sorting Programs Using SVS



Comparison of Energy Consumption of Sorting: A Case Study

- Question:
 - What is an energy efficient sorting on SVS systems?
- Our Answer[†]:

 - When the input size is large,

Average Energy Average Exec. Time	Hybrid << Heap << Quick Quick = Hybrid << Heap
Worst Energy Worst Exec. Time	Heap < Hybrid << Quick Heap < Hybrid <<< Quick

- Implications:
 - The proposed metrics enables this comparison
 - Energy optimization ≠ Performance optimization
 - Fastest on average \neq Energy optimal on average

[†]Only consider Quicksort, Heapsort and Hybrid sort

Concluding Remarks

- *Algorithmic* energy efficiency is *meaningful*.
- We propose
 - Energy efficient sorting algorithm
 - A measure for evaluating the optimal energy of algorithms
 - IntraDVS strategies using data flow information
 - ⇒ We can discuss which, either Quicksort or Heapsort, is more energy efficient.
- Future work
 - How to write energy efficient programs (partially published)
 - Compare energy efficiency of other algorithms (ongoing)
 - Expose a tradeoff with energy (ongoing)