

Novel Task Migration Framework on Configurable Heterogeneous MPSoC Platforms

Hao Shen Frédéric Pétrot*

System Level Synthesis Group, TIMA Laboratory

CNRS/Grenoble INP/UJF

43, Avenue Félix Viallet, 38031, Grenoble, France

hao.shen@imag.fr

frederic.petrot@imag.fr

TIMA Laboratory

Outline

Introduction

Task Migration Definition and Implementation

Task Migration Algorithms

Experimental Results

Conclusion and Future Works

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Trends (Power and Cost Constraints for Embedded Systems)

Multiple Processor System-on-Chip (MPSoC)

- Provide high Thread Level Parallelism (TLP)
- Provide high performance
- Power consumption and cost advantages

Configurable processor

- Different extended instruction set for different processors
- Dedicated compiler for each extended instruction set
- Provide high Instruction Level Parallelism (ILP)
- Power consumption and cost advantages

Heterogeneity

- Different configurable processor in a MPSoC for different kinds of applications
- Power consumption and cost advantages

Configurable Processors

Application Specific Instruction-set Processors (ASIP)

- Configuration for one application or a group of applications

Provide high Instruction Level Parallelism (ILP)

- SIMD instructions, MIMD instructions, specific instructions (such as Multiply-accumulate)

Configurable processors VS. RISC processors

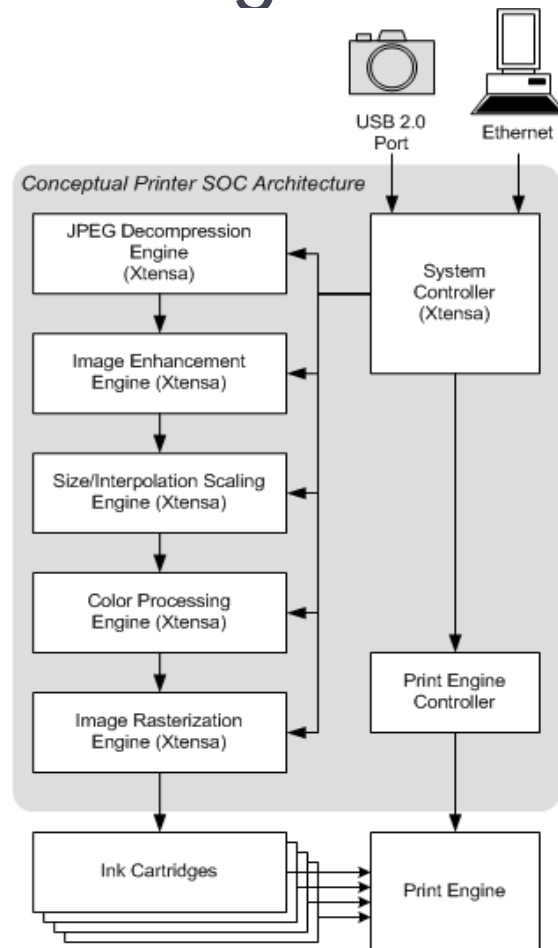
- Provide higher performance

Configurable processors VS. ASIC

- More flexibility and shorter time-to-market



One Existing Heterogeneous Multiple Configurable SoC Example



EPSON[®]



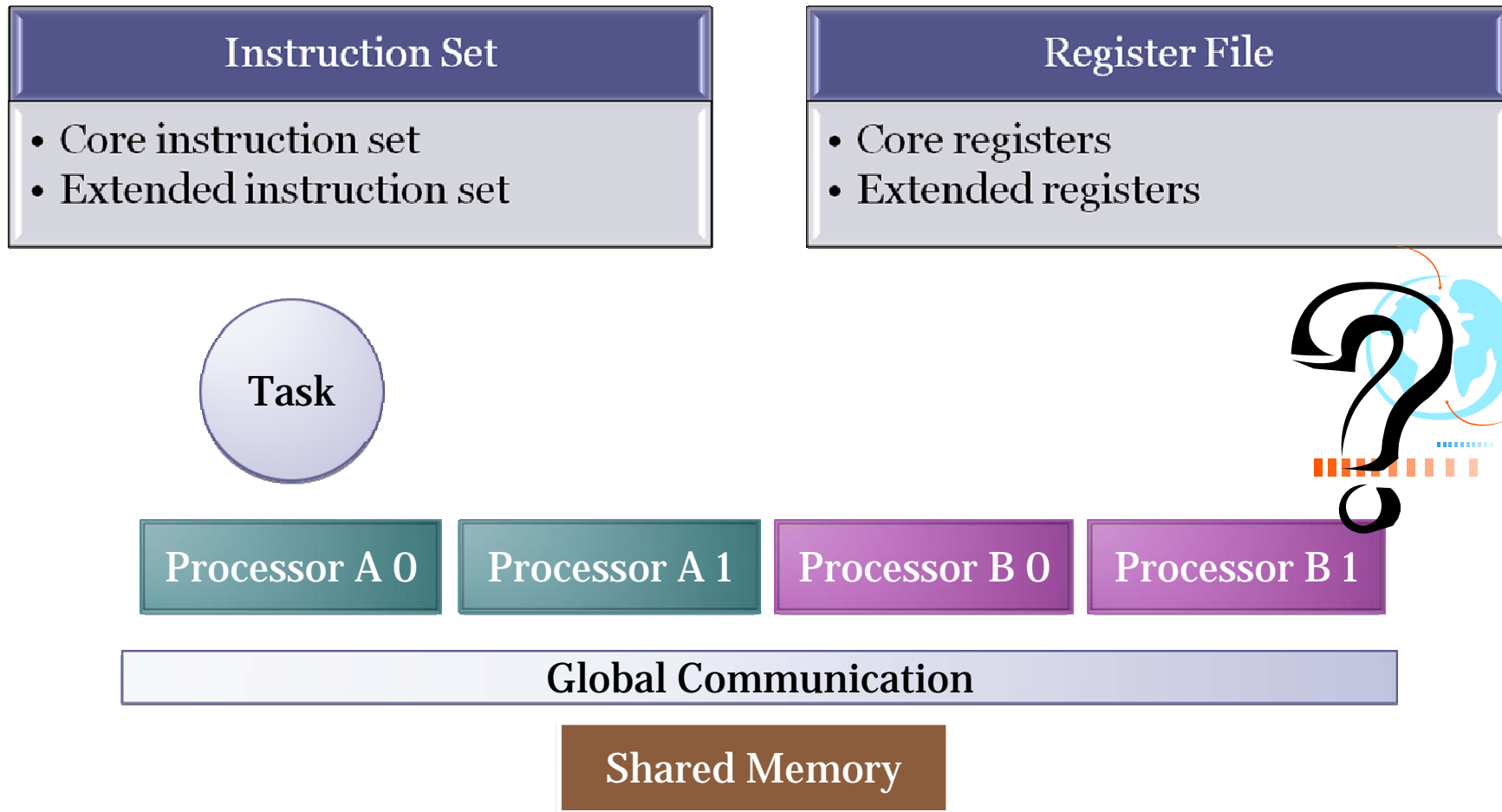
General printer solution

6 heterogeneous Xtensa processors

Fixed task mapping

Communication with FIFOs

What is the Problem of Configurable Processors?



Contribution of This Work

Task migration framework

- Support heterogeneous multiple configurable processor SoC
- Present realization details of this framework
- Add some formal description

Task migration algorithms

- Compare the efficiency of each algorithm
- Compare the migration cost of each algorithm

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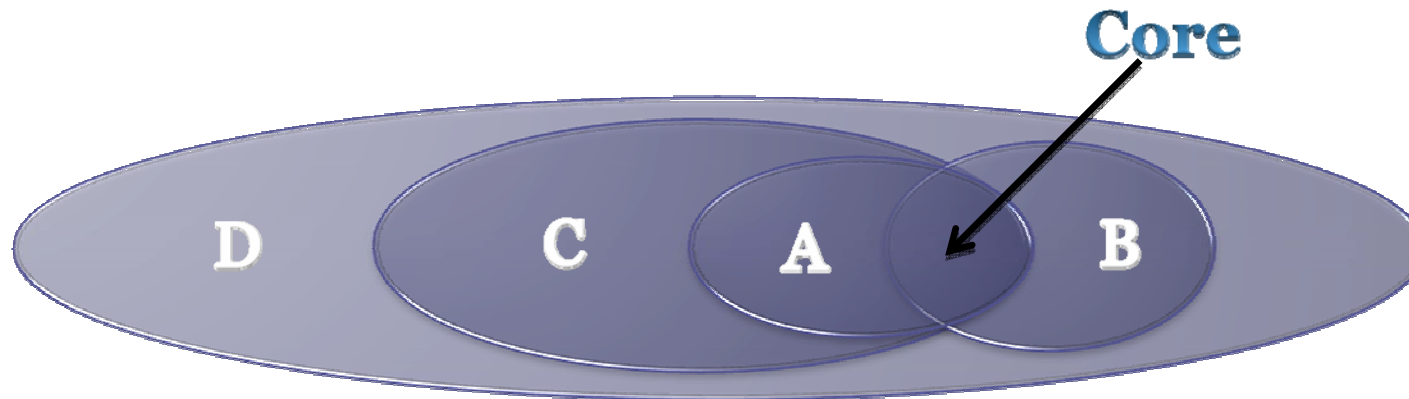
Task Migration Theory and Implementation

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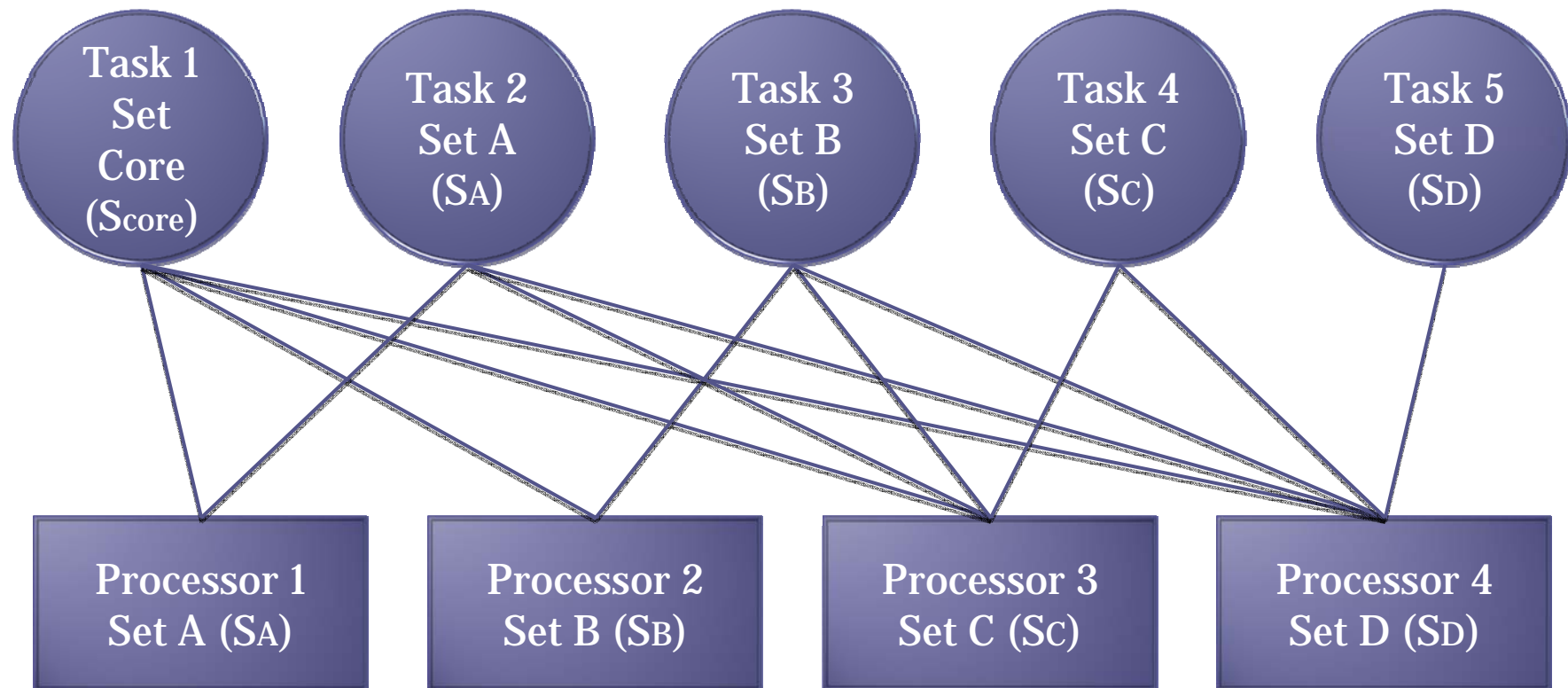
Instruction Set Relationship



$$\mathcal{S} = \{\text{Score}, S_A, S_B, S_C, S_D\}$$

$$\text{Score} \subset S_A, \dots, \text{Score} \subset S_D, S_A \subset S_C, S_A \subset S_D, S_C \subset S_D, S_B \subset S_D$$

Relationship Between Tasks and Processors



Instruction Set Identification

Assign ID for each CPU type (CPU_ISA_ID)

- Indicate the instruction set which it realizes

Assign ID for each task type (TASK_ISA_ID)

- Indicate the instruction set which it uses

Use bit operation to accelerate scheduling

- Use one bit to represent a standalone instruction set
- Test compatibility by using bit operations
- Save storage space for the OS realization

Scheduler Realization

Instruction set compatibility

- $(\text{CPU_ISA_ID} \mid \text{TASK_ISA_ID}) == \text{CPU_ISA_ID}$

Set	CPU_ISA_ID	Compatible Tasks	TASK_ISA_ID
Core	0x0000	Core	0x0000
A	0x0001	Core and A	0x0000, 0x0001
B	0x0010	Core and B	0x0000, 0x0010
C	0x0101	Core, A and C	0x0000, 0x0001, 0x0101
D	0x1111	Core, A, B, C and D	0x0000, 0x0001, 0x0010, 0x0101, 0x1111

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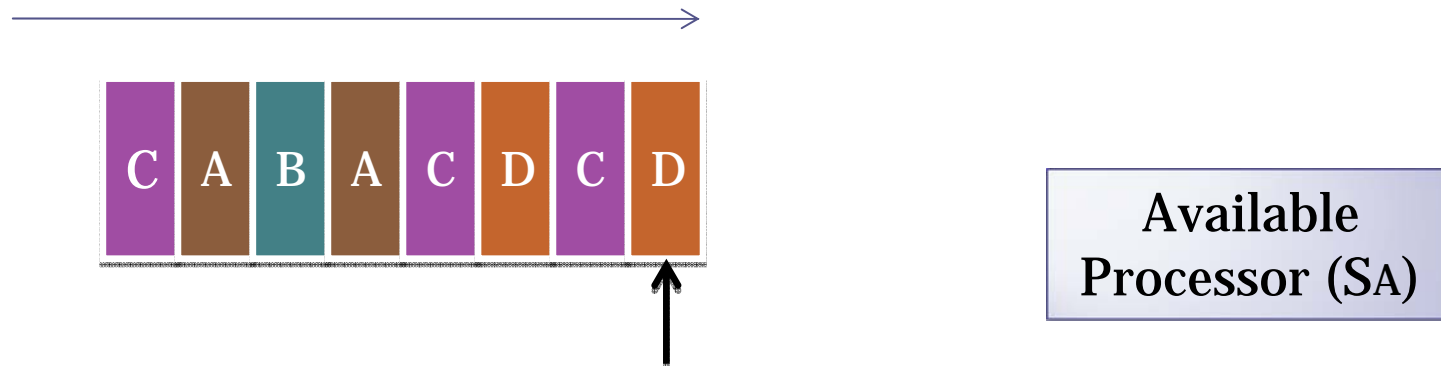
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First Match First Serve Algorithm

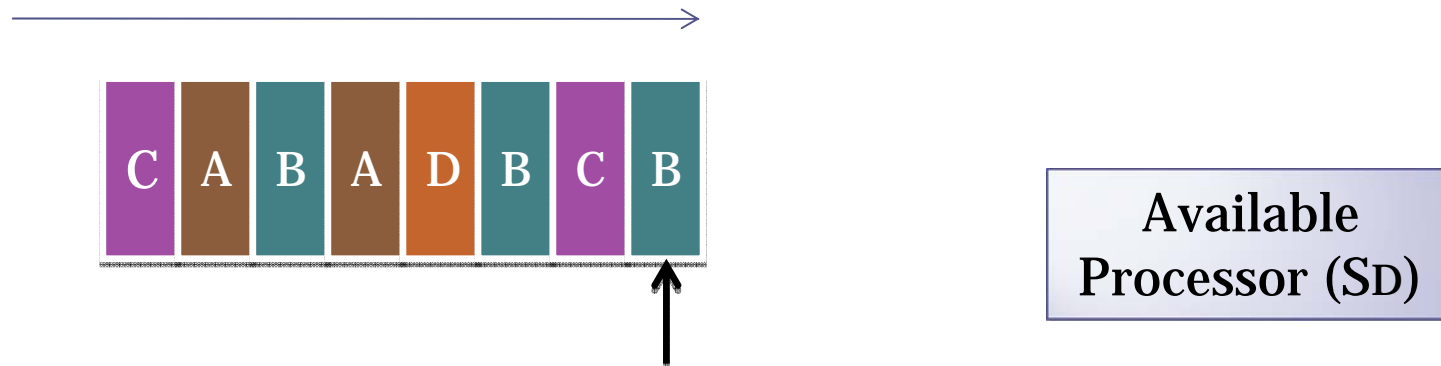


For $j=1$ to $|\mathbb{T}_{\text{queue}}|$ in the FIFO order

If $c(T_j, P_i) \in \mathbb{C}$ // T_j is compatible with P_i

Choose the task T_j

Most Compatible Algorithm



For all $T_j \in \mathbb{T}_{\text{queue}}$

If $c(T_j, P_i) \in \mathbb{C}$ // T_j is compatible with P_i

$\mathbb{T}_{\text{candidate}} = \mathbb{T}_{\text{candidate}} \cup T_j$ // Add T_j to candidate set

If $\mathbb{T}_{\text{candidate}} \neq \emptyset$

choose the task $T = \min(D(T_j \in \mathbb{T}_{\text{candidate}}, P_i))$

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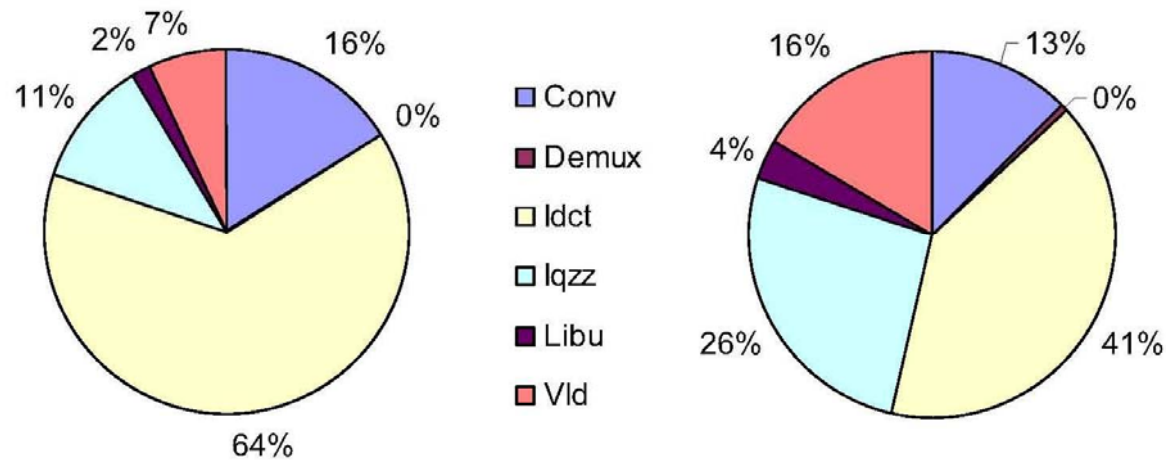
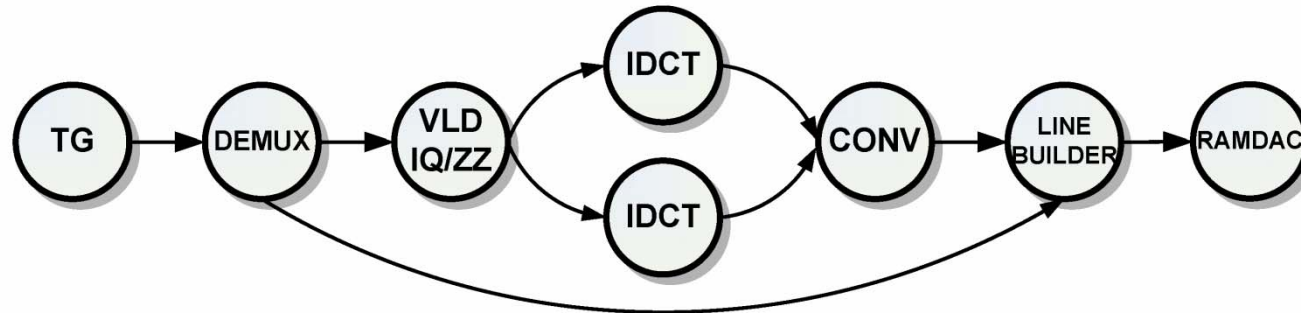
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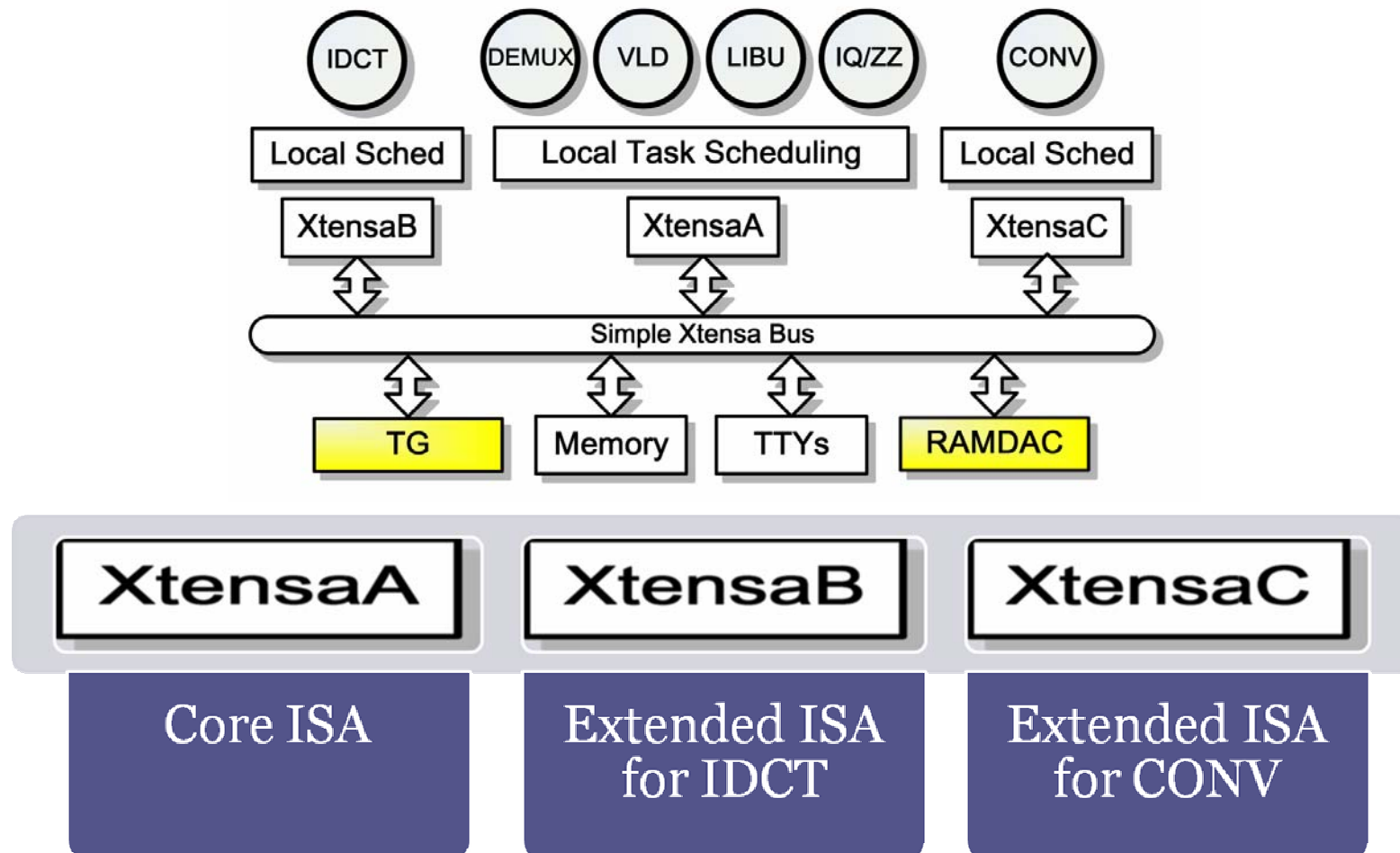
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Motion-JPEG Decoder Example and the Optimization



Heterogeneous MPSoC Architecture



Performance and Cost Advantages

	Fixed Task Assignment	FMFS Algorithm	Most Comp Algorithm	SMP Task Scheduling
Frame/s	1.44	2.88	2.70	2.88
CPUs Gate number	341,773	341,773	341,773	611,295
Perf/Cost	0.50	1.00	0.94	0.56

Heterogeneous architecture VS. homogeneous architecture

- May achieve the same performance
- May need smaller chip size (higher performance/cost ratio)

Task migration VS. fixed task assignment

- Need the same chip size
- Provide higher system performance (shorten the processor waiting time)

Different task migration algorithms have different performance

- FMFS requires less computation resource during migration

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Conclusion and Future Works

A task migration framework

- Support configurable processors
- Support heterogeneous MPSoC architectures
- Support several migration algorithms

Future works

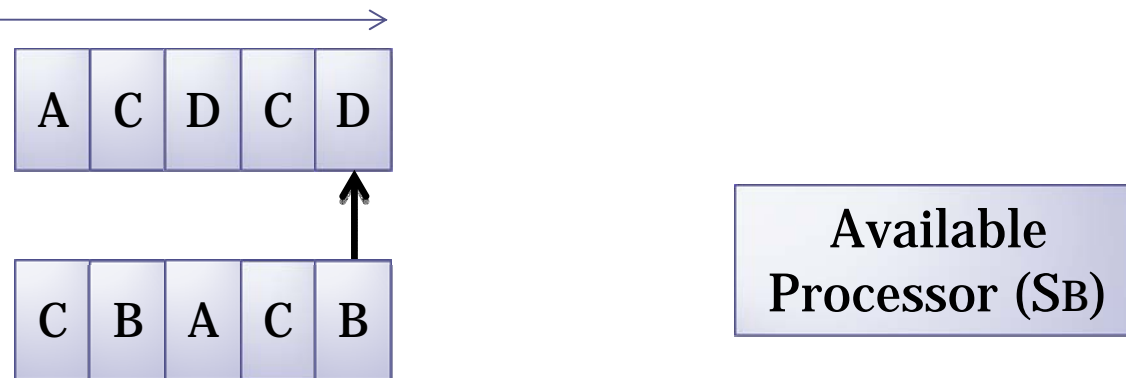
- Formalization this framework
- More complex case studies (benchmarks)

Questions & Answers

Hao.Shen@imag.fr

Frederic.Petrot@imag.fr

Priority Based Most Compatible Algorithm



For $k = 1$ to n // n queues with different priorities

For all $T_j \in \mathbb{T}_{\text{queue}}$

If $c(T_j, P_i) \in \mathbb{C}$ // T_j is compatible with P_i

$\mathbb{T}_{\text{candidate}} = \mathbb{T}_{\text{candidate}} \cup T_j$

If $\mathbb{T}_{\text{candidate}} \neq \emptyset$

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Time

- 20' presentation
- 5' question