

A New Compilation Technique for SIMD Code Generation across Basic Block Boundaries

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Jan. 19, 2010

Outline

Background

- **SIMD Code Generation across Basic Block Boundaries**
- **Experimental Results**
- **Conclusion**

Background

- **Spread of digital signal processing applications**
 - CODEC, speech/image recognition, etc.
 - Require higher computational power for microprocessors

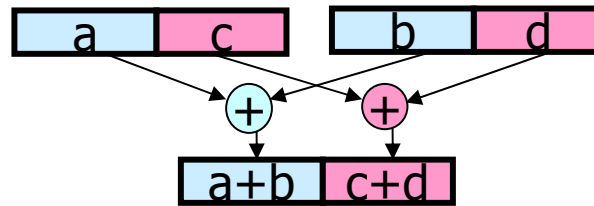


- **SIMD instruction sets**

- Provide both high performance and reasonable energy consumption

SIMD instructions

- **Perform multiple operations on data in parallel**
 - Take registers having several small data



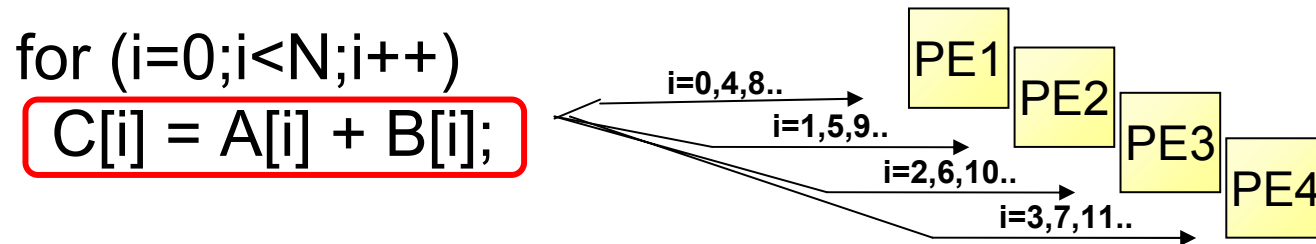
SIMD add

- **Difficulties in the utilization of SIMD instructions**
 - Need to extract parallel operations from sequential programs

SIMD Code Generation Techniques

- **For loops** [Bik, et al. 2002]

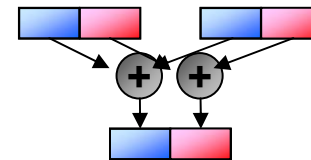
- Map same operations in different iterations into one instruction



- **For basic blocks** [larsen, et al. 2000]

- Map same operations in a basic block into one instruction

$$\begin{aligned} C[i] &= A[i] + B[i] \\ C[i+1] &= A[i+1] + B[i+1] \end{aligned}$$



SIMD Code Generation with if-conversion

- **For programs with control flow**

- Remove conditional statements using “if-conversion”

```
if(x==0) {
    c[i] = a[i]+b[i];
    c[i+1] = a[i+1]+b[i+1];
} else {
    c[i] = a[i]-b[i];
    c[i+1] = a[i+1]-b[i+1];
}
```

```
t0 = a[i]+b[i];
t1 = a[i+1]+b[i+1];
t2 = a[i]-b[i];
t3 = a[i+1]-b[i+1];

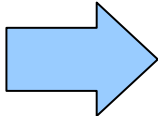
c[i] = (x==0) ? t0 : t2;
c[i+1] = (x==0) ? t1 : t3;
```

Limitations:

- If-conversion may decrease performance
- Not applicable when “if-else” statements cannot be removed by if-conversion

Challenge: SIMD Code Generation with Control Flow

```
a0=P0[0];  
a1=P0[1];  
b0=P1[0];  
b1=P1[1];  
if(mode) {  
  a0=a0+b0;  
  a1=a1+b1;  
} else {  
  a0=a0-b0;  
  a1=a1-b1;  
}  
P2[0]=a0;  
P2[1]=a1;
```



```
A=P0[0:1];  
B=P1[0:1];  
if(mode) {  
  A=A+B;  
} else {  
  A=A-B;  
}  
P2[0:1]=A;
```

Key point :
the way to keep data dependency
between basic blocks

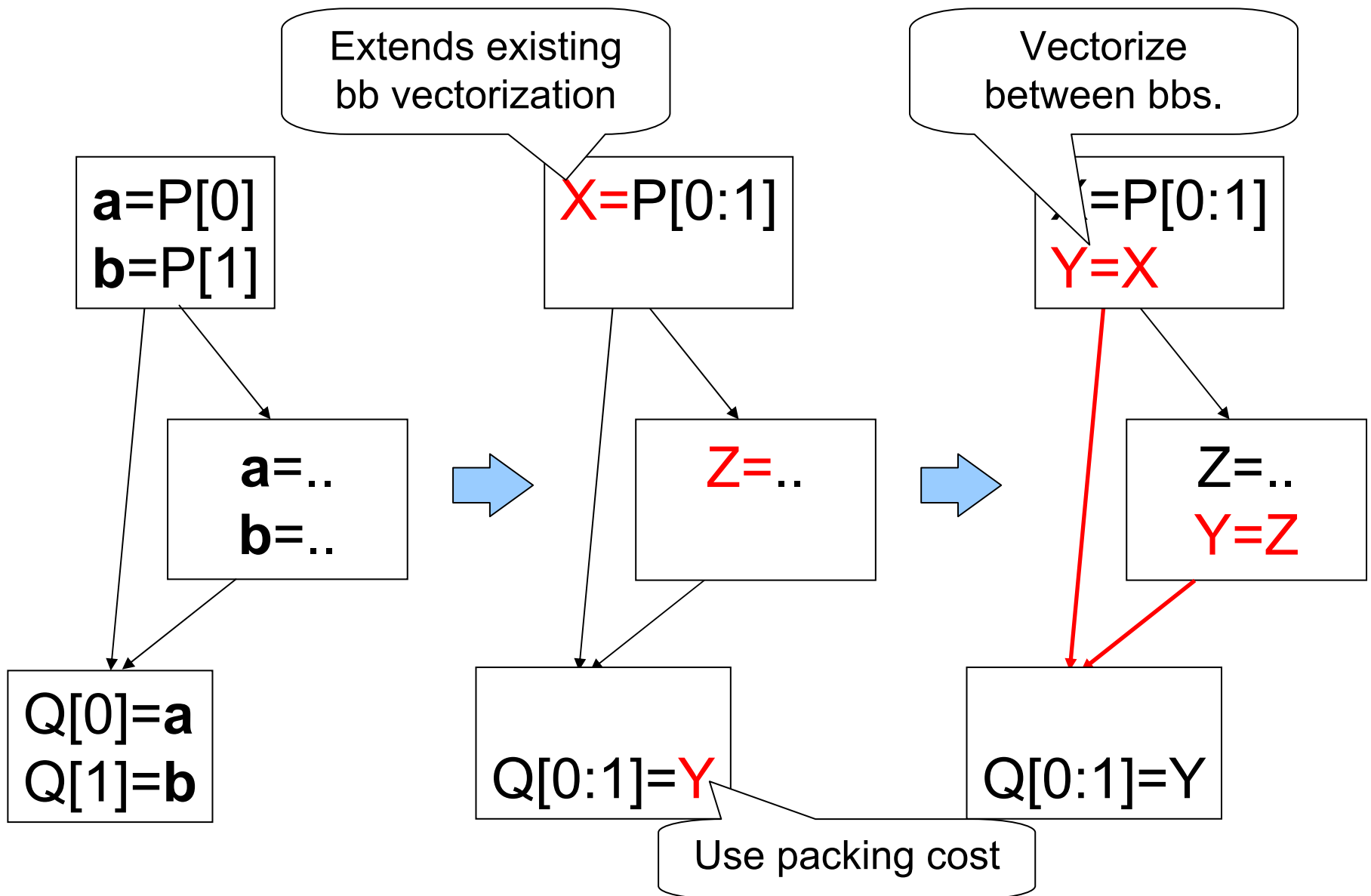
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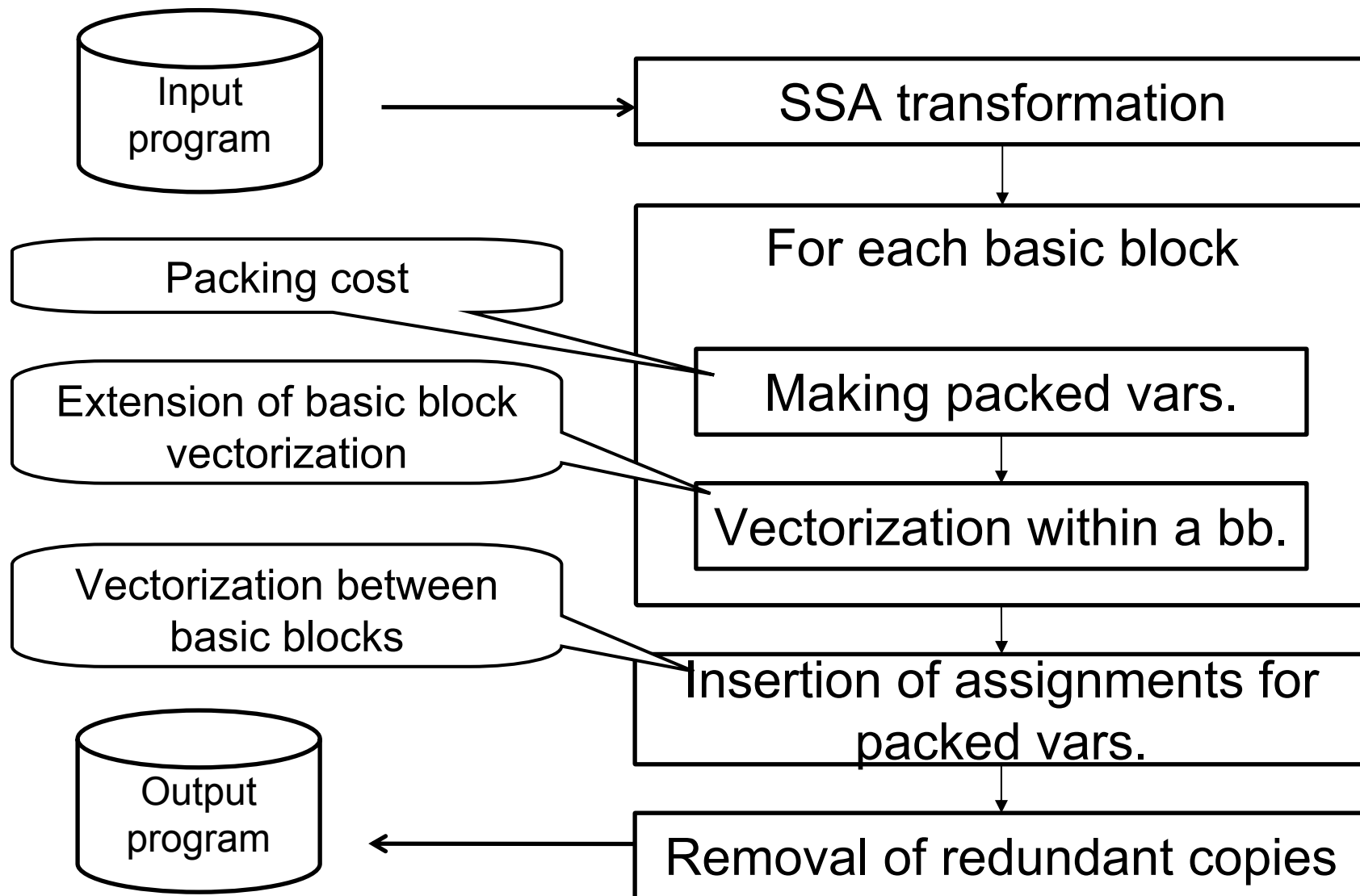
SIMD Code Generation across Basic Block Boundaries

- **Generates SIMD code without modifying control flow**
 - Deal with data dependency between basic blocks
- **Basic concepts**
 - Extension of basic block vectorization
 - Vectorization between basic blocks
 - Packing cost

Basic concepts



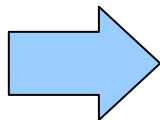
Proposed Code Generation Flow



SSA transformation

- **Remove implicit dependency due to variables**
 - Use existing SSA transformation technique

```
a0=P0[0];
a1=P0[1];
b0=P1[0];
b1=P1[1];
if(mode) {
  a0=a0+b0;
  a1=a1+b1;
} else {
  a0=a0-b0;
  a1=a1-b1;
}
P2[0]=a0;
P2[1]=a1;
```



```
a00=P0[0];
a10=P0[1];
b00=P1[0];
b10=P1[1];
if(mode) {
  a01=a00+b00;
  a11=a10+b10;
} else {
  a02=a00-b00;
  a12=a10-b10;
}
a03=∅(a01, a02);
a13=∅(a11, a12);
P2[0]=a03;
P2[1]=a13;
```

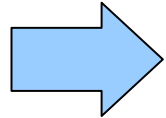
Vectorization within a Basic Block

- **Make packed variables based on packing cost, then, vectorize each basic block**

```
a00=P0[0];
a10=P0[1];
b00=P1[0];
b10=P1[1];
if(mode) {
  a01=a00+b00;
  a11=a10+b10;
} else {
  a02=a00-b00;
  a12=a10-b10;
}
a03=Ø(a01, a02);
a13=Ø(a11, a12);
P2[0]=a03;
P2[1]=a13;
```

A1

B1



```
A0=P0[0:1];
B0=P1[0:1];
if(mode) {
  A2=A1+B1
} else {
  A4=A3+B2
}
P2[0:1]=A5
```

A0 : (a00,a10)
B0 : (b00,b10)

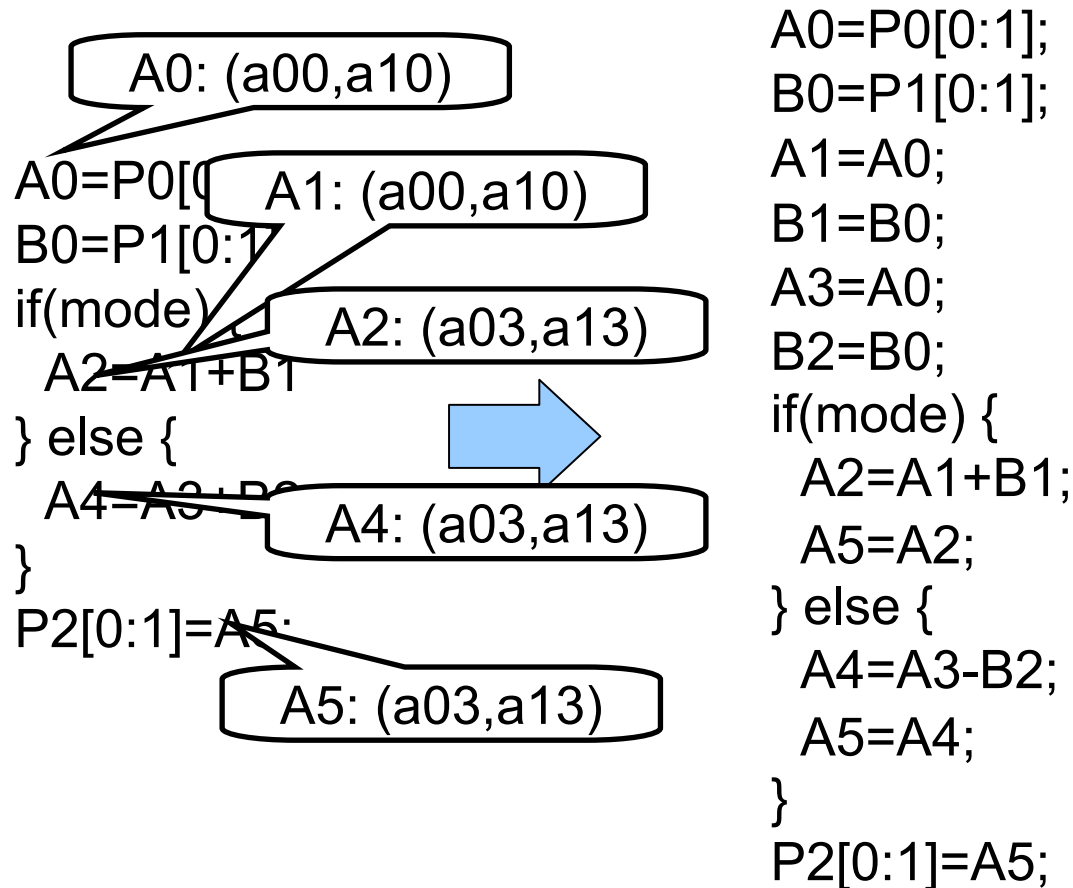
(a00,a10) : 1
(a00,b00) : 2
(a00,b10) : 2
(a10,b00) : 2
(a10,b10) : 2
(b00,b10) : 1

A1

B1

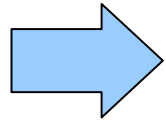
Insertion of assignments for packed vars.

- Look up packed vars having same contents
- Then, insert an assignment statement



Removal of Redundant Copies

```
A0=P0[0:1];
B0=P1[0:1];
A1=A0;
B1=B0;
A3=A0;
B2=B0;
if(mode) {
  A2=A1+B1;
A5=A2;
} else {
  A4=A3-B2;
A5=A4;
}
P2[0:1]=A5;
```



```
A0=P0[0:1];
B0=P1[0:1];
if(mode) {
  A5=A0+B0
} else {
  A5=A0-B0
}
P2[0:1]=A5;
```

- A lot of redundant statements may be inserted
- > use traditional copy removal technique

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Compiler Implementation

- **Target processor**

- Toshiba MeP with IVC2 embedded processor core
 - 3 issue VLIW architecture
 - 8 byte/4 halfword/2 word SIMD instruction set

- **Compiler**

- C compiler for MeP with IVC2 was used as a base compiler
- Proposed technique was implemented in the base compiler

Average Performance Evaluation

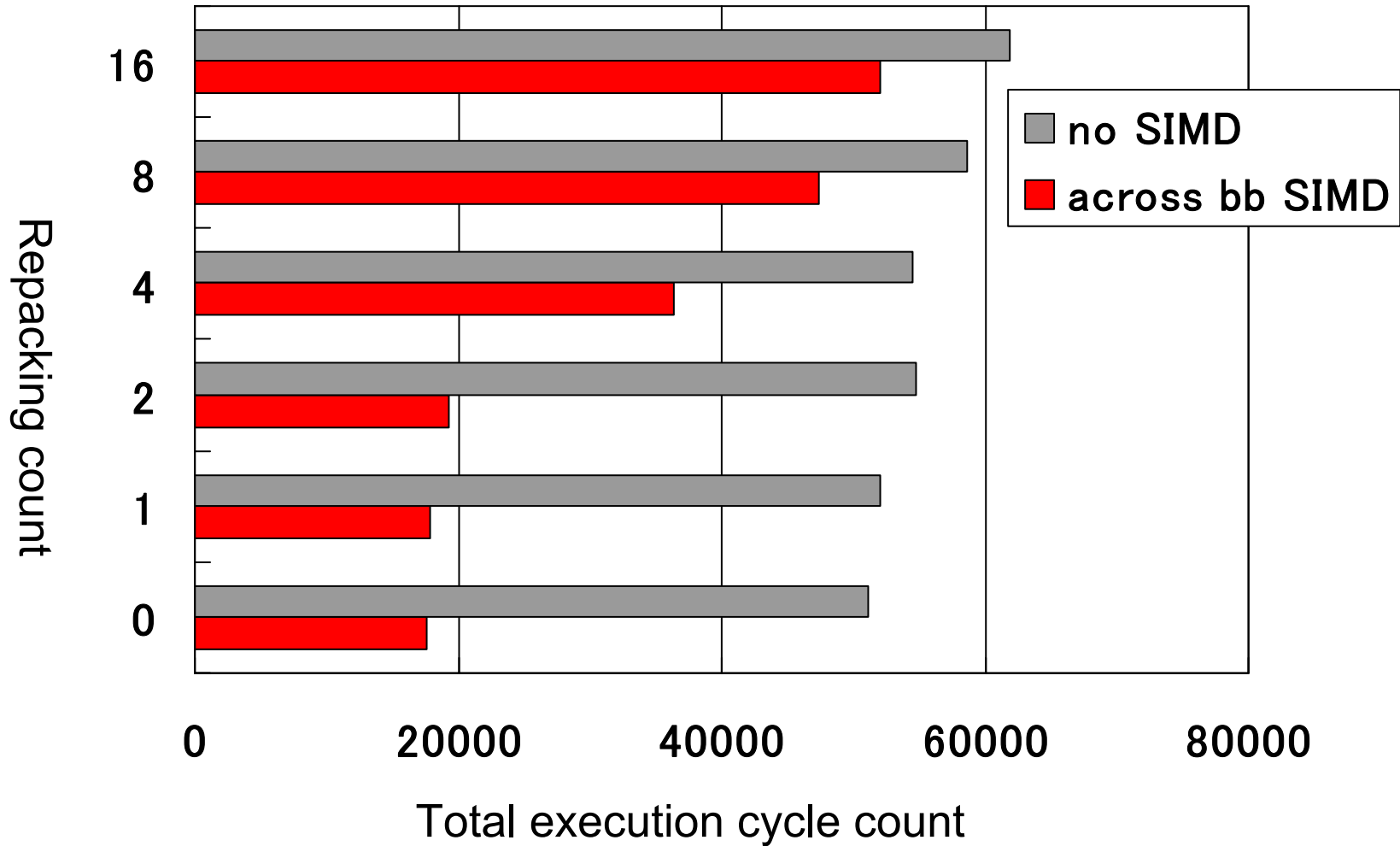
- **Randomly generated programs**

- Composed of:
 - 5 “if” statements
 - 2 SIMD operations in each basic block
 - 0-16 repacking of SIMD data
- 4 elements in a SIMD datum
- Use 10 programs for each 0,1,2,4,8,16 repacking

- **Evaluation method**

- Each program was executed 32 times
- Compared total execution cycle count over 10 programs with different repacking count

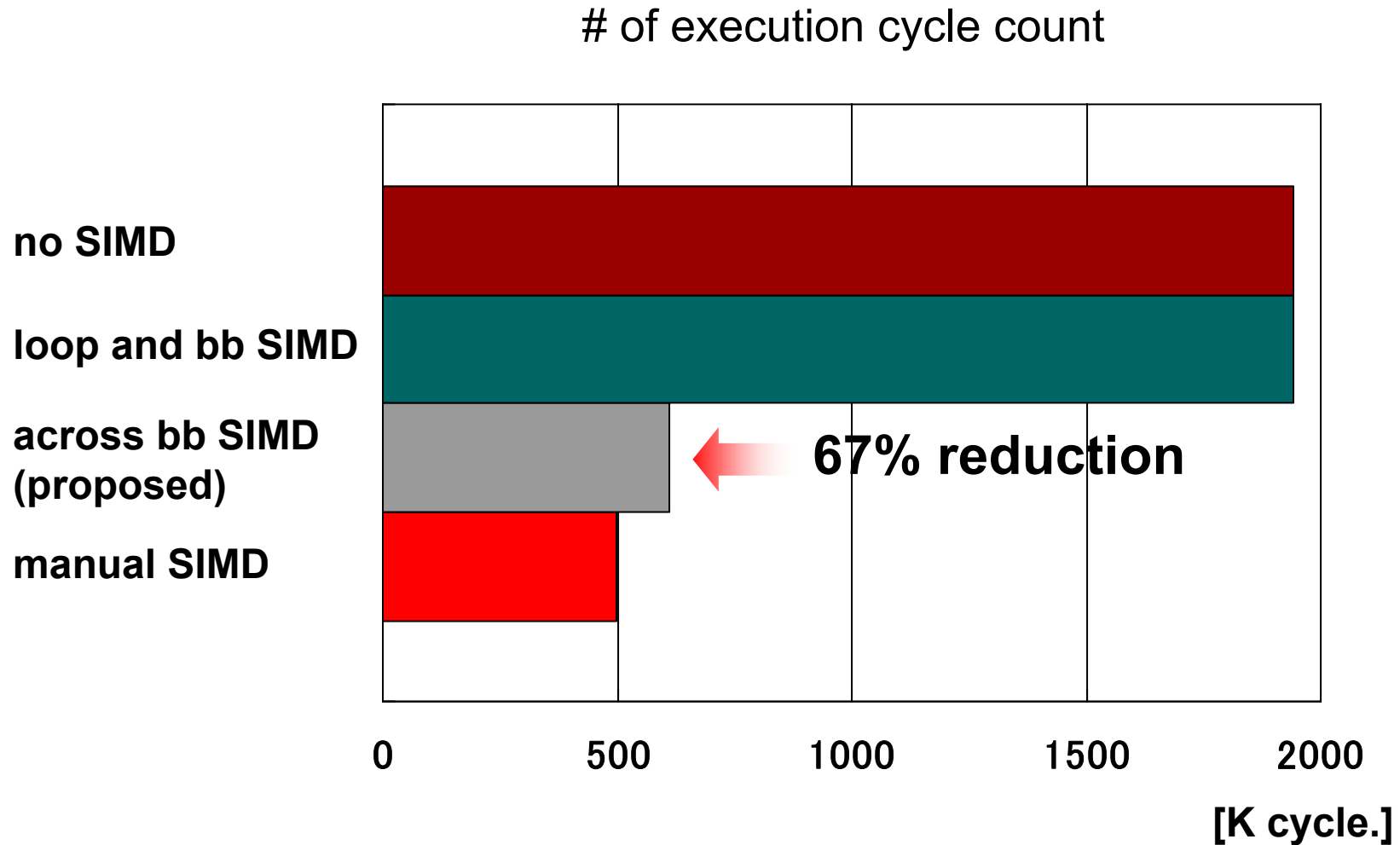
Result



Case Study: H.264 decoder inter prediction

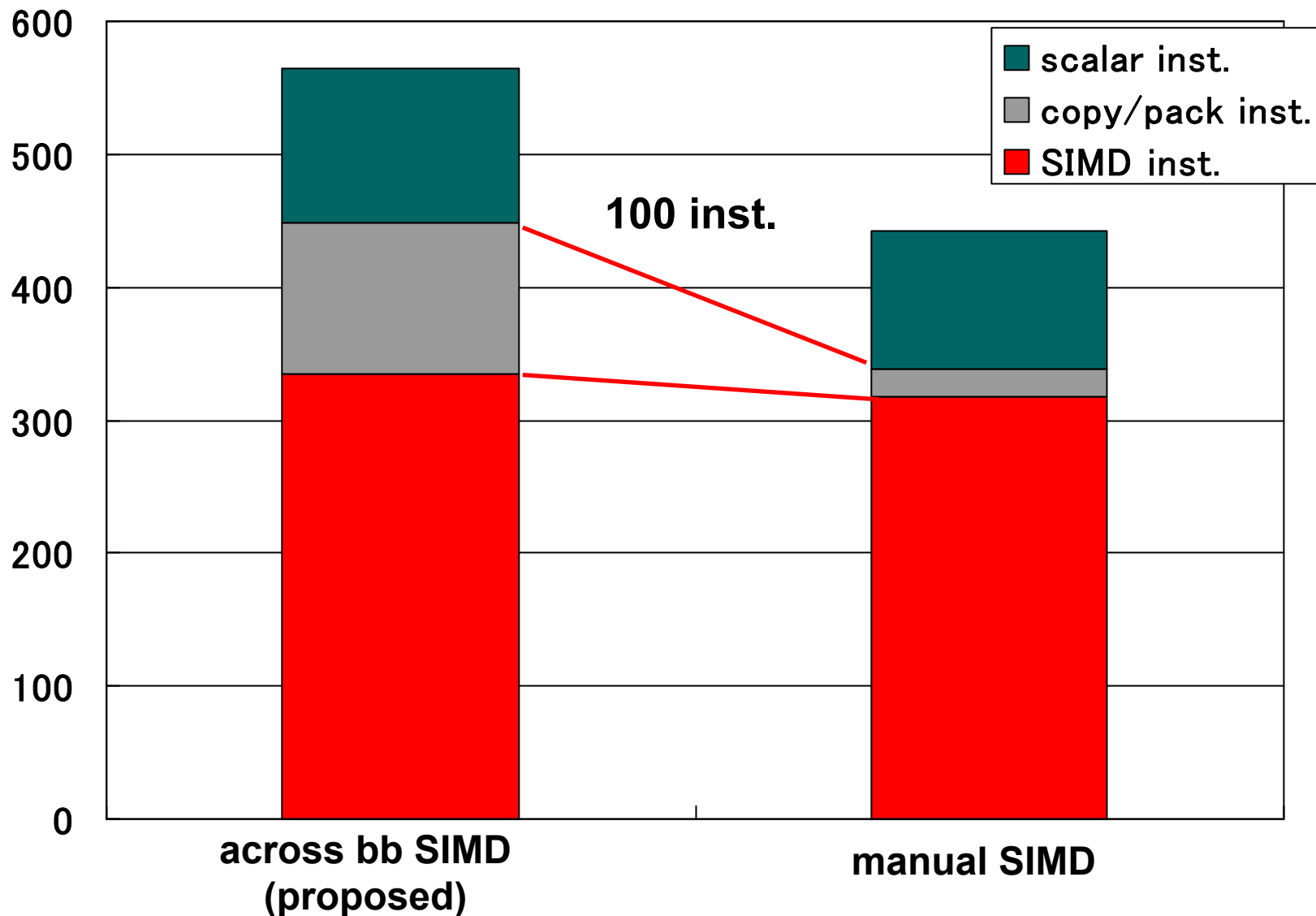
- **Reference software developed for MeP**
 - before MeP specific optimization
 - Inter prediction:
 - Consumes about 16% of execution cycles in entire decoding process
 - Contains 16 if-statements
- **Compared different compilation**
 - “no SIMD”
 - “loop and bb SIMD”
 - “across bb SIMD” (proposed)
 - “manual SIMD”

Comparison of Execution Cycle Count




Breakdown of Static Instruction Counts

[# of inst]



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Conclusion

- **New SIMD vectorization technique across basic block boundaries was proposed**
- **Performance of inter-prediction was improved**
 - Showed a 67% reduction in execution cycles
- **Future work**
 - Enhance vectorization to reduce copy/pack instructions