

Newfoundland & Labrador, Canada

4R-2



#### Symmetry-Aware TCG-Based Placement Design under Complex Multi-Group Constraints for Analog Circuit Layouts

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- Introduction
- Prior Work
- Symmetric-feasible TCG
- Symmetric Packing
- Symmetric-feasible TCG perturbation
- Experimental Results
- Conclusions



# Analog Placement

- Placement problem
  (1) Problem definition
  (2) Goal
- Symmetric constraint
  - (1) Sensitive cells placed on opposite sides
  - (2) Reduce the effect of parasitic mismatches
  - (3) Reduce the circuit sensitivity



# Symmetric Constraints

- Symmetric types Symmetric pair cell Self symmetric Symmetric group
- Flow
  - Perturb
  - Packing
  - Simulate annealing
- Comparison
  - performance
  - time complexity of perturbation
  - time complexity of packing



# **Our Contributions**

Our Contributions

- (1)introduce conditions to verify the symmetric feasibility of TCG enclosing multiple groups;
- (2) propose an efficient contour-based packing scheme;
- (3) a set of random perturbation operations with time complexity of O(n) is defined.



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Prior Work

- Previous placement applications
   -SP
  - -Transitive closure graph (TCG)
  - -TCG-S
  - -O-tree
  - -B-tree
  - -HB\*-tree



Prior Work

• TCG



• SP

Uses an ordered pair of sequence which is  $\alpha$ - and  $\beta$ sequence to encode the placement .The topological relationship between two cells *a* and *b* can be derived from an SP:

•if  $\alpha_a^{-1} < \alpha_b^{-1}$  and  $\beta_a^{-1} < \beta_b^{-1}$ , then cell *a* is to the left of cell *b*; •if  $\alpha_a^{-1} < \alpha_b^{-1}$  and  $\beta_b^{-1} < \beta_a^{-1}$ , then cell *a* is on the top of cell *b*,



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#### Symmetric feasible conditions

Let  $(G_h, G_v)$  be the TCG representation of a placement containing two symmetry groups  $\Gamma$  and  $\Phi$ . For the multi-group situation, we can define the following conditions.

**Definition 1**: For  $(a, a') \in \Gamma$ ,  $(b, b') \in \Gamma$ ,  $(c, c') \in \Phi$ , and  $(d, d') \in \Phi$ , a TCG representation is symmetric-feasible if the following four conditions are satisfied. For intra-group of  $\Gamma$  (the same for  $\Phi$ )

$$in G_h: a \models b \iff a' \models b'; \qquad (1)$$

$$in G_v: a \perp b \iff b' \perp a'; \qquad (2)$$

For inter-group between  $\Gamma$  and  $\Phi$ 

in 
$$G_h : a \models c$$
 and  $a \uparrow \models c \land \neq > d \models b$  and  $d \uparrow \models b \land$ ; (3)  
in  $G_v : a \perp c \lt \neq > c \land \perp a \land$ ; (4)

where  $\langle \neq \rangle$  denotes the two cases before and after this symbol cannot simultaneously appear in the same TCG.



#### Symmetric feasible conditions



below  $\Phi$ ; (III)  $\Gamma$  and  $\Phi$  are intermingled; (IV)  $\Phi$  is placed within  $\Gamma$ 

• Lemma: Any placement containing multiple symmetry groups can be represented with a symmetric-feasible TCG.

# Symmetric TCG



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Fig. 2. (a) TCG; (b) topological order; (c) symmetric placement with separate self-symmetric halves; (d) final symmetric placement



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# Symmetric Packing Flow

Beg	zin
1	Construct the topological order of the TCG;
2	Generate packSeq, which contains only the symmetric cells;
3	Use the dummy axis cells to separate packSeq into sub-sequences;
4	Create a list made up of the sub-sequences;
5	Do initial packing for the first sub-sequence, update the axis position
	X <sub>axis</sub> ;
6	For (the <i>list</i> is not empty)
7	Do initial packing for the second sub-sequence;
8	Follow <i>packSeq</i> one by one. Compare the Y coordinates of each two cells in one symmetric pair such as $(a, a')$ , and make $Y_a = Y_{a'} = max(Y_a, Y_{a'})$ and shift the packing symmetric cells that have vertical relationship with the shifted cells;
9	Follow <i>packSeq</i> one by one, for every two cells in one symmetric pair such as $(a, a')$ ; calculate $\Delta X_a =  X_a - X_{axis} $ and $\Delta X_{a'} =  X_{a'} - X_{axis} $ . Then shift one symmetric cell to make $\Delta X_a = \Delta X_{a'} = max (\Delta X_a)$ $\Delta X_{a'}$ . Also tune-up the corresponding cells with the same $\Delta X_a$ :
10	Do final packing for the first and second sub-sequences:
11	Remove the first sub-sequence from the sub-sequence list:
12	End for
13	Consider all symmetric cells as preplaced cells and final-pack the whole sequence;
14	Post-process self-symmetric cells to merge two split parts to one unit;
En	d
E:	2 Commentation and the floor

Fig. 3. Symmetric packing flow



### Symmetric Packing



- Step 1:packing preparation; Step 2:initial packing (a-d); Step 3:tune-up operation on symmetric cells (f-h); Step 4:final packing (i)
- Two Schemes : (1)general flow (2) simplified flow
- Complexity is  $O(g \cdot m \cdot lgm)$ .
- Lemma: Any symmetric-feasible TCG containing multiple symmetry groups can be packed to a symmetric placement in polynomial time.



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### Symmetric TCG Perturbation

- A. Cluster Edge Move-Reverse and Edge Move Operations
- in-in (in-out) degree
- in-in (or in-out) cluster
- Cluster edge move-reverse (or cluster edge move) operation
- All the operations in part B can be done with Cluster edge movereverse (or cluster edge move) operation.
- Lemma: Without losing TCG transitive-closure property, the topological relationship between any two vertices in a TCG can be modified by a cluster edge move or cluster edge move-reverse operation, which takes O(n) time.



### Symmetric TCG Perturbation

- **B.** Five Perturbation Operations
- Vertex Rotation
- Symmetric Swap

This operation is that one vertex in a symmetric pair swaps position with its symmetric counterpart.

• Symmetric-Cell Move

This operation is to change the horizontal or vertical relationship between symmetric vertices within one symmetry group or from different symmetry groups.

• Asymmetric-Cell Move

This operation is to perturb the topology relationship between an asymmetric vertex and other vertices.

• Symmetry-Group Move



#### Symmetric TCG Perturbation



**Theorem 1**: Using operations above, the solution space of the symmetric-feasible TCGs can be fully explored. Each operation takes at almost O(n) time, where *n* is the number of cells.



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### Experimental Results

- We implement this symmetry TCG scheme based on a simulated annealing algorithm. The program was coded in C++ under Linux operating system and tested on a 2.1 GHz PC with several benchmarks.
- Our cost function for evaluation is obtained from the following equation:

 $Cost = \alpha_{area} \cdot Area + \sum_i \beta_i WireLength_i$ 

- Test Cases
  - (1) MCNC benchmarks
  - (2) biasynth\_2p4g of 65 cells and lnamixbias\_2p4g of 110 cells
  - (3) modified 65 cells and 110 cells, OTA



#### Experimental Results

I.MCNC BENCHMARKS																
Circuit	cell s	grou ps		Abs		SP		SPWD		SP+LP	HB	HB*-tree		CG-1	S-TCG-2	
ante	9	2	0	Cost 12		0.7%	112.2%		107.7%		106.4%	101.2%		100	.0%	50.74
apic			Т	ime 10		00%	767%		466.7%		400.0%	100%		119	9%	3
	22		0	Cost 1		8.9%	109.3%		108.5%		107.4%	103.7%		99.	0%	1.33
amiss	33		Т	ïme	1670%		1190%		243.4%		354.8%	83.8%		121	1%	62
ami49	10	2	0	Cost 12		9.1% 107.9		9% 108.7%			107.1%	102.5%		99.2	2%	42.21
	49		Т	Time 22		.50% 136		0%	201.0%		228.0%	90.4%		118	3%	107
							]	I.CIF	<b>RCUIT TE</b>	ST						
Circuit	t Ce	ell gr s p	ou s			Al	bs		SP		SPWD	SP+LP		HB*·	-tree	S-TCG-2
biasynt	h	5	,	Co	st	127%		109.8%			107.4%	104%		102.	1%	5.42
_2p	C	0 3	)	Tir	ne	215	0%		611%		154%	221%		89.3	3%	131
lnamixbi	as 1	10 4		Co	ost	119	)%	1	14.6%		109%	109%		103.	.2%	51.41
_2p	1	10 .	)	Time		243	0%	1	130%		543%	920%		82.5	5%	287
mod_	6	5	2	Co	ost	123	3% 1		07.8%		107.1%	10	6%	104.	.2%	5.53
biasynt	h		,	Tir	ne	213	7%		621%		153%	224%		88.	5%	144
mod_lna	mi   1	10		Cost		120%		114.1%		111%		11	0%	105.	.1%	52.23
xbias			,	Tir	me 244		0% 1		1110%		572%	934%		86.3	3%	294
				Co	ost 119		)% 1		14.9%		107.2%	105%		105.	.5%	27.53
			,	Time		2084%		724%			167%	207%		90.4	4%	263



#### **Experimental Results**



Fig. 7. Final placement of circuit lnamixbias\_2p4g



Fig. 8. Final placement of circuit OTA



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### Conclusions

- In this paper, we introduce a scheme to handle the symmetric constraints in the analog layout placement.
  - An efficient perturbation strategy was proposed to achieve random state conversion in O(n) time without losing TCG symmetric-feasibility and validity.
  - Although the HB\*-tree scheme slightly outperforms our approach in terms of area in some larger circuits, our proposed approach is able to cover any placement configurations.



#### Thank you and Questions?