Symmetry-Aware TCG-Based Placement Design under
Complex Multi-Group Constraints for Analog Circuit Layouts

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## Outline

- Introduction
- Prior Work
- Symmetric-feasible TCG
- Symmetric Packing
- Symmetric-feasible TCG perturbation
- Experimental Results
- Conclusions


## Analog Placement

- Placement problem
(1) Problem definition
(2) Goal
- Symmetric constraint
(1) Sensitive cells placed on opposite sides
(2) Reduce the effect of parasitic mismatches
(3) Reduce the circuit sensitivity


## Symmetric Constraints

- Symmetric types

Symmetric pair cell
Self symmetric
Symmetric group

- Flow

Perturb
Packing
Simulate annealing

- Comparison
performance
time complexity of perturbation
time complexity of packing


## Our Contributions

Our Contributions

- (1)introduce conditions to verify the symmetric feasibility of TCG enclosing multiple groups;
- (2) propose an efficient contour-based packing scheme;
- (3) a set of random perturbation operations with time complexity of $O(n)$ is defined.


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## Prior Work

- Previous placement applications
-SP
-Transitive closure graph (TCG)
-TCG-S
-O-tree
-B-tree
-HB*-tree


## Prior Work

- TCG


- SP

Uses an ordered pair of sequence which is $\alpha$ - and $\beta$ sequence to encode the placement .The topological relationship between two cells $a$ and $b$ can be derived from an SP:
-if $\alpha_{a}^{-1}<\alpha_{b}{ }^{-1}$ and $\beta_{a}^{-1}<\beta_{b}{ }^{-1}$, then cell $a$ is to the left of cell $b$; -if $\alpha_{a}^{-1}<\alpha_{b}^{-1}$ and $\beta_{b}^{-1}<\beta_{a}^{-1}$, then cell $a$ is on the top of cell $b$,

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## Symmetric feasible conditions

Let $\left(G_{h}, G_{v}\right)$ be the TCG representation of a placement containing two symmetry groups $\Gamma$ and $\Phi$. For the multi-group situation, we can define the following conditions.

Definition 1: For $\left(a, a^{\prime}\right) \in \Gamma,\left(b, b^{\prime}\right) \in \Gamma,\left(c, c^{\prime}\right) \in \Phi$, and $\left(d, d^{\prime}\right) \in \Phi$, a TCG representation is symmetric-feasible if the following four conditions are satisfied. For intra-group of $\Gamma$ (the same for $\Phi$ )
in $G_{h}: a \vdash b\langle\neq\rangle a^{\prime} 卜 b^{\prime}$;
in $G_{v}: a \perp b\langle\neq\rangle b^{\prime} \perp_{a^{\prime}}$;
For inter-group between $\Gamma$ and $\Phi$ in $G_{h}: a \mid-c$ and $a^{\prime} \mid c^{\prime}\langle\neq\rangle+b$ and $d^{\prime} 卜 b^{\prime} ;$
in $G_{v}: a^{\perp} c<\neq>c^{\prime} \perp^{\prime} a^{\prime} ;$
where $\langle\neq\rangle$ denotes the two cases before and after this symbol cannot simultaneously appear in the same TCG.

## Symmetric feasible conditions


(I)

(II)

(III)

(IV)

Fig. 1. (I) group $\Gamma$ is placed at the left of group $\Phi$; (II) $\Gamma$ is placed below $\Phi$; (III) $\Gamma$ and $\Phi$ are intermingled; (IV) $\Phi$ is placed within $\Gamma$

- Lemma: Any placement containing multiple symmetry groups can be represented with a symmetric-feasible TCG.


## Symmetric TCG


(a)
( $\left.a, d, c, b_{l}, D N 1, c^{\prime}, e, a^{\prime}, b_{r}, g, h, f, D N 2, g^{\prime}, f^{\prime}\right)$
(b)

(c)

(d)

Fig. 2. (a) TCG; (b) topological order; (c) symmetric placement with separate self-symmetric halves; (d) final symmetric placement

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## Symmetric Packing Flow

## Begin

1 Construct the topological order of the TCG;
2 Generate packSeq, which contains only the symmetric cells;
3 Use the dummy axis cells to separate packSeq into sub-sequences;
4 Create a list made up of the sub-sequences;
5 Do initial packing for the first sub-sequence, update the axis position $X_{\text {ausis }}$;
6 For (the list is not empty)
7 Do initial packing for the second sub-sequence;
8 Follow packSeq one by one. Compare the Y coordinates of each two cells in one symmetric pair such as ( $a, a$ ), and make $Y_{a}=Y_{a^{\prime}}=$ $\max \left(Y_{a}, Y_{a}\right)$ and shift the packing symmetric cells that have vertical relationship with the shifted cells;
9 Follow packSeq one by one, for every two cells in one symmetric pair such as $(a, a)$; calculate $\Delta X_{a}=\left|X_{a}-X_{\alpha x i s}\right|$ and $\Delta X_{a^{\prime}}=\left|X_{a^{\prime}}-X_{a x i s}\right|$. Then shift one symmetric cell to make $\Delta X_{a}=\Delta X_{a^{\prime}}=\max \left(\Delta X_{\infty}\right.$ $\Delta X_{a}$ ). Also tune-up the corresponding cells with the same $\Delta X_{a}$;
10 Do final packing for the first and second sub-sequences;
11 Remove the first sub-sequence from the sub-sequence list,
12 End for
13 Consider all symmetric cells as preplaced cells and final-pack the whole sequence;
14 Post-process self-symmetric cells to merge two split parts to one unit; End
Fig. 3. Symmetric packing flow

## Symmetric Packing



Fig. 4. Examples of the packing

- Step 1:packing preparation; Step 2:initial packing (a-d); Step 3:tune-up operation on symmetric cells (f-h); Step 4:final packing (i)
- Two Schemes : (1)general flow (2) simplified flow
- Complexity is $O(g \cdot m \cdot \operatorname{lgm})$.
- Lemma: Any symmetric-feasible TCG containing multiple symmetry groups can be packed to a symmetric placement in polynomial time.


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## Symmetric TCG Perturbation

A. Cluster Edge Move-Reverse and Edge Move Operations

- $\quad$ in-in (in-out) degree
- in-in (or in-out) cluster
- Cluster edge move-reverse (or cluster edge move) operation
- All the operations in part B can be done with Cluster edge movereverse (or cluster edge move) operation.
- Lemma: Without losing TCG transitive-closure property, the topological relationship between any two vertices in a TCG can be modified by a cluster edge move or cluster edge move-reverse operation, which takes $O(n)$ time.


## Symmetric TCG Perturbation

B. Five Perturbation Operations

- Vertex Rotation
- Symmetric Swap

This operation is that one vertex in a symmetric pair swaps position with its symmetric counterpart.

- Symmetric-Cell Move

This operation is to change the horizontal or vertical relationship between symmetric vertices within one symmetry group or from different symmetry groups.

- Asymmetric-Cell Move

This operation is to perturb the topology relationship between an asymmetric vertex and other vertices.

- Symmetry-Group Move


## Symmetric TCG Perturbation



Fig. 6. Example of asymmetric move

Theorem 1: Using operations above, the solution space of the symmetric-feasible TCGs can be fully explored. Each operation takes at almost $O(n)$ time, where $n$ is the number of cells.

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## Experimental Results

- We implement this symmetry TCG scheme based on a simulated annealing algorithm. The program was coded in $\mathrm{C}++$ under Linux operating system and tested on a 2.1 GHz PC with several benchmarks.
- Our cost function for evaluation is obtained from the following equation:

$$
\text { Cost }=\alpha_{\text {area }} \text { Area }+\sum_{i} \beta_{i} \text { WireLength } h_{i}
$$

- Test Cases
(1) MCNC benchmarks
(2) biasynth_2p4g of 65 cells and lnamixbias_2p4g of 110 cells
(3) modified 65 cells and 110 cells, OTA


## Experimental Results

I.MCNC BENCHMARKS

| Circuit | $\begin{gathered} \text { cell } \\ \mathrm{s} \end{gathered}$ | $\begin{gathered} \text { grou } \\ \text { ps } \end{gathered}$ |  | Abs | $\boldsymbol{S P}$ | SPWD | $S P+L P$ | HB**tree | S-TCG-1 | S-TCG-2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| apte | 9 | 2 | Cost | 120.7\% | 112.2\% | 107.7\% | 106.4\% | 101.2\% | 100.0\% | 50.74 |
|  |  |  | Time | 1000\% | 767\% | 466.7\% | 400.0\% | 100\% | 119\% | 3 |
| ami33 | 33 | 2 | Cost | 118.9\% | 109.3\% | 108.5\% | 107.4\% | 103.7\% | 99.0\% | 1.33 |
|  |  |  | Time | 1670\% | 1190\% | 243.4\% | 354.8\% | 83.8\% | 121\% | 62 |
| ami49 | 49 | 2 | Cost | 129.1\% | 107.9\% | 108.7\% | 107.1\% | 102.5\% | 99.2\% | 42.21 |
|  |  |  | Time | 2250\% | 1360\% | 201.0\% | 228.0\% | 90.4\% | 118\% | 107 |

II.CIRCUIT TEST

| Circuit | $\begin{array}{\|c} \hline \text { cell } \\ \mathrm{s} \\ \hline \end{array}$ | $\begin{gathered} \text { grou } \\ \text { ps } \end{gathered}$ |  | Abs | SP | SPWD | $S P+L P$ | HB**tree | S-TCG-2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| biasynth$\ldots 2 \mathrm{p}$ | 65 | 3 | Cost | 127\% | 109.8\% | 107.4\% | 104\% | 102.1\% | 5.42 |
|  |  |  | Time | 2150\% | 611\% | 154\% | 221\% | 89.3\% | 131 |
| $\begin{gathered} \text { lnamixbias } \\ 2 p \end{gathered}$ | 110 | 5 | Cost | 119\% | 114.6\% | 109\% | 109\% | 103.2\% | 51.41 |
|  |  |  | Time | 2430\% | 1130\% | 543\% | 920\% | 82.5\% | 287 |
| mod_ | 65 | 3 | Cost | 123\% | 107.8\% | 107.1\% | 106\% | 104.2\% | 5.53 |
| biasynth |  |  | Time | 2137\% | 621\% | 153\% | 224\% | 88.5\% | 144 |
| mod_lnami xbias | 110 | 5 | Cost | 120\% | 114.1\% | 111\% | 110\% | 105.1\% | 52.23 |
|  |  |  | Time | 2440\% | 1110\% | 572\% | 934\% | 86.3\% | 294 |
| OTA | 69 | 5 | Cost | 119\% | 114.9\% | 107.2\% | 105\% | 105.5\% | 27.53 |
|  |  |  | Time | 2084\% | 724\% | 167\% | 207\% | 90.4\% | 263 |

## Experimental Results



Fig. 7. Final placement of circuit lnamixbias_2p4g


Fig. 8. Final placement of circuit OTA

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## Conclusions

- In this paper, we introduce a scheme to handle the symmetric constraints in the analog layout placement.
- An efficient perturbation strategy was proposed to achieve random state conversion in $O(n)$ time without losing TCG symmetric-feasibility and validity.
- Although the HB*-tree scheme slightly outperforms our approach in terms of area in some larger circuits, our proposed approach is able to cover any placement configurations.


## Thank you

 and
## Questions?

