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Adjustable Contiguity of Run-Time Task Allocation in Networked Many-Core Systems

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Outline

☞ Introduction.

☞ Motivation.

☞ Contiguity Adjustable Square Allocation.

☞ Results.

Introduction

Introduction

Many-Core Systems

- Connected via NoC

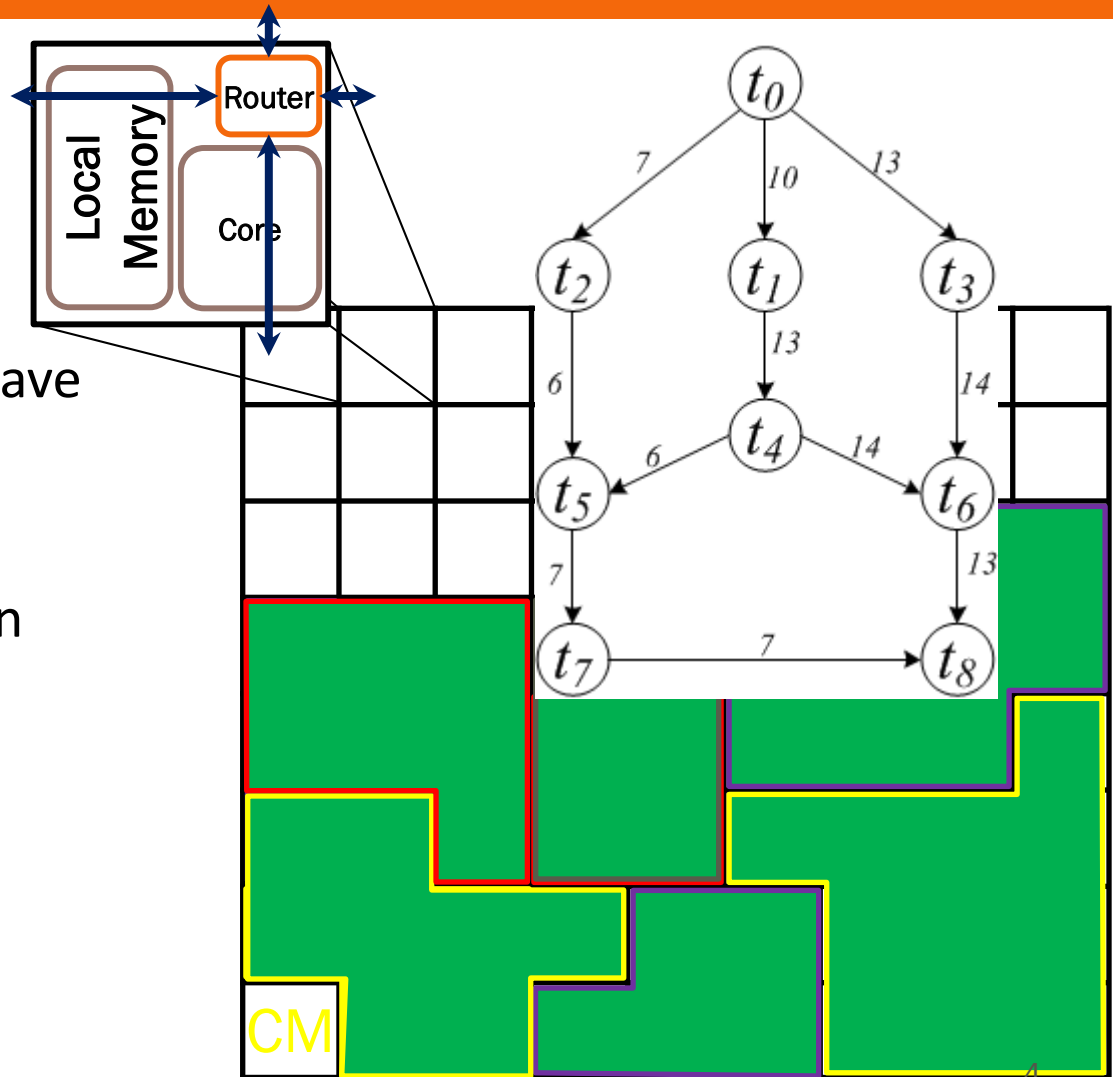
Applications: set of communicating tasks.

Applications enter and leave the system at Run-Time.

Question: How to map an Application?

Answer: Contiguously

- Less network power
- Less network congestion



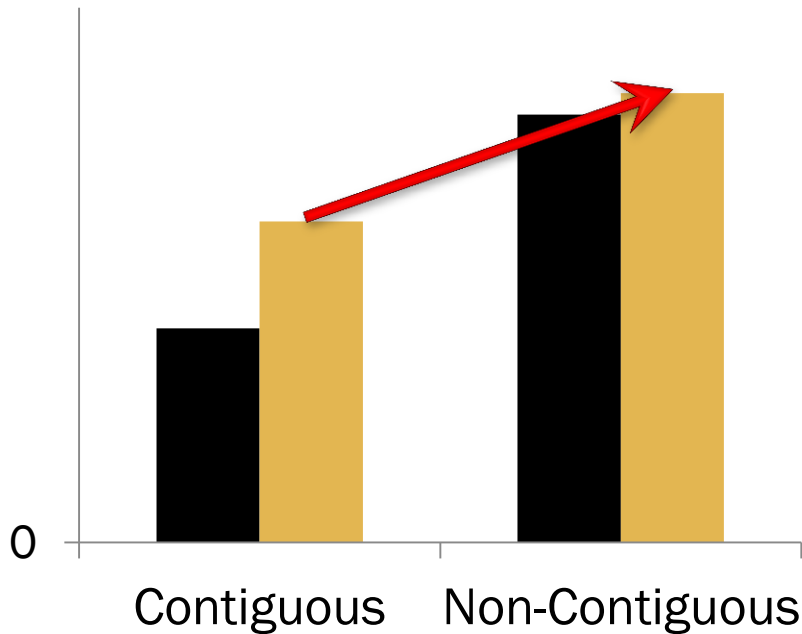
Motivation

Motivation

Store-and-Forward

Forced to limit the system to **only contiguous** allocations.

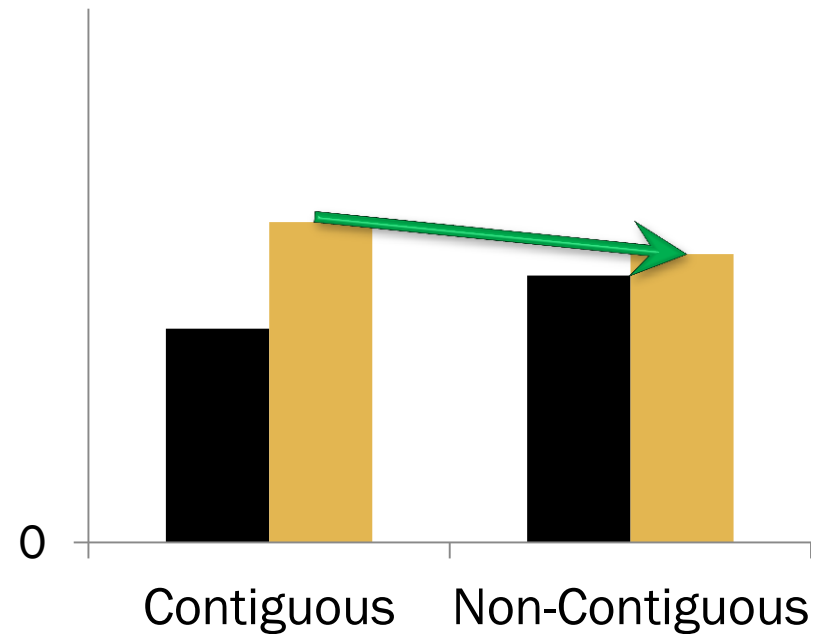
■ Execution ■ Turn-Around



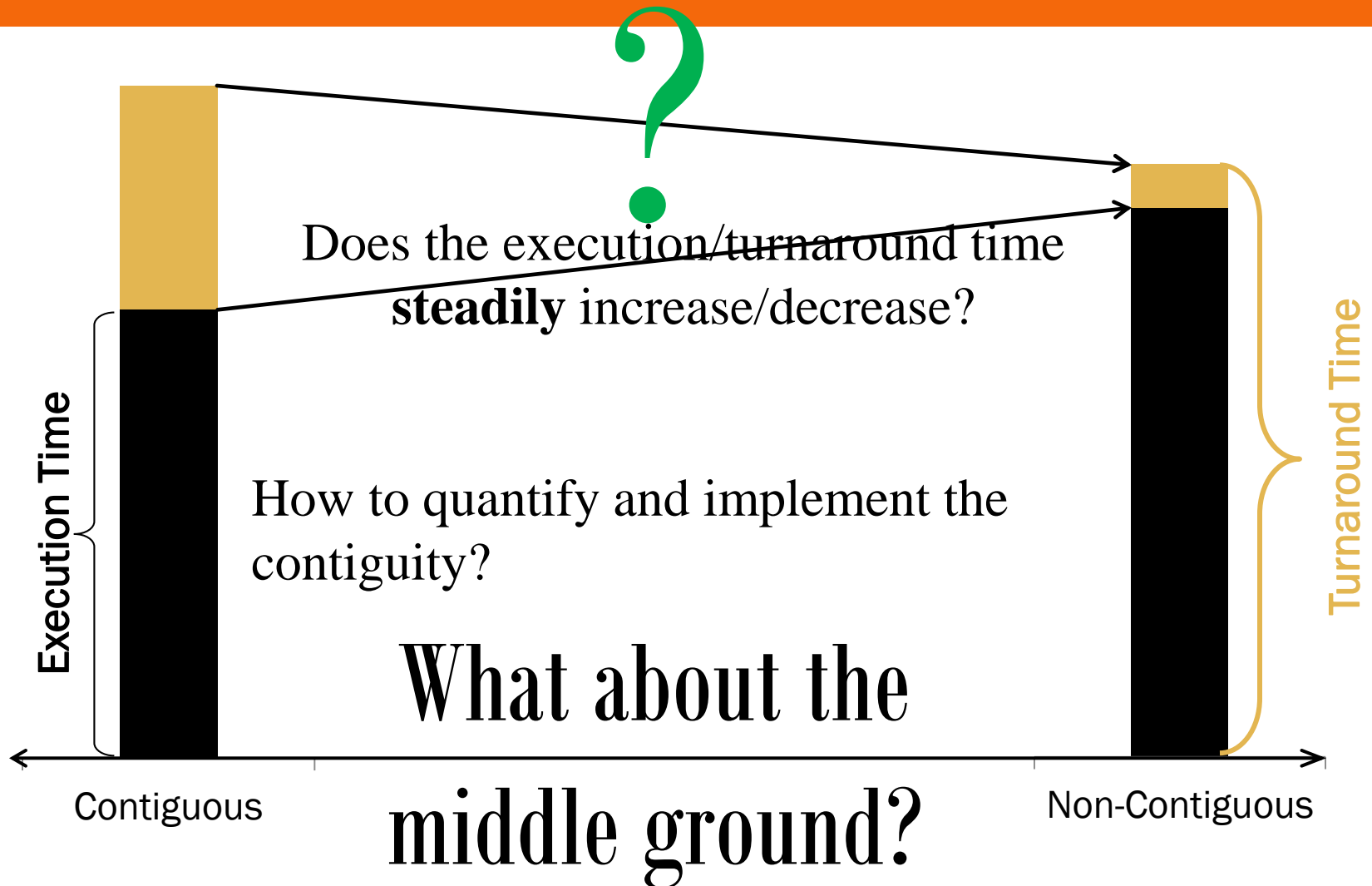
Wormhole

Non-contiguous allocation: **higher throughput** with **more network power**

■ Execution ■ Turn-Around



Motivation



CASqA ^{α}

Contiguity Aadjustable Square Allocation

CASqA^α

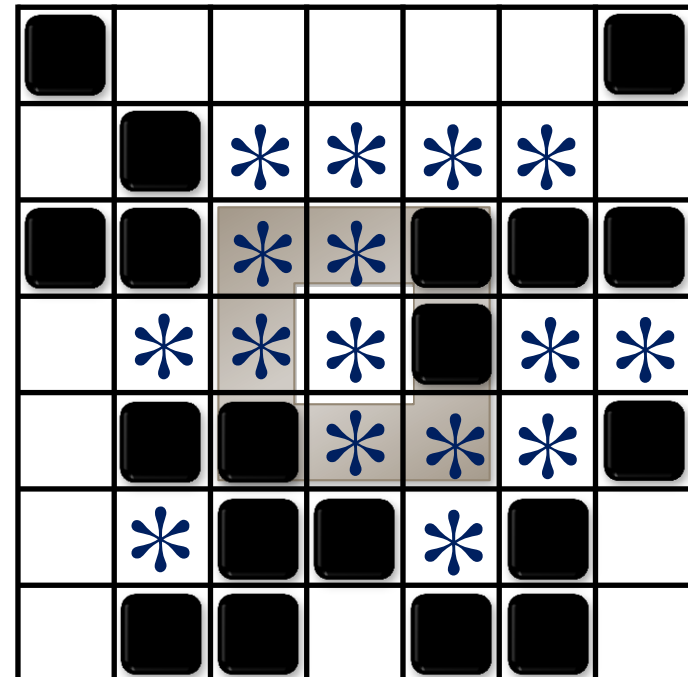
∞ $0.0 \leq \alpha \leq 1.0$:

- level of “desired contiguity” / “allowed

α adaptively determines the **maximum exploration radius**.

For an application with $|T|$ tasks:

1. Start from a first node
2. until $|T|$ available nodes are found:
 - Increase the **exploration radius** (allocation dispersion)



CASqA^α

∞ R_{max} :

- Maximum allowed exploration radius.
- Initial value:

$$R_{max} = \left\lfloor \frac{\lceil \sqrt{|T|} \rceil}{2} \right\rfloor$$

∞ τ (expansion threshold):

- Initial value: $|T| \times \alpha$.
- if less than τ nodes are needed: $R_{max}++$
- $\tau \times = \alpha$ on each increase.

∞ $\alpha=0 \rightarrow \tau=0 \rightarrow$ don't explore beyond initial R_{max} .

∞ $\alpha=1 \rightarrow \tau=|T| \rightarrow$ expand exploration until allocated.

CASqA^α- example

A 9-task application enters the system with the following run-time configuration.

∞ Initial

$$R_{max} = \left\lceil \frac{\lceil \sqrt{|9|} \rceil}{2} \right\rceil = 1$$

If $\alpha \geq 0.56$
 $\rightarrow \tau > 5$

nodes.

∞ $\alpha=0.5$, initial $\tau=4.5$.

R_{max}	τ	#alloc	#unalloc	expand
1	4.5	4	5	NO

■						■
	■	*	*			
■	■	*	*	■	■	■
*	*	*	*	■		
*	■	■	*	*		■
*	*	■	■		■	
*	■	■		■	■	

CASqA ^{α} - details

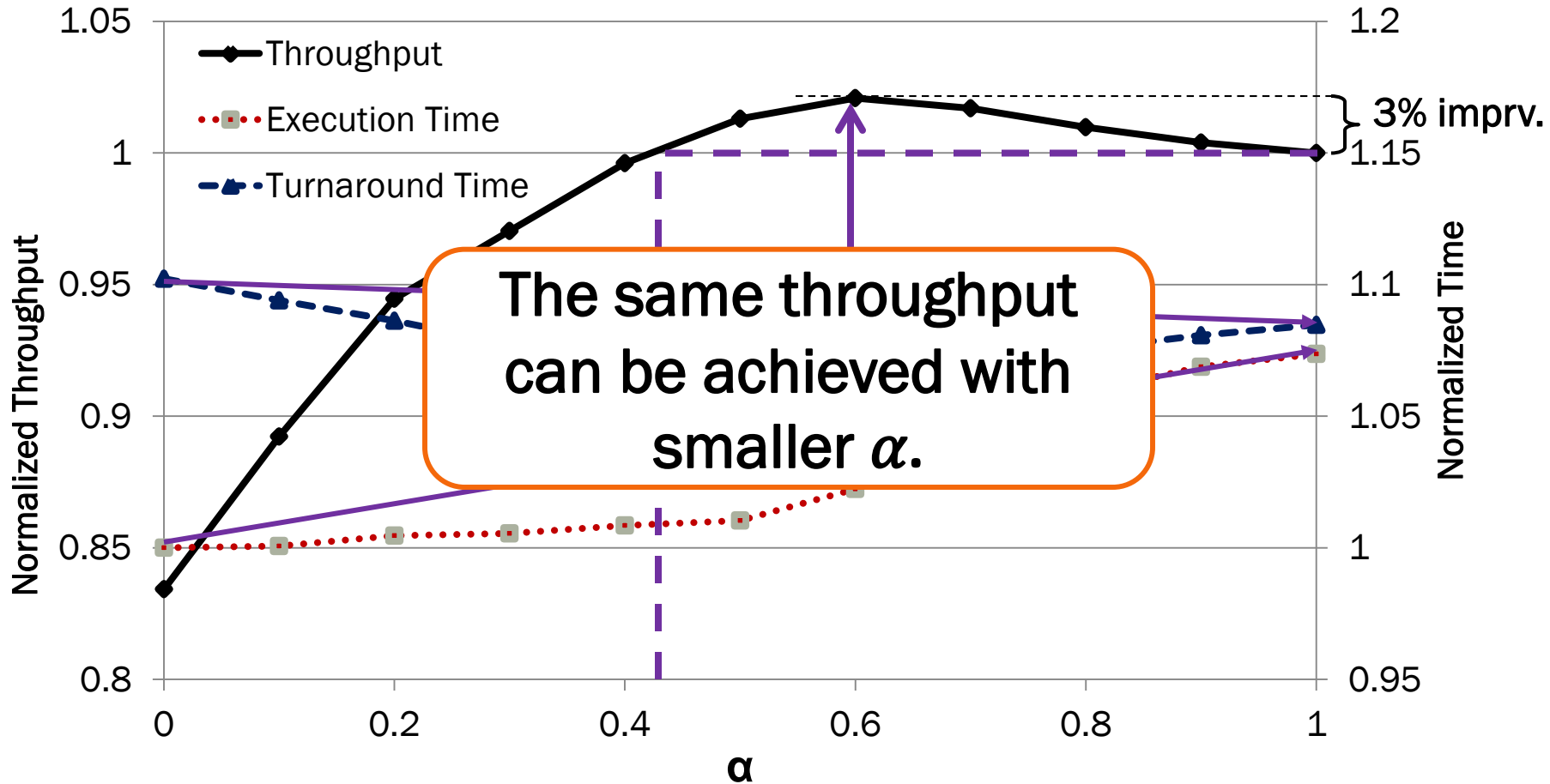
- ✎ ICEB is defined as a unified metric to model both congestion and energy dissipation of the network.
- ✎ ICEB metric is used by CASqA to arrange task placement within a selected area.
- ✎ First node of allocation is selected using our previous approach, SHiC.
- ✎ The source code can be downloaded at: :
<http://users.utu.fi/mofana/CASqA.html>

Experimental Results

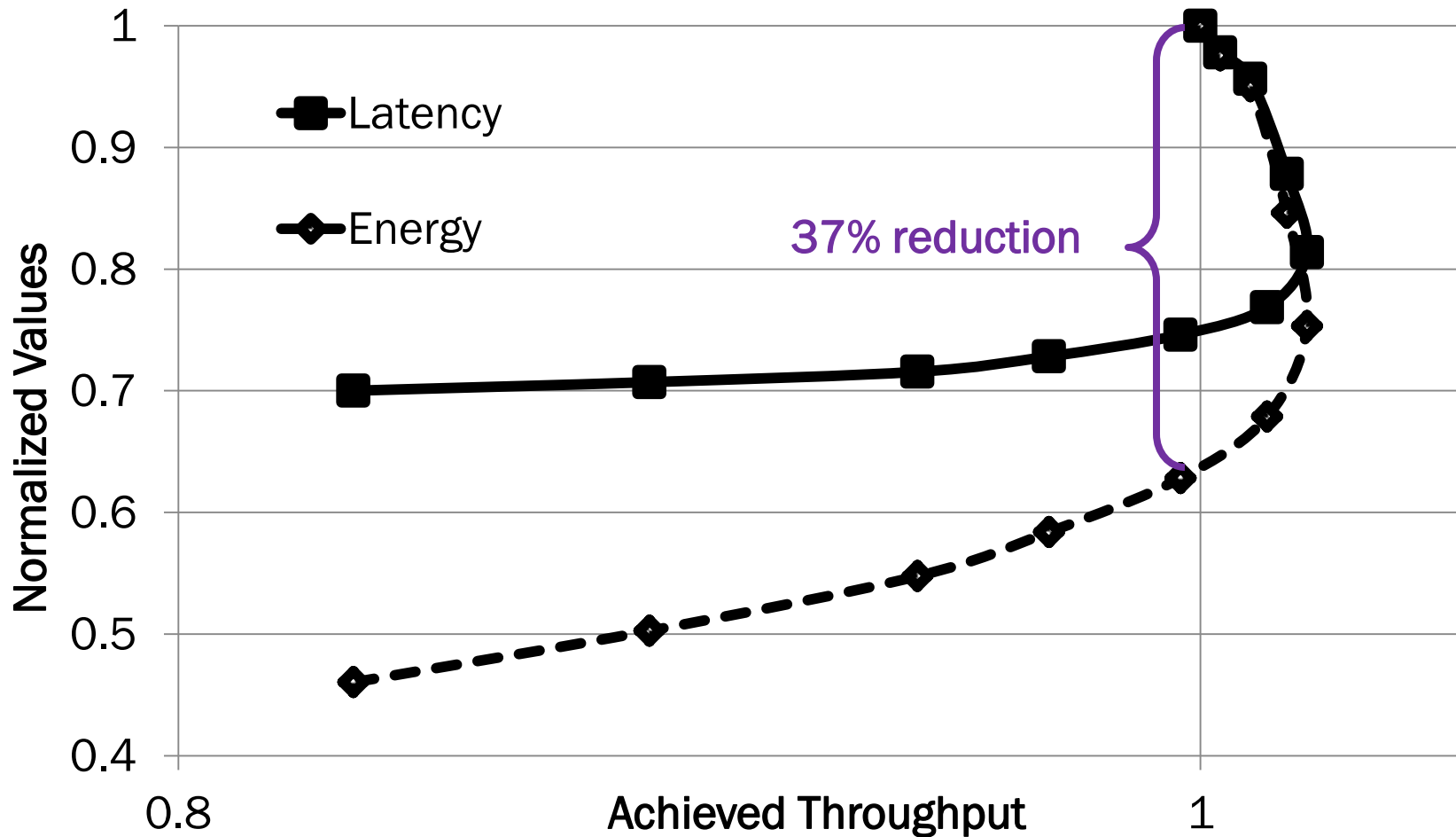
Setups

- ⌘ A 16x16 mesh interconnect.
- ⌘ Cycle-Accurate System-C network model.
- ⌘ Message-Passing.
- ⌘ Communication-Intensive applications with 4 to 35 tasks.
- ⌘ Random sequence of application.
- ⌘ Upon a mapping failure, new *first nodes* are selected until the current application being mapped.

Results- Throughput



Results- Network Metrics



Results- Comparison

Mapping	NMRD	E_{norm}	L_{avg}	%Congestion
CASqA ^{0.5}	1.13	1	41	30.52
CASqA ^{0.5**}	1.13	1.25	48.23	37.7
CoNA	1.69	1.48	50.32	38.4
NN	1.97	1.56	55.53	41.98

** arbitrary task placement within the selected area (without ICEB).

Conclusion

- ✎ Investigation of the middle ground of the spectrum between strictly contiguous and unlimited non-contiguous allocation.
- ✎ The communication distance is yet a matter in today's on-chip networks.
- ✎ CASqA provides the tune to adjust the power/throughput optimal dispersion level, α .
- ✎ The optimum α depends on the application specifications.