

ANALYTICAL PLACEMENT FOR RECTILINEAR BLOCKS

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Outline

1. Background
2. Rectilinear block placement
3. Related works
4. Proposed method
5. Experiments
6. Conclusion remarks

Background

- * LSI production method: much improved
 - * Large number of elements on one chip
- * Design method: not so improved
- * Gap between the production and the design: critical
- * Using IP modules: much promising
 - * Problem: Shape of IP modules may be rectilinear
- * Fast rectilinear placement: important

Rectilinear blocks placement

- * **[Input]** Set of blocks including Rectilinear blocks, Net-list, Chip outline
- * **[Output]** Block placement
- * **[Objective]** Minimization of total wire-length
- * **[Constraint]** No overlap

Framework of previous works

- * Basically, topological relation based
 - * SP, BSG, B*-tree, ...
 - * Utilization of stochastic optimization
 - * SA, ...
-
1. Divide the rectilinear block into a set of rectangles
 2. Add constraints to remain the shape in the perturbation
 3. Enhance the position calculation to recover the shape

Characteristic of topological relation based method

- * **Pros:** No overlap guarantee
 - * from the nature of topological relation
- * **Cons:** Slow convergence
 - * from the nature of stochastic optimization method

Fast placement method: Necessary

Analytical placement is promising

Analytical placement

- * Utilization of the gradient of the objective function
 - * Requirement: differentiable objective function
 - approximation of max and min functions
- * Objective function: non-linear, in general
 - * Requirement: no constraint
 - consideration of relaxed problem

Relaxed problem

- * **[Input]** Set of blocks including Rectilinear blocks, Net-list, Chip outline
- * **[Output]** Block placement
- * **[Objective]** Minimization of total wire-length and less overlap
 - * To obtain less overlap
 - * Density function: widely used
 - * hard to capture overlap among rectilinear blocks
 - * Direct consideration: **Overlap Removable Length**

Approximation of max and min functions

- * Objective function: max and min function included
 - * total wire-length
 - * overlap removable length
- * max and min function: not differentiable
 - * need to approximate them to be differentiable

* Log-Sum-Exponential (LSE) :

$$\text{LSE} = t \log \sum_i e^{\frac{x_i}{t}}$$

* **Pros:** Fast convergence

* **Cons:** Numerical unstableness

t : smoothing parameter

* Stable LSE [Funatsu, 2009]: Numerical stability

Stable-LSE (SLSE)

$$\text{SLSE} = X_{\max} + t \log \sum_i e^{\frac{(x_i - X_{\max})}{t}}$$

where X_{\max} is the maximum number of $\{x_i\}$

- * for all i , $x_i - X_{\max} \leq 0$
- * at least one i , $x_i - X_{\max} = 0$

thus, SLSE is numerical stable

Overlap Removable Length (ORL)

- * Calculate the necessary length to remove the overlap

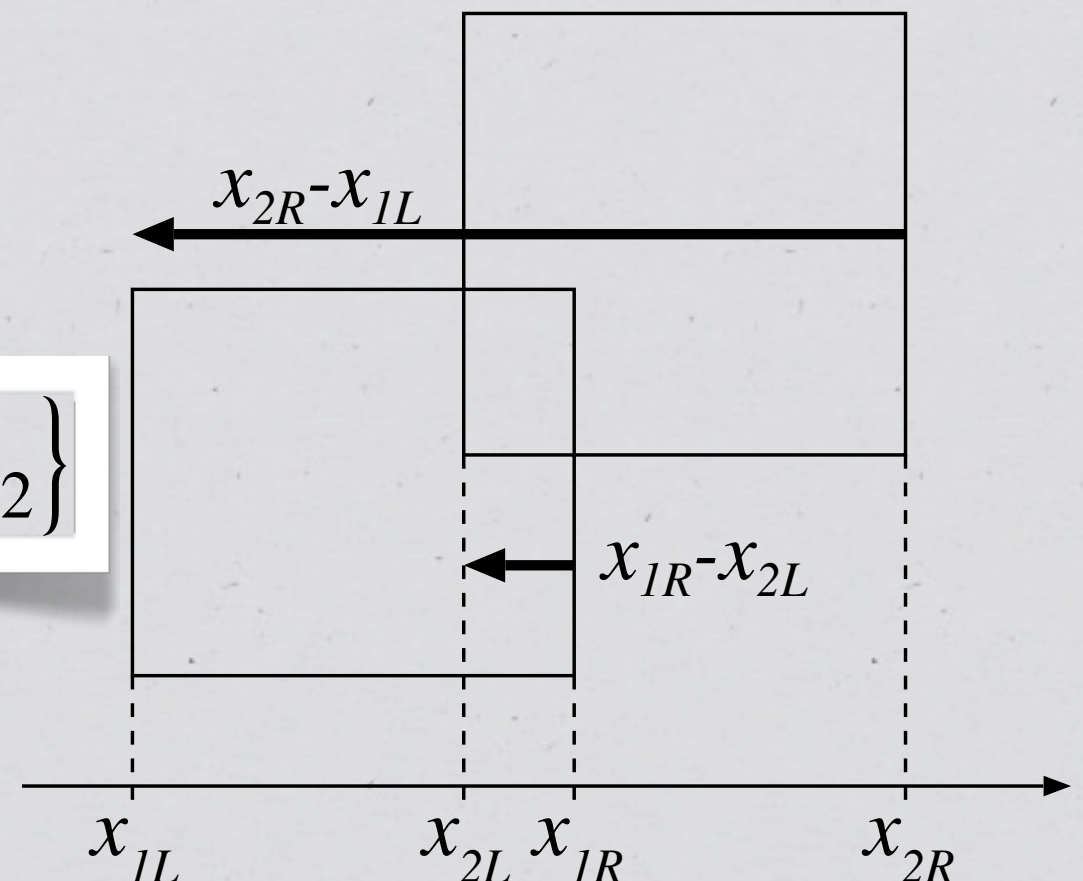
- * for x-coordinate:

$$\text{ORL}_{1,2}^X = \max\{\min\{x_{1R} - x_{2L}, x_{2R} - x_{1L}\}, 0\}$$

- * similar way for y-coordinate

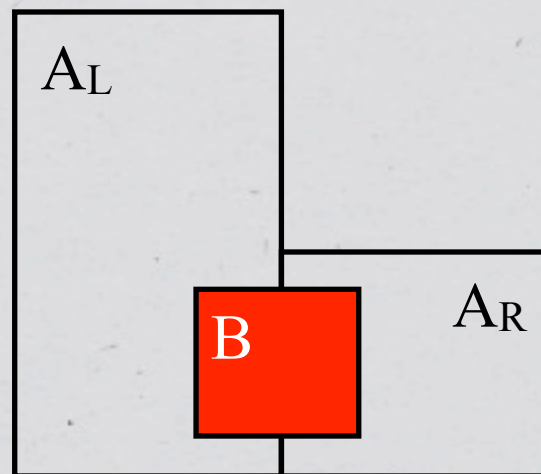
- * ORL for block 1 and 2:

$$\text{ORL}_{1,2} = \min\{\text{ORL}_{1,2}^X, \text{ORL}_{1,2}^Y\}$$



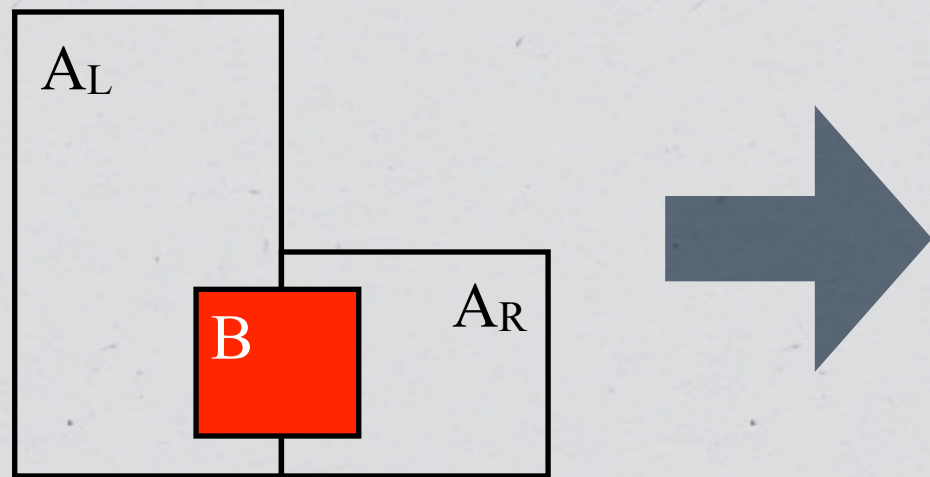
Utilization of ORL: rectangle decomposition

- * Similar to topological relation based method
- * Rectilinear block \rightarrow a set of rectangle blocks
- * seems to be easy to handle them
- * however, there is a main drawback:



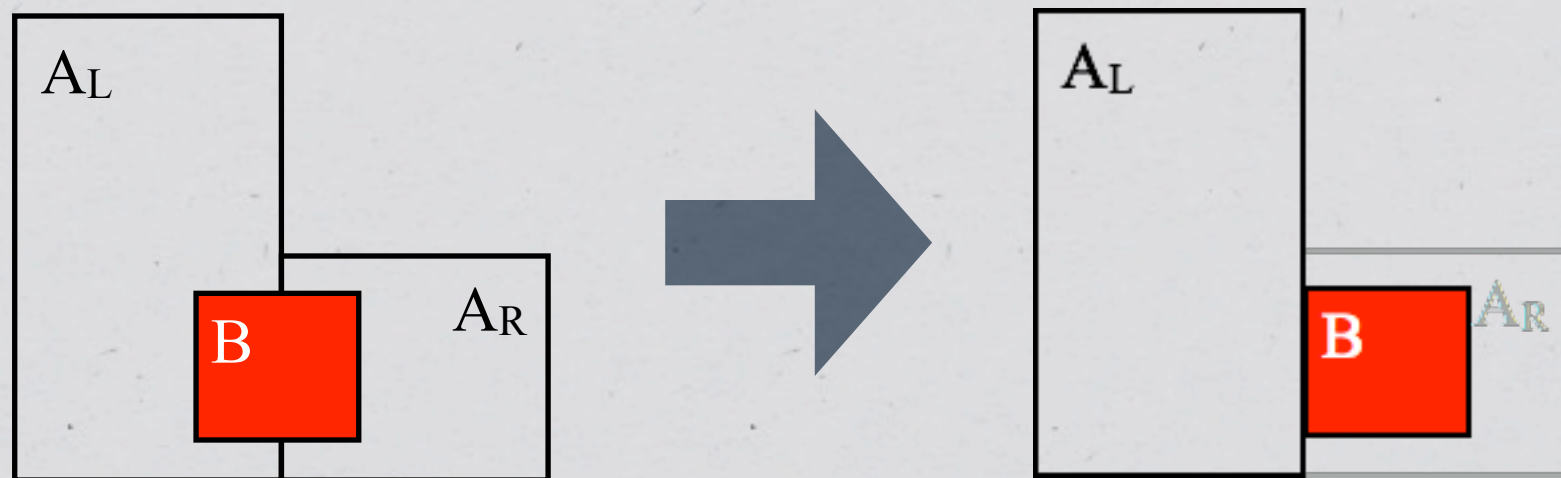
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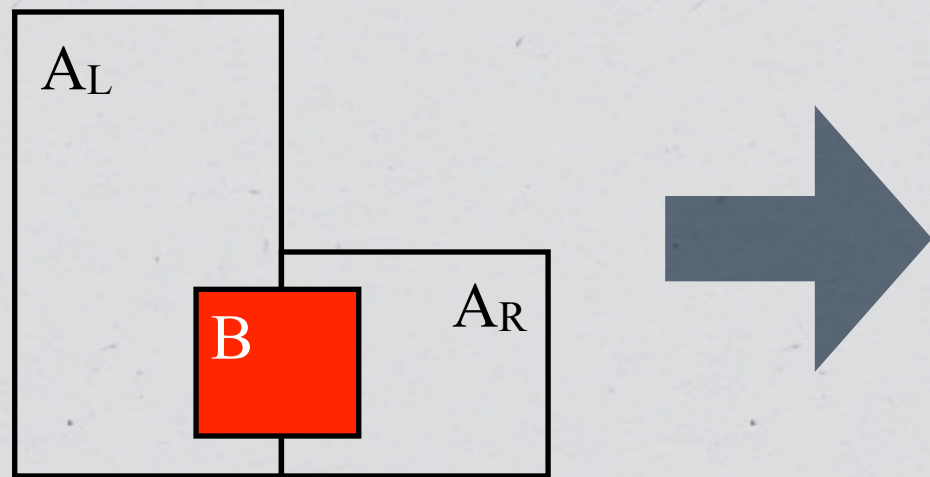
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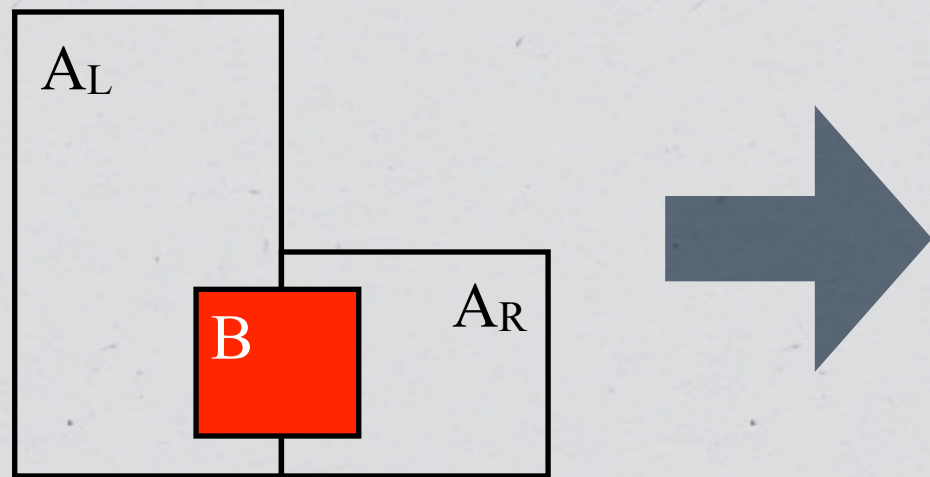
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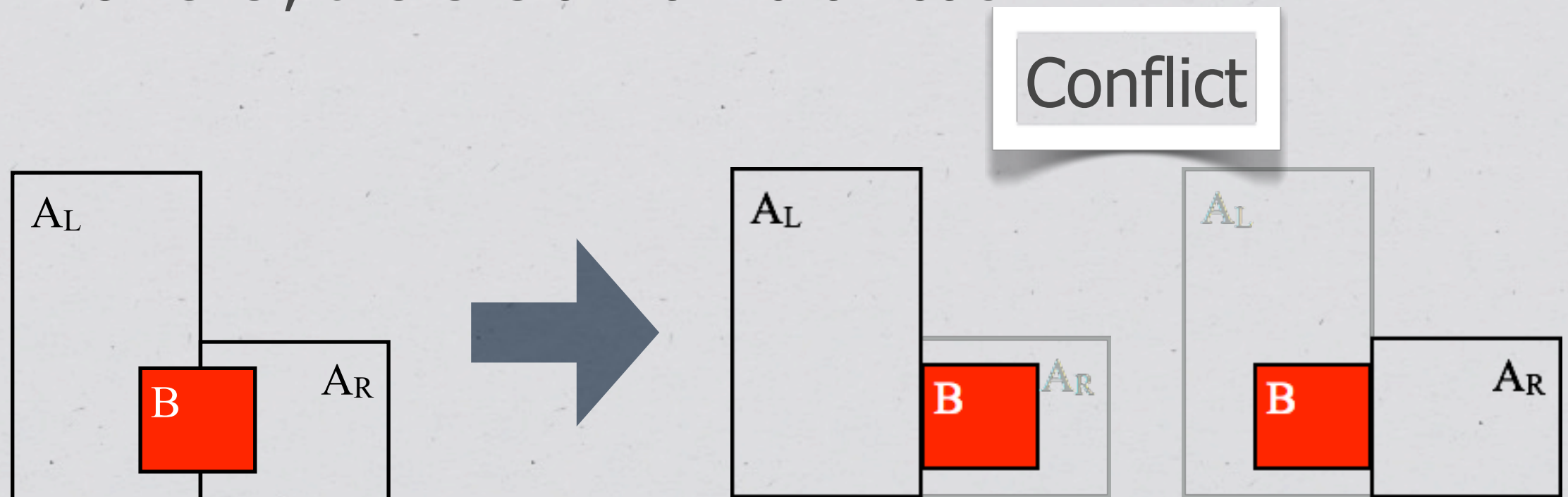
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Analysis of ORL

- * from the definition of ORL:

$$\text{ORL}_{1,2} = \min \{ \text{ORL}_{1,2}^X, \text{ORL}_{1,2}^Y \}$$

- * Equivalent to the minimum length to remove overlap for any directions
 - * for rectangles, consideration of only x- and y-directions is enough
- * For rectilinear blocks, enhance ORL to seek the minimum length for any directions

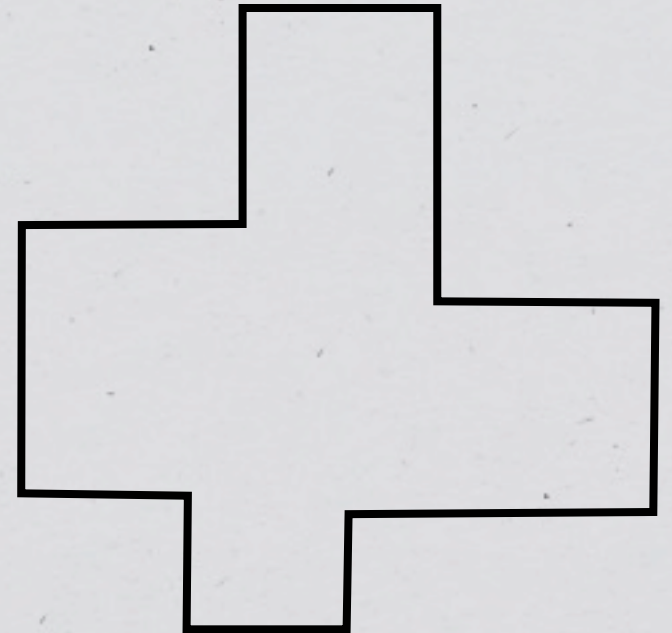
Enhancement of ORL

Pre-process

1. Calculate each side profile of each block
2. For each block pair, enumerate the non-overlap conditions

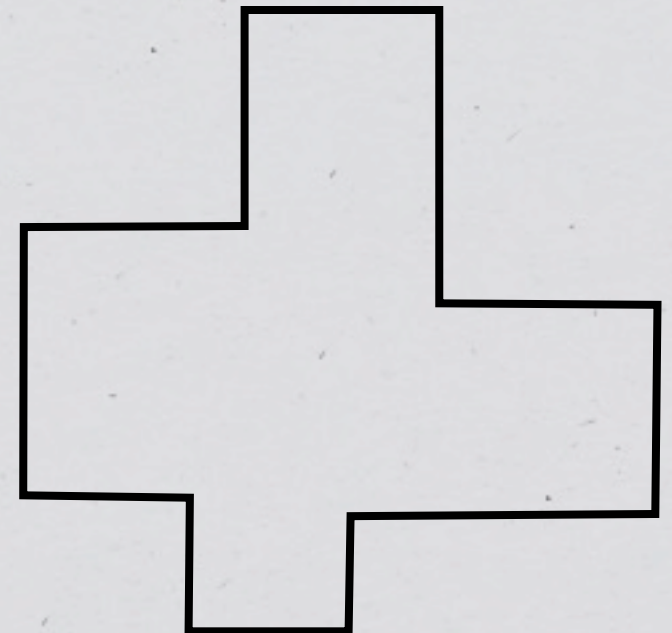
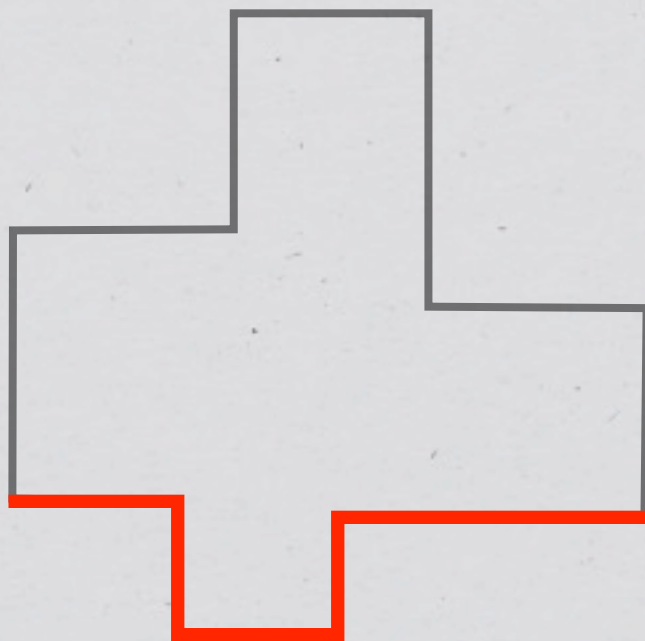
Side profile

* Outer shape of each side



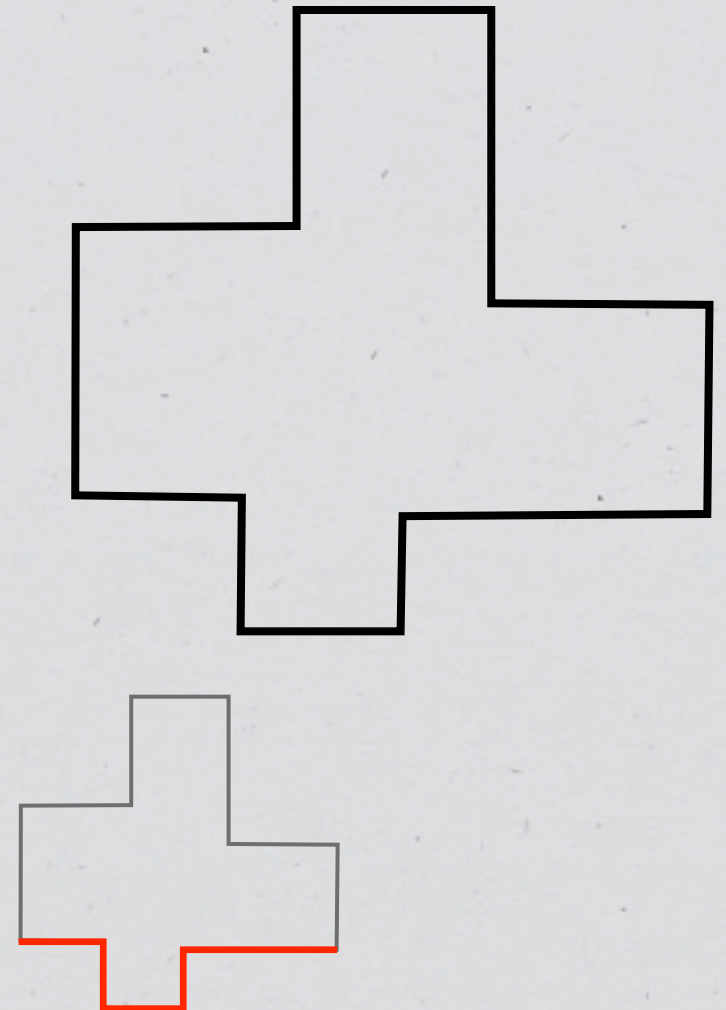
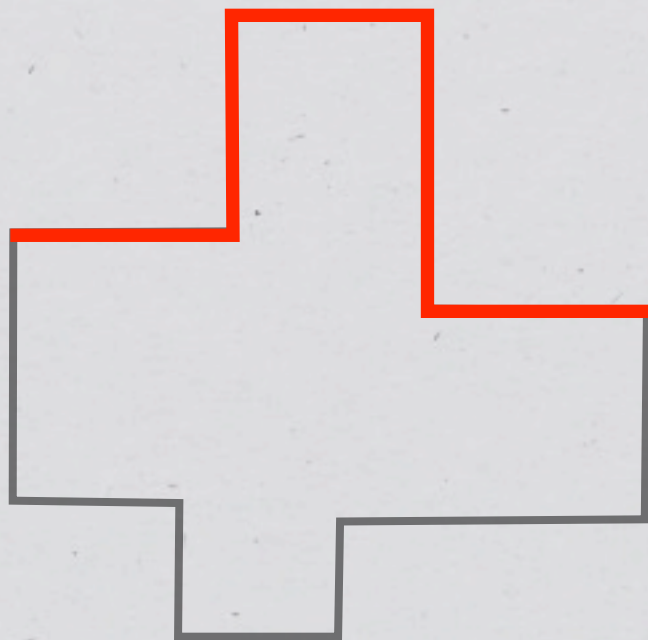
Side profile

- * Outer shape of each side
- * Bottom profile



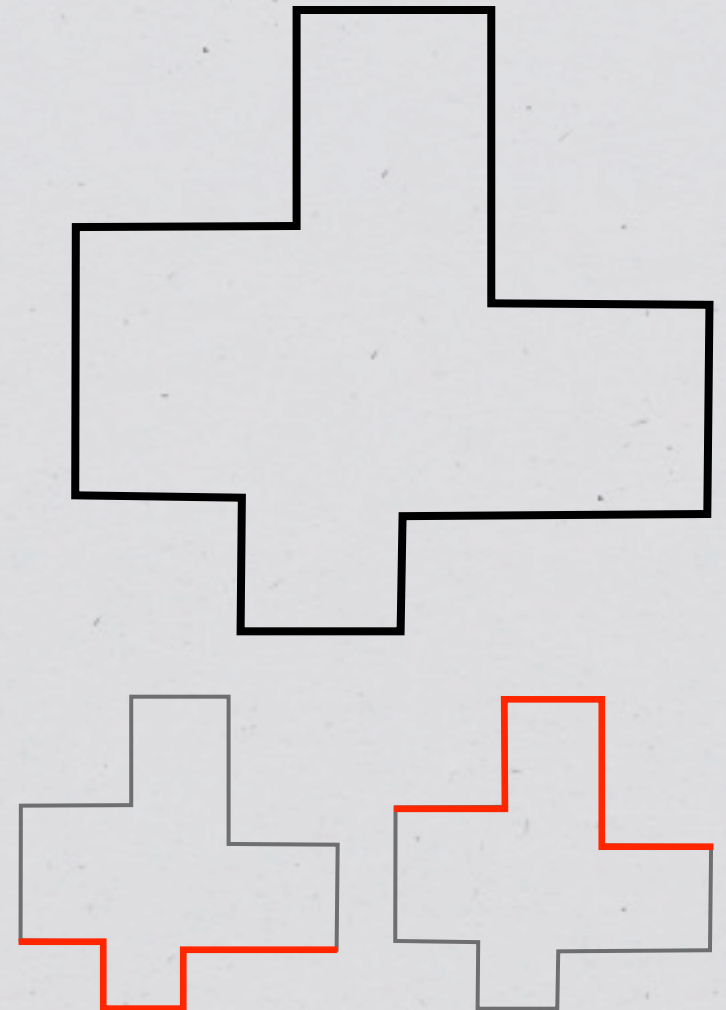
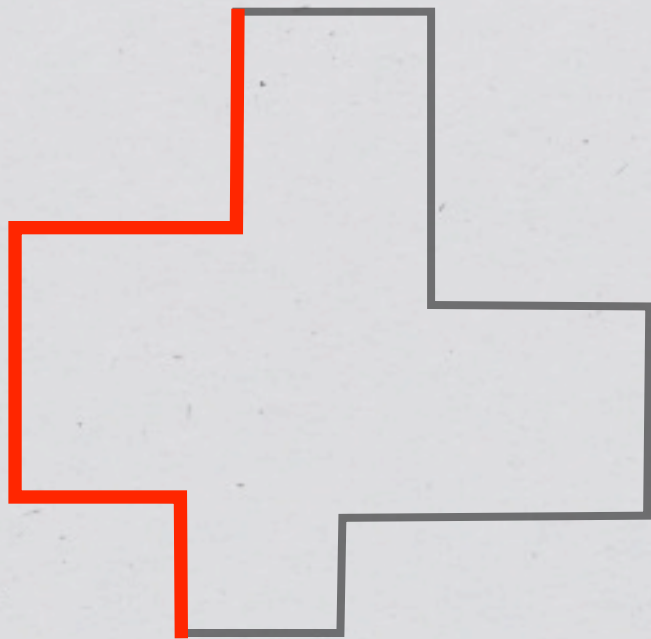
Side profile

- * Outer shape of each side
- * Bottom profile
- * Top profile



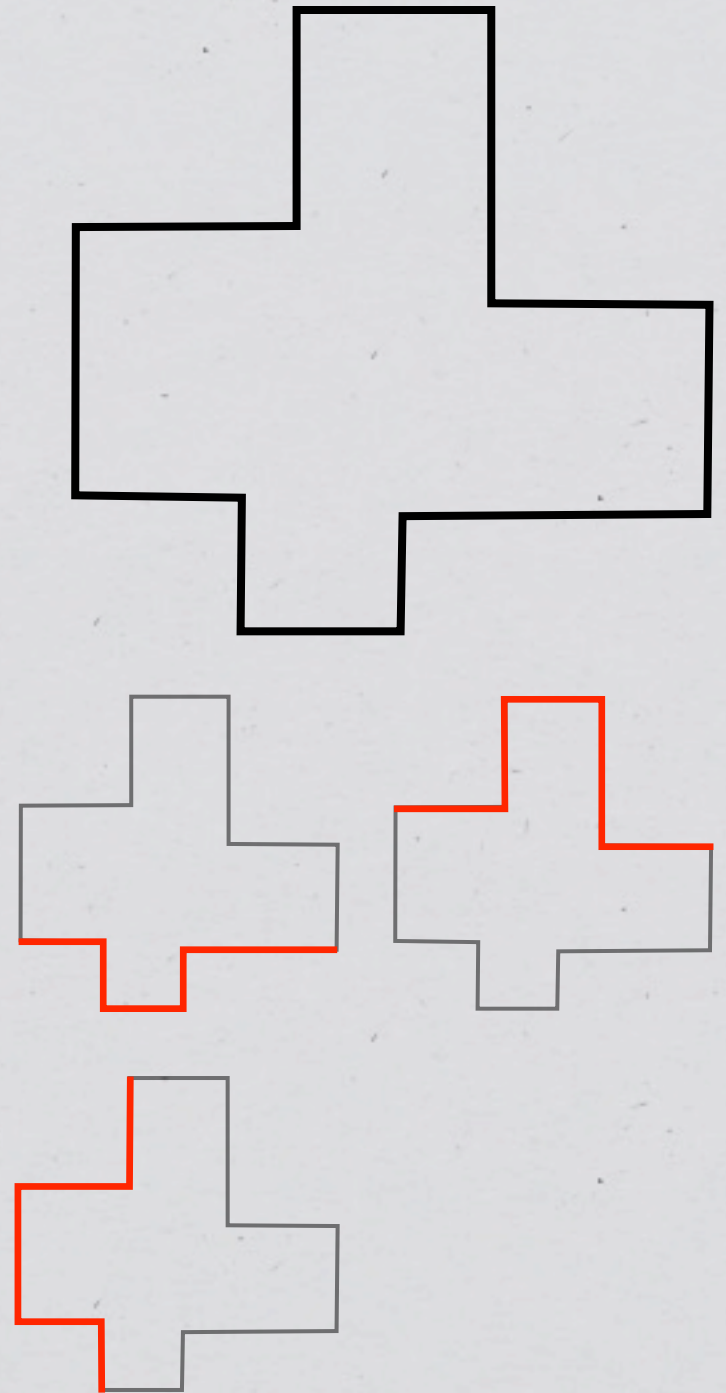
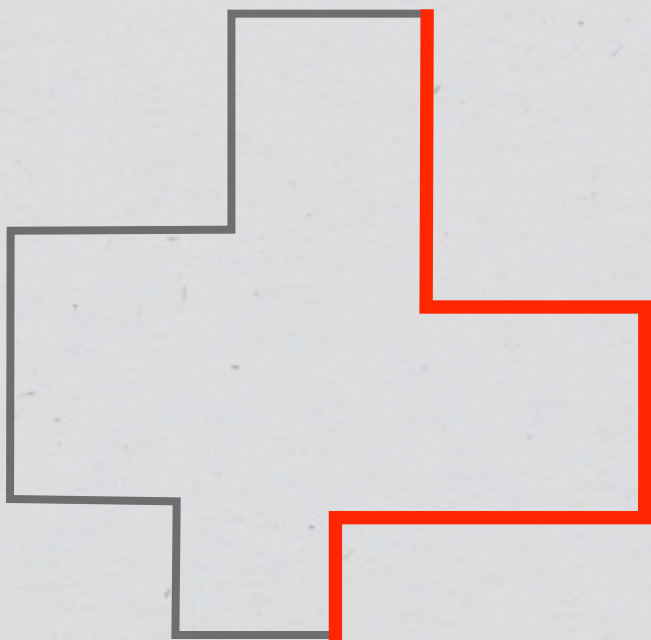
Side profile

- * Outer shape of each side
- * Bottom profile
- * Top profile
- * Left profile



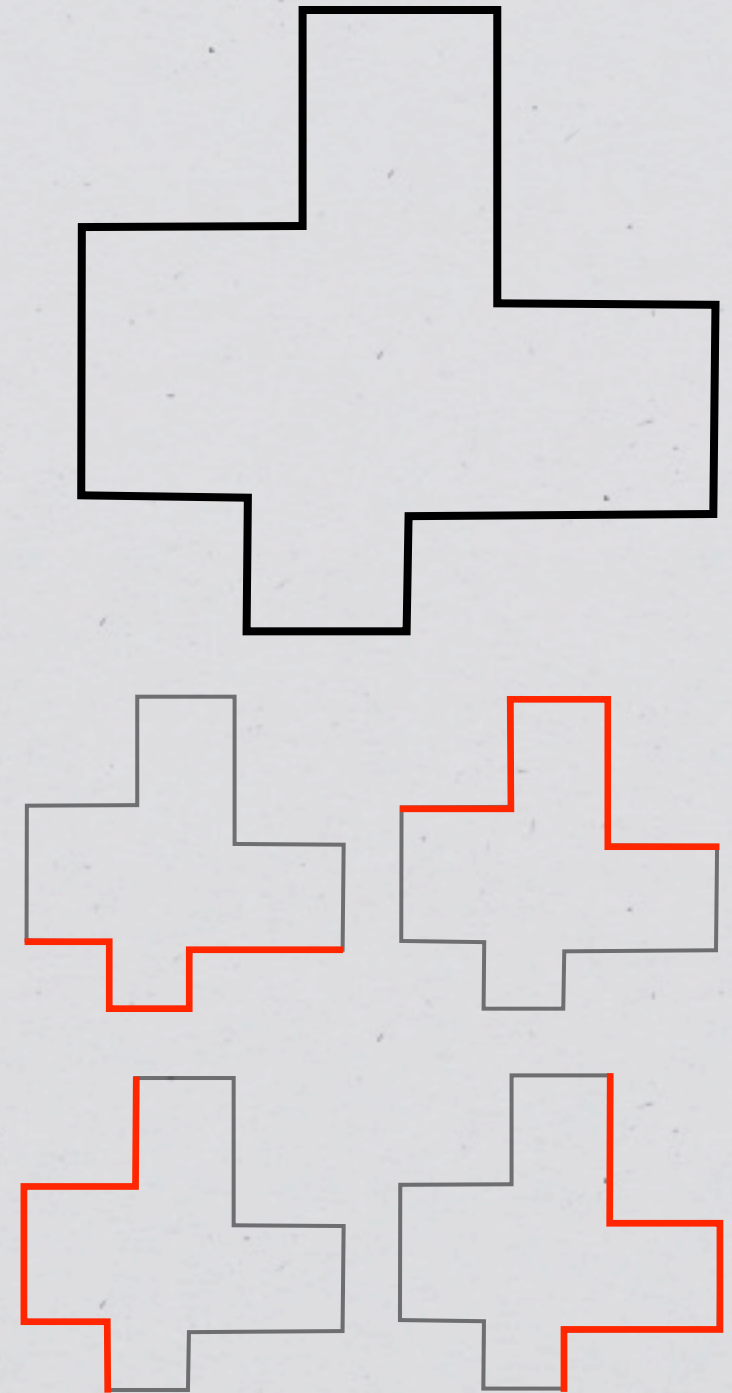
Side profile

- * Outer shape of each side
- * Bottom profile
- * Top profile
- * Left profile
- * Right profile



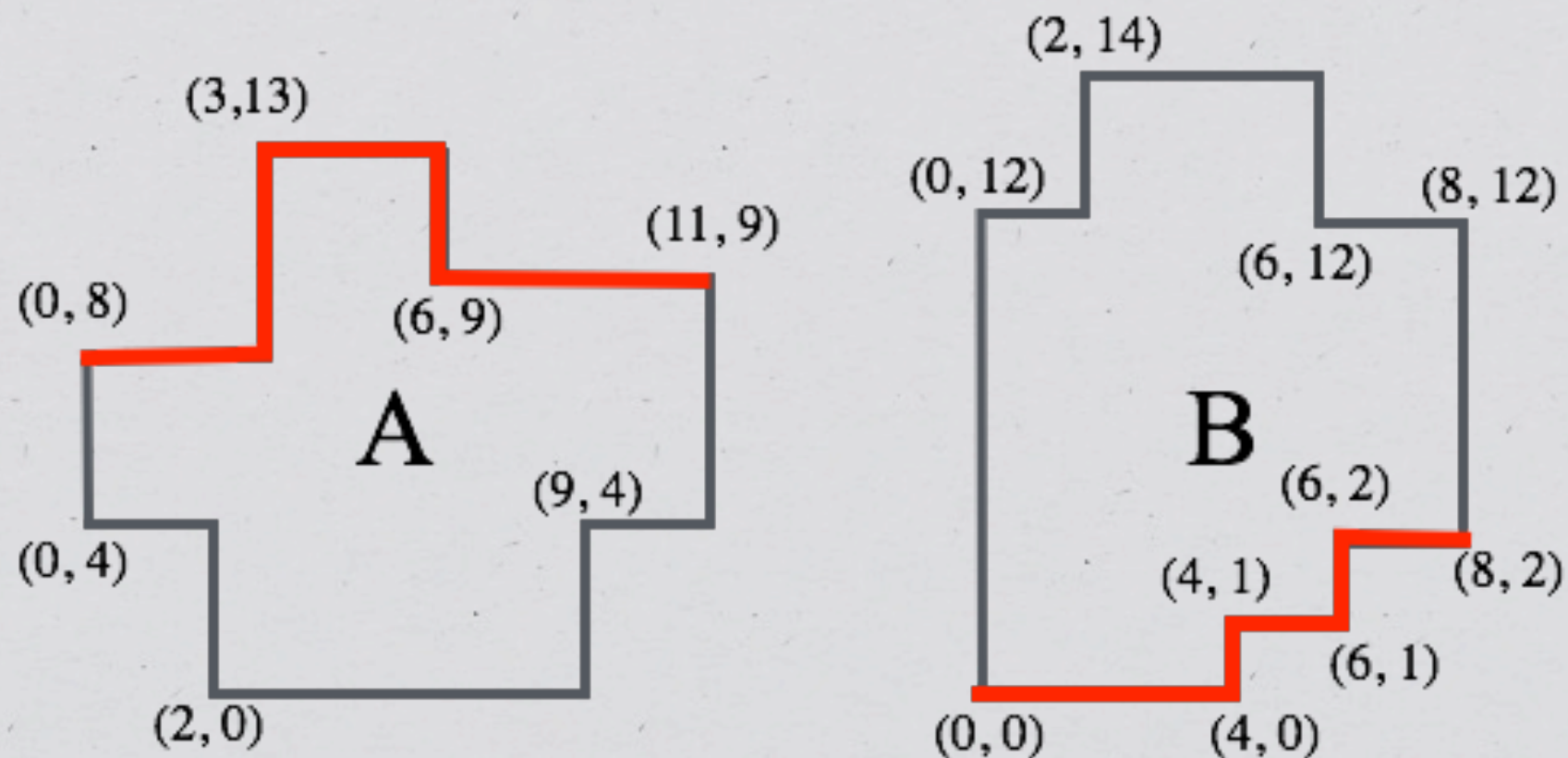
Side profile

- * Outer shape of each side
 - * Bottom profile
 - * Top profile
 - * Left profile
 - * Right profile
- * Complexity: $O(m)$
 - * m : # corners



Enumerate Conditions

- * Non-overlap conditions:
- * represented by relative position



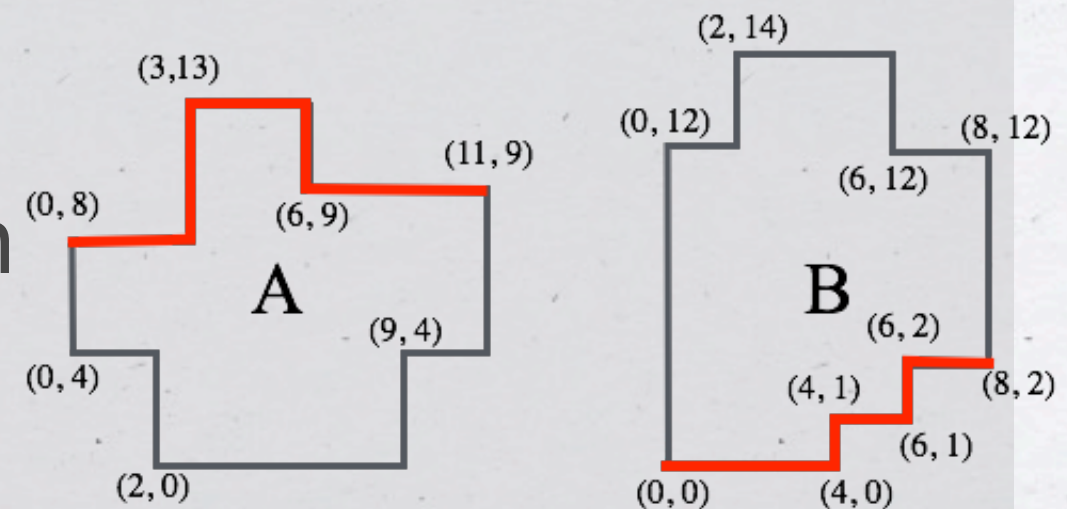
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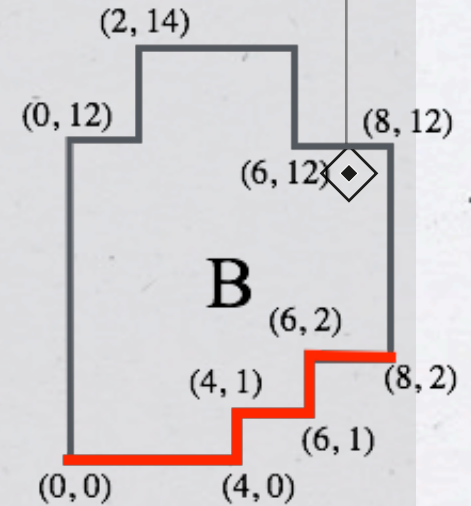
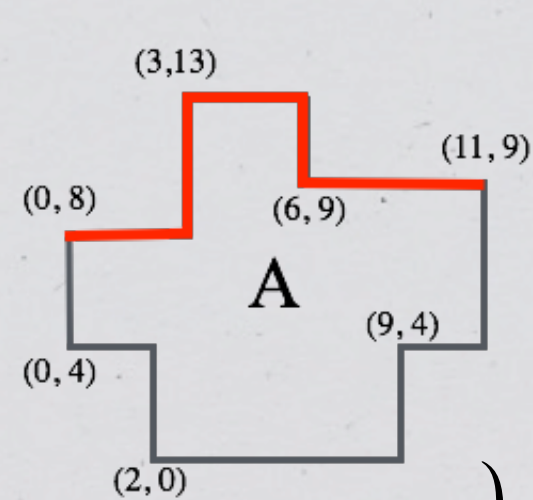
* represented by relative position

$$\begin{aligned}
 &\Delta_x < -8, \\
 &-8 \leq \Delta_x < -6, \Delta_y \geq 6, \\
 &-6 \leq \Delta_x < -5, \Delta_y \geq 7, \\
 &-5 \leq \Delta_x < -3, \Delta_y \geq 11, \\
 &-3 \leq \Delta_x < -1, \Delta_y \geq 12, \\
 &-1 \leq \Delta_x < 6, \Delta_y \geq 13, \\
 &6 \leq \Delta_x < 11, \Delta_y \geq 9, \\
 &11 \leq \Delta_x
 \end{aligned}$$

where $\begin{cases} \Delta_x = x_B - x_A, \text{ and} \\ \Delta_y = y_B - y_A \end{cases}$



ORL



$$\min \left\{ \begin{array}{l} \max \{ \Delta_x + 8, 0 \}, \\ \max \{ \min \{ -\Delta_x - 8, \Delta_x + 6 \}, 0 \} + \max \{ -\Delta_y + 6, 0 \}, \\ \max \{ \min \{ -\Delta_x - 6, \Delta_x + 5 \}, 0 \} + \max \{ -\Delta_y + 7, 0 \}, \\ \max \{ \min \{ -\Delta_x - 5, \Delta_x + 3 \}, 0 \} + \max \{ -\Delta_y + 11, 0 \}, \\ \max \{ \min \{ -\Delta_x - 3, \Delta_x + 1 \}, 0 \} + \max \{ -\Delta_y + 12, 0 \}, \\ \max \{ \min \{ -\Delta_x - 1, \Delta_x - 6 \}, 0 \} + \max \{ -\Delta_y + 13, 0 \}, \\ \max \{ \min \{ -\Delta_x + 6, \Delta_x - 11 \}, 0 \} + \max \{ -\Delta_y + 9, 0 \}, \\ \max \{ \min \{ -\Delta_x - 8, \Delta_x + 6 \}, 0 \} + \max \{ \Delta_y + 8, 0 \}, \\ \max \{ \min \{ -\Delta_x - 6, \Delta_x - 7 \}, 0 \} + \max \{ \Delta_y + 14, 0 \}, \\ \max \{ \min \{ -\Delta_x + 7, \Delta_x - 11 \}, 0 \} + \max \{ \Delta_y + 8, 0 \}, \\ \max \{ -\Delta_x + 11, 0 \} \end{array} \right\}$$

Optimization framework

* Objective function: $\sum \text{WL} + \alpha \left(\sum \text{ORL}^2 + \sum \text{ORL_wcb} \right)$

where

- * WL: wire length with HPWL
- * ORL: Overlap removable length
- * ORL_wcb: Overlap removable length with chip boundary

Optimization flow

1. Construct an initial placement randomly.
2. Set the smoothing parameter $t = 10$.
 - 2.1. Optimize the placement with $\alpha = 0$.
 - 2.2. Optimize the placement with larger α , iteratively.
3. Set the smoothing parameter $t = 0.01$.
 - 3.1. Optimize the placement with larger α , iteratively.

Experiments

- * CPU: Intel Core i5-4570, 3.2GHz
- * Memory: 4GB
- * OS: LinuxMint 17 (qiana)
- * gcc: version 4.8.2
- * Non-linear programming solver: liblbfgs 1.10

- * Benchmarks

name	#blocks	#rectilinear blocks	#nets
n100	100	10	885
n200	200	156	1585
n300	300	243	1893

Experimental results

* Average of 100 trials

* WL Comparison: Rectangle placement with B*-tree

* n100: 32.06, n200: 58.33, n300 71.00

name	ORL	WL	Runtime [sec]
n100	0.204	32.44	5.98
n200	0.248	64.05	56.95
n300	0.770	73.25	132.41

Experimental results

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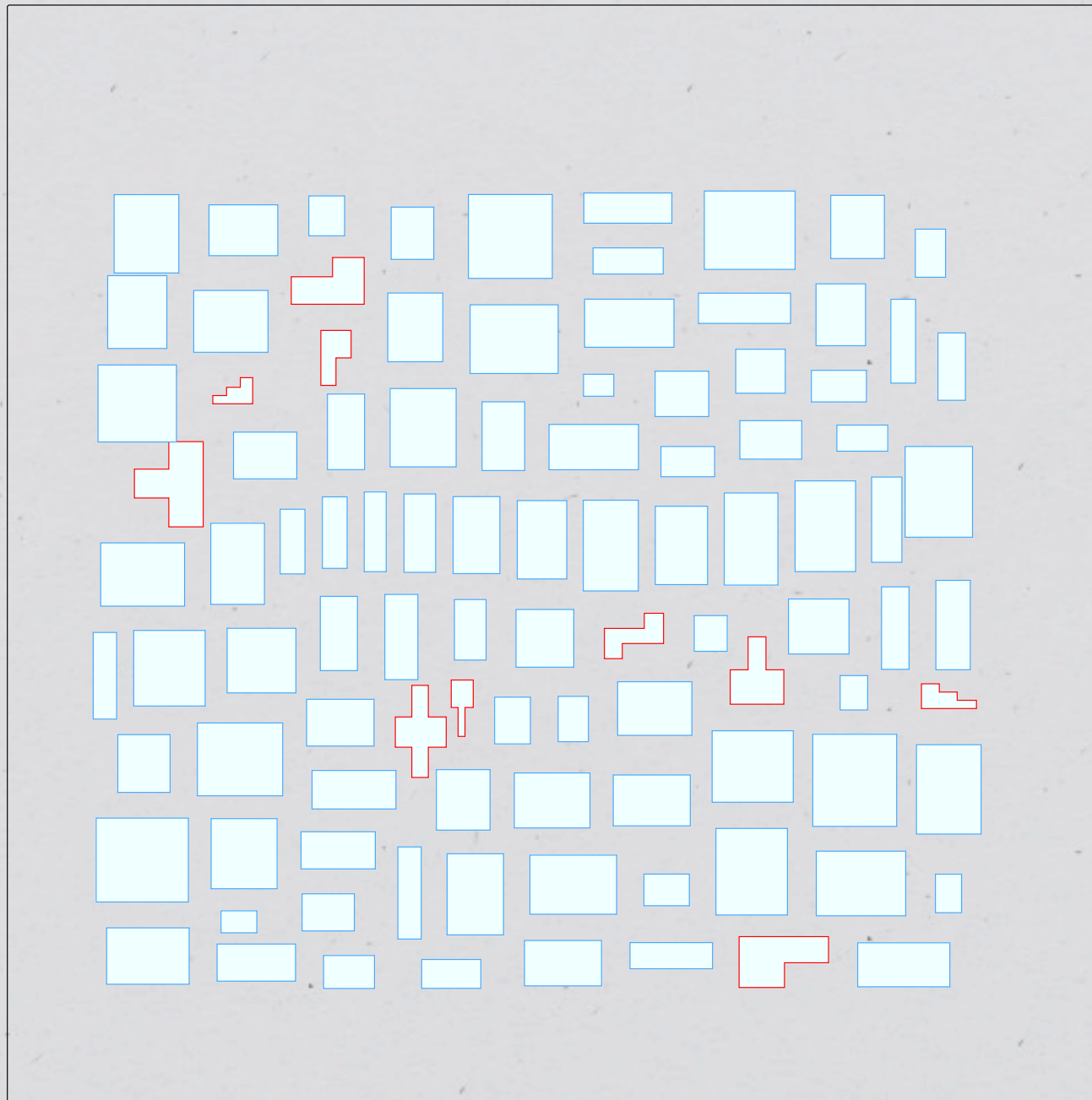
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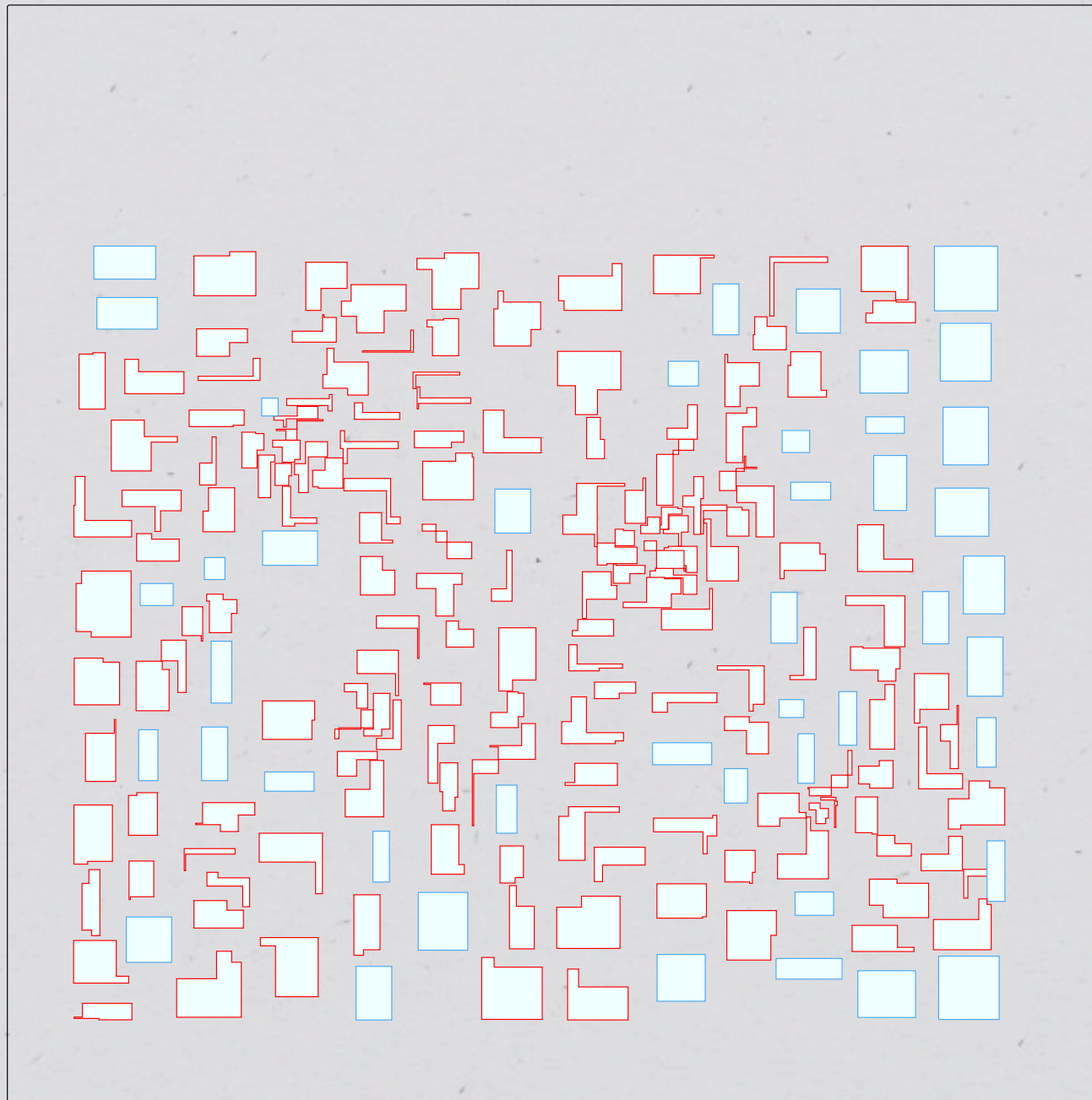
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Efficient

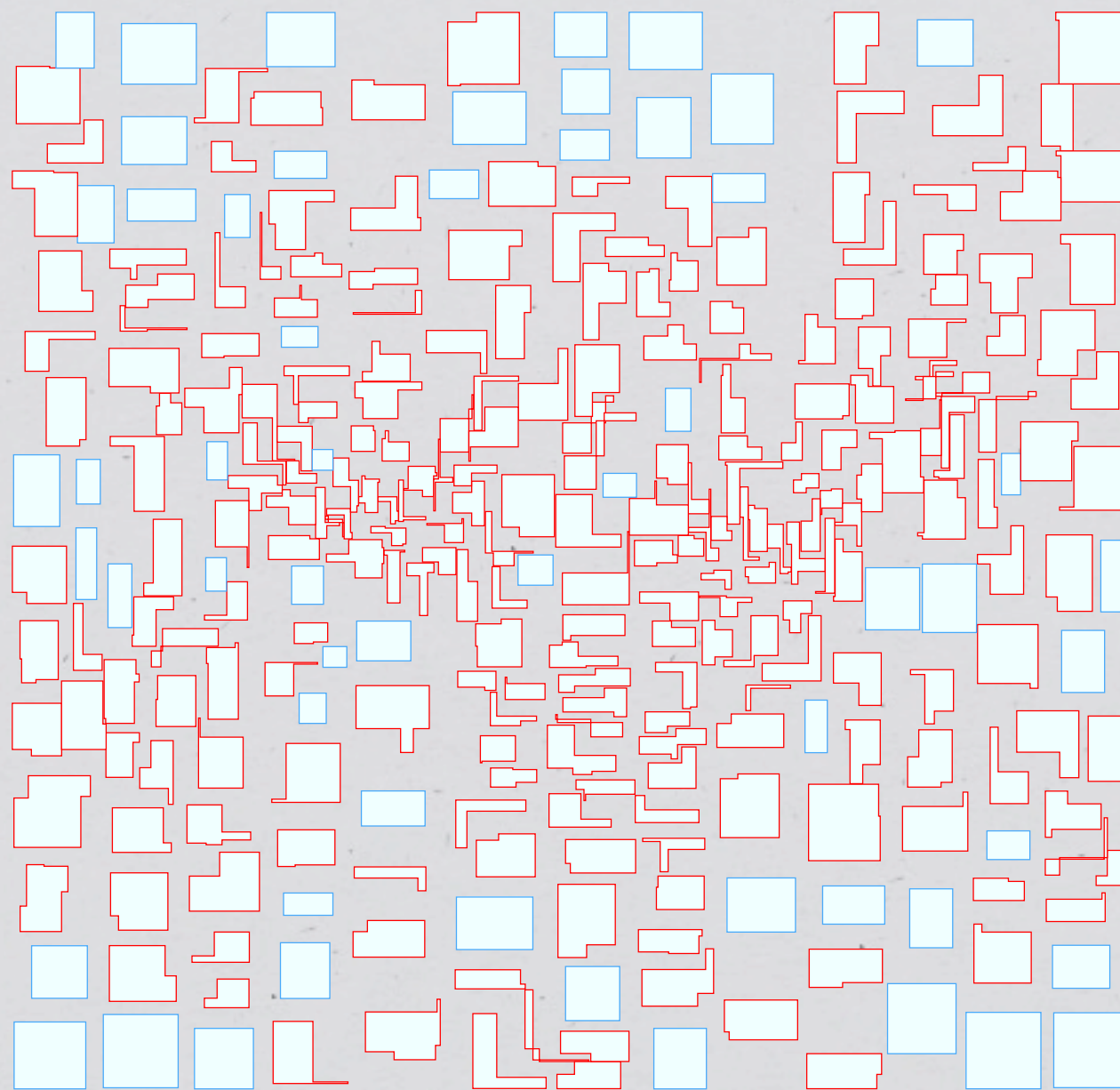
n100



n200



n300



Conclusions

- * Proposition of the analytical placement for the rectilinear blocks.
- * to remove overlap: ORL
- * max and min approximation: Stable-LSE
- * Confirmation of the proposed method, empirically

Future works

- * Refinement of the speed
- * Consideration of the routability
- * Application to other problems
 - * the proposed method: not limited rectilinear blocks
 - * applicable to the problem representing the non-overlap conditions with the differentiable functions