## An Effective Legalization Algorithm for Mixed-Cell-Height Standard Cells

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## Outline

## Introduction

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Proposed Method

Experimental Results


## Conclusions

## Multiple-Row-Height Cells

- Multiple-row-height cells [Beak et al., SPIE'08]
- Minimize cell area
- Minimize intra-cell routing space

Cell area : 54 grids
Cell area : 48 grids
VSS


Single-row-height \& double-row-height cells with an identical function

## Modern Placement Flow



## Legalization

- Given
- A global placement result with cell location ( $x_{i}^{\prime}, y_{i}^{\prime}$ )
- Objective
- Minimize

$$
\sum_{i=1}^{n}\left(\left|x_{i}-x_{i}^{\prime}\right|+\left|y_{i}-y_{i}^{\prime}\right|\right)
$$

- No cell overlaps \& cells are in the chip
- Cells are aligned in rows



## Power-rail Alignment

- Power-rail alignment issues
- An odd-row-height cell alignment can be achieved by vertical cell flipping
- Two types of even-row-height cells: VDD or VSS runs along its top and bottom boundaries



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## Window-based Legalization

- Cells partially inside the window are non-local cells
- Rows are divided by non-local cells into continuous segments
- Find insertion intervals to insert a triple-rowheight cell [Chow et al., DAC'16]



## Single-Row-Height Cells Legalization

- Not much work on legalization with multiple-rowheight standard cells is reported in the literature; however, single-row-height standard cells legalization has been studied for a long time
- Tetris Legalization [Hill, Patent 2002]
- Abacus Legalization [Spindler et al., ISPD'08]


## Modern Legalizers: Tetris vs. Abacus



Global placement

Tetris


Greedy heuristic
$\Rightarrow$ Larger displacement

Small displacement
Abacus w/o violation


Dynamic programming
$\Rightarrow$ Smaller displacement

## Abacus Legalization

- Legalize a cell at a time from left to right in a row
- Minimize quadratic displacement (QD)

Displacement weight for module $i \quad$ Abutment constraint in one row



## Abacus Legalization

- Legalize a cell at a time from left to right in a row

$\mathrm{QD}: \min \left(x_{1}-x_{1}^{\prime}\right)^{2}+\left(x_{2}-x_{2}^{\prime}\right)^{2} \| \underset{6}{6}$ QD: $\min 2\left(x_{1}-\frac{x_{1}^{\prime}+\left(x_{2}^{\prime}-w_{1}\right)}{2}\right)^{2}$


Unknown legal $x$-pos

## Inequality to equality

$$
\Rightarrow x_{1}=\frac{3+(6-5)}{2}=2, x_{2}=7
$$

## Abacus Extension

- Single-row Abacus ignores other rows

- If we consider the maximal cell height as a new row, we extend the dynamic programming method directly. It causes large number of dead space



## Insufficiency of Previous Works

- Window-based legalization insufficiency
- Decide the order of cells arbitrarily
- Find consecutive vacant rows difficultly
- Incur large displacement when moving to other window
- Abacus insufficiency
- Inhibit legalization in one row when cells legalized in another row
- Incur vertical overlap chain


## Previous Works Comparisons



## Contributions

- Propose a multiple-row-height cells legalization based on Abacus
- Remedy Abacus' insufficiencies \& extend its advantages
- Introduce a dead-space-aware cost function for multiple-row-height cells legalization
- Achieve a smaller displacement compared to a leading academic legalizer
- About 50\% smaller $\triangle$ HPWL than the state-of-the-art work


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## Problem Formulation

- Given
- A global placement result with multiple-row-height cells in location ( $x_{i}^{\prime}, y_{i}^{\prime}$ )
- Objective
- Minimize $\sum_{i=1}^{n}\left(\left|x_{i}-x_{i}^{\prime}\right|+\left|y_{i}-y_{i}^{\prime}\right|\right)+D \times \alpha$ $\left(x_{i}, y_{i}\right)$ : legal location of cell $i$ $D$ : total dead space
$\alpha$ : user-defined parameter
- No cell overlaps \& cells are in the chip
- Cells are aligned in rows
- Cells orient to power rails with VDD/VSS constraints


## Multi-Height Cells Legalization



## Cell and Row Selection

 ckelbisstyoung to be legalized

- Friycetherychosestincerinolf doereespareding cost



## PlaceRow Method

- Solve QD by dynamic programming approach:
- solve sub problems optimally to obtain the final solution


Cell 2 overlaps with previous cell


Cluster cell 1 and cell 2 and move the cluster to new global x-pos

## Minimum Cost Selection

- Legalize the overlap cells, and get the cost
- Select the minimum cost row

|  | Order |  |
| :---: | :---: | :---: |
|  | 6 | Cost : 11 |
|  | 4 | Cost: 7 |
|  | 2 | Cost: 3 |
|  | The closest row 1 | Cost : 1 Minimum |
|  | 3 | Cost: 5.2 |
| O | 5 | Cost: 9 |
|  | 7 | Cost : 13 |

## Execution Time Reduction

- Speed up this process by only placing cells into their neighboring rows



## Dead Space Consideration

- Add the area of dead space into the cost function

$$
-\min \sum_{i=1}^{n}\left(\left|x_{i}-x_{i}^{\prime}\right|+\left|y_{i}-y_{i}^{\prime}\right|\right)+\alpha \times D
$$

$\alpha$ is user-defined parameter
$D$ is total dead space



Original cost function


Dead space-aware cost function

## Multiple Overlaps Solution

- Choose the row with maximum overlapping area when overlapping with previous cells in different rows
- Do Multi-PlaceRow on that row


Cell 5 overlaps with cell 3 and cell 2


Cell 5 clusters with cell 3


Clustered cells then cluster with cell 2

## Multi-PlaceRow Analysis

- If every multiple-row-height cell is never clustered and only the horizontal movement is considered, the solution generated by MultiPlaceRow is optimal



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## Experimental Settings

- Platform
- C++ programming language
- 64-bit Linux machine
- Intel Xeon 2.93 GHz CPU with 48GB memory
- Comparison
- Comparison with Chow et al., DAC'16
- Benchmarks
- Same designs adapted by Chow et al.


## Experimental Results

## - $52 \%$ smaller $\triangle H P W L$ than the ILP method

- $59 \%$ smaller $\Delta H P W L$ than the state-of-the-art work

| Benchmark | \#S. Cell | \#D. Cell | Density | GP HPWL <br> (m) | Disp. (sites) |  |  | $\triangle \mathrm{HPWL}$ |  |  | Runtime (s) |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | ILP | DAC'16 | Ours | ILP | DAC'16 | Ours | ILP | DAC'16 | Ours |
| des_perf_1 | 103842 | 8802 | 0.91 | 1.43 | 2.13 | 3.32 | 3.46 | 2.61\% | 2.85\% | 0.96\% | 4098.7 | 7.0 | 18.0 |
| des_perf_a | 99775 | 8513 | 0.43 | 2.57 | 0.66 | 0.96 | 0.68 | 0.11\% | 0.28\% | 0.14\% | 193.8 | 2.6 | 6.1 |
| des_perf_b | 103842 | 8802 | 0.50 | 2.13 | 0.62 | 0.85 | 0.64 | 0.12\% | 0.31\% | 0.16\% | 250.8 | 2.4 | 6.2 |
| fft_1 | 121913 | 5500 | 0.46 | 5.25 | 0.45 | 0.47 | 0.47 | 0.09\% | 0.10\% | 0.10\% | 206.0 | 1.9 | 7.8 |
| fft_2 | 30297 | 1984 | 0.84 | 0.46 | 1.58 | 1.81 | 1.55 | 2.25\% | 1.66\% | 0.93\% | 776.8 | 1.1 | 1.3 |
| fft_a | 30297 | 1984 | 0.50 | 0.46 | 0.66 | 0.86 | 0.64 | 0.55\% | 0.87\% | 0.68\% | 72.7 | 0.4 | 0.8 |
| fft_b | 28718 | 1907 | 0.25 | 0.75 | 0.60 | 0.64 | 0.64 | 0.32\% | 0.33\% | 0.33\% | 38.2 | 0.3 | 0.8 |
| matrix_mult_1 | 28718 | 1907 | 0.28 | 0.95 | 0.73 | 0.80 | 0.62 | 0.32\% | 0.33\% | 0.27\% | 61.9 | 0.4 | 1.0 |
| matrix_mult_2 | 152427 | 2898 | 0.80 | 2.39 | 0.49 | 0.53 | 0.48 | 0.36\% | 0.28\% | 0.22\% | 967.4 | 3.9 | 9.1 |
| matrix_mult_a | 152427 | 2898 | 0.79 | 2.59 | 0.45 | 0.49 | 0.44 | 0.30\% | 0.22\% | 0.17\% | 825.0 | 4.0 | 8.9 |
| matrix_mult_b | 146837 | 2813 | 0.42 | 3.77 | 0.27 | 0.33 | 0.27 | 0.09\% | 0.14\% | 0.09\% | 150.7 | 1.3 | 9.3 |
| matrix_mult_c | 143695 | 2740 | 0.31 | 3.43 | 0.25 | 0.30 | 0.25 | 0.09\% | 0.13\% | 0.09\% | 127.8 | 1.3 | 8.9 |
| pci_bridge32_a | 143695 | 2740 | 0.31 | 3.29 | 0.27 | 0.29 | 0.27 | 0.11\% | 0.11\% | 0.11\% | 139.0 | 1.4 | 9.0 |
| pci_bridge32_b | 26268 | 3249 | 0.38 | 0.46 | 0.88 | 0.95 | 0.88 | 0.52\% | 0.58\% | 0.63\% | 49.4 | 0.3 | 0.8 |
| superblue12 | 25734 | 3180 | 0.14 | 0.98 | 0.95 | 0.96 | 0.52 | 0.12\% | 0.13\% | 0.12\% | 15.3 | 0.2 | 0.7 |
| superblue11_a | 861314 | 64302 | 0.43 | 42.94 | 1.85 | 1.94 | 1.86 | 0.15\% | 0.15\% | 0.14\% | 3073.6 | 23.4 | 80.2 |
| superblue12 | 1172586 | 114362 | 0.45 | 39.23 | 1.45 | 1.63 | 1.63 | 0.18\% | 0.22\% | 0.25\% | 5079.0 | 106.5 | 91.1 |
| superblue14 | 564769 | 47474 | 0.56 | 27.98 | 2.56 | 2.62 | 2.38 | 0.22\% | 0.22\% | 0.16\% | 3360.6 | 17.1 | 70.3 |
| superblue16_a | 625419 | 55031 | 0.48 | 31.35 | 1.61 | 1.73 | 1.68 | 0.10\% | 0.12\% | 0.11\% | 2470.7 | 21.7 | 61.0 |
| superblue19 | 478109 | 27988 | 0.52 | 20.76 | 1.52 | 1.60 | 1.68 | 0.14\% | 0.14\% | 0.11\% | 1848.8 | 10.9 | 44.6 |
|  |  |  |  | Average | 1.00 | 1.16 | 1.05 | 0.44\% | 0.46\% | 0.29\% | 1190.3 | 10.4 | 21.8 |
|  |  |  |  | N. <br> Average | 0.95 | 1.11 | 1.00 | 1.52 | 1.59 | 1.00 | 54.6 | 0.48 | 1.0 |

## Layout

- The benchmark circuit fft_2


Legalization result


A partial layout of legalization result

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## Conclusions

- Develop an effective algorithm for the mixed-cell-height standard cells legalization problem
- Derive a dead-space-aware objective function and an optimization scheme to handle this issue
- Achieve best wirelength among all published methods in reasonable running time as shown in experiment results, e.g., about 50\% smaller wirelength increase than the state-of-the-art work which is a best paper nominee at DAC'16


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## Multi-PlaceRow Analysis - Case I

- The to-be-inserted cell 5 keeps its position while there is no overlap



## Multi-PlaceRow Analysis - Case II

- The to-be-inserted cell 5 overlaps with previous cell 4 in row 2



## Multi-PlaceRow Analysis - Case III-A

- The to-be-inserted cell 5 overlaps with previous cell 4 in row 2 and cell 3 in row 1



## Multi-PlaceRow Analysis - Case III-B

- The overlapping area between cell 3 and cell 5 is larger than that between cell 4 and cell 5
- Cluster cell 5, cell 3, and cell 1
- After the moving, cell 5 overlaps cell 4
- Add cell 4 to the previous cluster
- After all cells are non-overlapping, cell 1 exits its cluster to restore to original position



## Future Work: Order of Cells Determination

- Determine the legalizing order of cells before legalization, instead of deciding order only by x-coordinate

Determine by x -coordinate


Determine by preprocess


## Future Work: Vertical Power Rail Awareness

- Vertical power rails are usually implemented by Metal 2 in process, and some designs also use Metal 2 to implement the power rails of multiple-row-height cells
- We should avoid vertical power rails overlapping with the power rails of multiple-row-height cells


