# Crosstalk Analysis using Reconvergence Correlation

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## Outline

- Problem Statement
- Relative Timing Windows
- Crosstalk Analysis using Relative Timing Windows
- Probabilistic Analysis
- Results and Conclusions

#### **Problem Statement**



- Timing window at N1 and N2 overlap
  - Hence considered do be switching together
- But, if N switches at time t,
  - N1 and N2 can only witch at [t+1, t+2] and [t+3, t+5] respectively.
  - N1 and N2 can never switch together
- Pessimism in current methodology
- Current methodology ignores the correlation (Reconvergence) between N1 and N2

#### Timing Windows (Classical)

- Timing windows at Ni TW(Ni) =  $(t^{i}_{min}, t^{i}_{max})$
- Delay from Ni to N D(Ni →N)= $(d^{i}_{min}, d^{i}_{max})$
- Timing Windows at N

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$$\Gamma W(N) = \bigcup_{i=1}^{n} \left( t_{\min}^{i} + d_{\min}^{i}, t_{\max}^{i} + d_{\max}^{i} \right)$$

- Due to addition, relative information is lost,
- We need a way to preserve this relative information

#### **Relative Timing Windows**

- In relative timing windows based approach, delays and timing windows of previous net are preserved
- TW at N relative to Ni

 $\mathsf{TW}(\mathsf{N}/\mathsf{N}i) = \{\mathsf{TW}(\mathsf{N}i), \mathsf{D}(\mathsf{N}i \rightarrow \mathsf{N})\}$ 



#### **Crosstalk Analysis**

• To find if A1 and A2 can switch together,

- Find latest common divergence points in fan-in cone of A1,A2. Let these points are set 𝒫.
- Find relative TW of A1,A2 relative to all points in  $\Phi$

$$TW(A1) = \bigcup_{i \in \phi} \{TW(i), D(i \rightarrow A1)\}$$
$$TW(A2) = \bigcup_{i \in \phi} \{TW(i), D(i \rightarrow A2)\}$$

Common part (TW(i)) can be ignored, hence nets A1 and A2 can switch together if D(i  $\rightarrow$ A1) and D(i  $\rightarrow$ A2) overlap



- To find if N1 and N2 can switch together,
  - Find latest divergence point (N in this case)
  - Find relative TW of N1,N2 relative to N TW(N1/N)={TW(N),(1,2)} TW(N2/N)={TW(N),(3,5)}
  - Common part (TW(N)) can be ignored hence nets N1 and N2 don't switch together

#### **Probabilistic Analysis**

- Let victim has N aggressors
- All switching events are uniformly distributed in time period T.
- If width of timing window of  $i_{th}$  aggressor is  $\tau_i$ .
- Probability that all nets can switch together

$$P_{old} = \left(\prod_{i=1}^{N} \frac{\tau_i}{T}\right) \left(\sum_{j=1}^{N} \frac{T}{\tau_j}\right)$$

**Conventional Approach** 

Let width of TW at divergence point is d, Prob. will be

$$P_{new} = \left(\prod_{i=1}^{N} \frac{\tau_i - d}{T}\right) \left(\sum_{j=1}^{N} \frac{T}{\tau_j - d}\right)$$

**Our Approach** 

#### Results (Crosstalk Glitch)



Part 1 : Nets below threshold (using conventional and proposed approach) = 139634

Part 2 : Violation pruned with proposed approach = 2053

Part 3 : New violations due to proposed approach = 0

Part 4 : Nets above threshold (using existing and proposed approach) = 14906

on a 65nm design

### Results (Crosstalk Delay)



on a 65nm design

## Conclusions

- Current approach does not consider correlation between nets
  - Can lead to pessimistic analysis
- We proposed use of relative timing windows to address this pessimism
- We analytically found the effectiveness of the approach using probabilistic methods
- Our approach pruned many false violation on a real design

## Thank You

Q&A