

# Closed form solution for optimal buffer sizing

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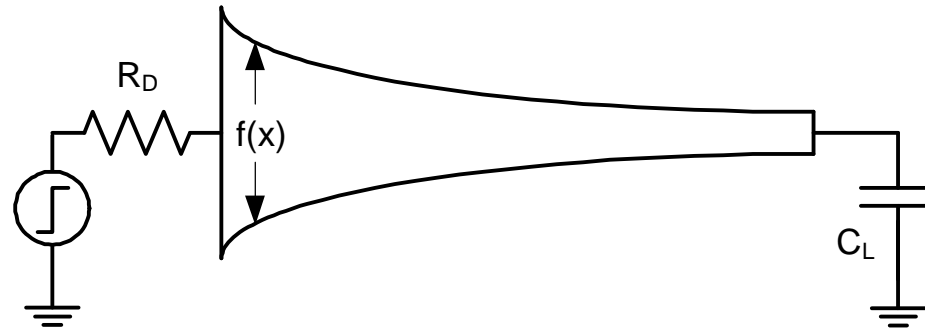
ASPDAC 2006  
January 25, 2006

# Outline

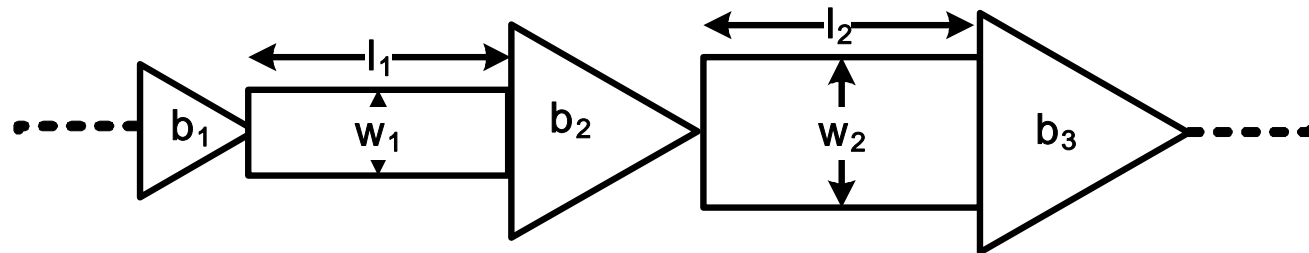
- Introduction
- Assumptions
- Recurrence relation
- Closed form solution
- Testing procedure and results
- Conclusion

# Some delay minimization techniques

- Wire sizing

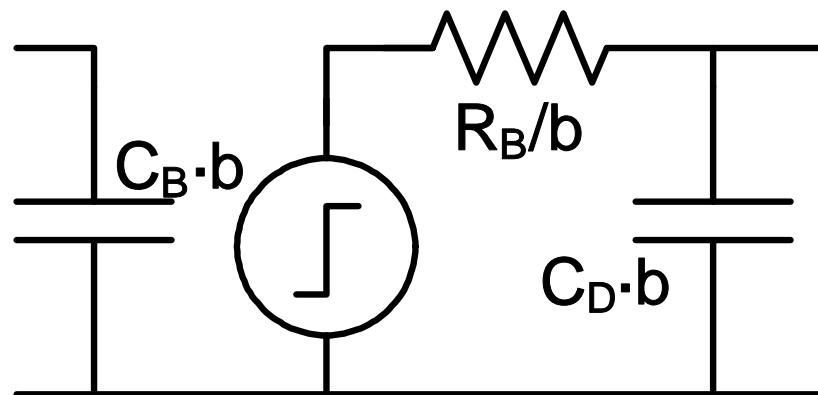


- Buffer insertion



# Buffer sizing

- Can also use buffers of different sizes
- Larger buffer has smaller output resistance, but represents larger capacitive load to previous wire segment
- Use a very simple buffer model (buffer of size  $b$ )

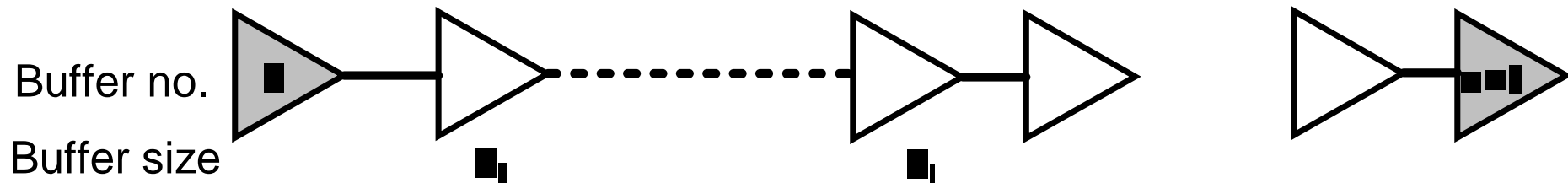


# Model assumptions

- Single interconnection wire with fixed width
- No wire sizing allowed
- Elmore delay model
- Use      model for wire segments
- Allow buffer sizing
- Number of buffers given
- Buffers are equally spaced along the wire
- Driver (resistance) and load (capacitance) are buffers of fixed size

# Interconnect with buffers

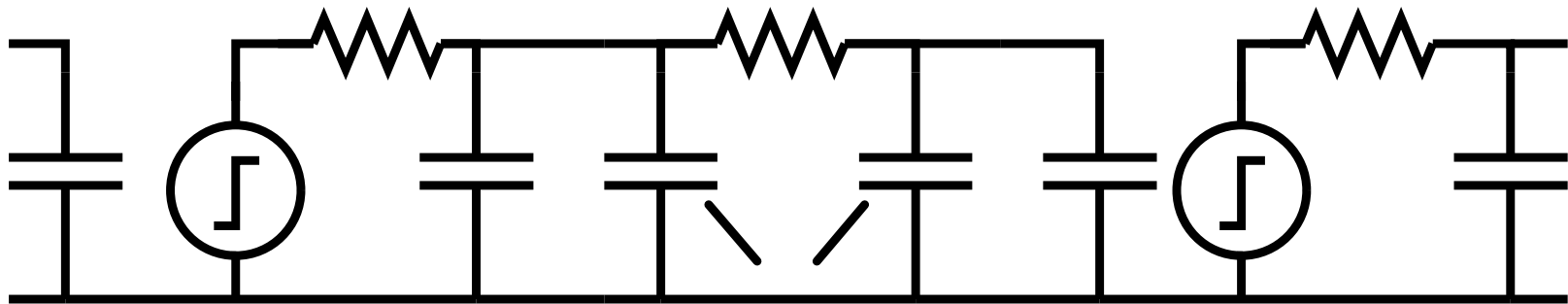
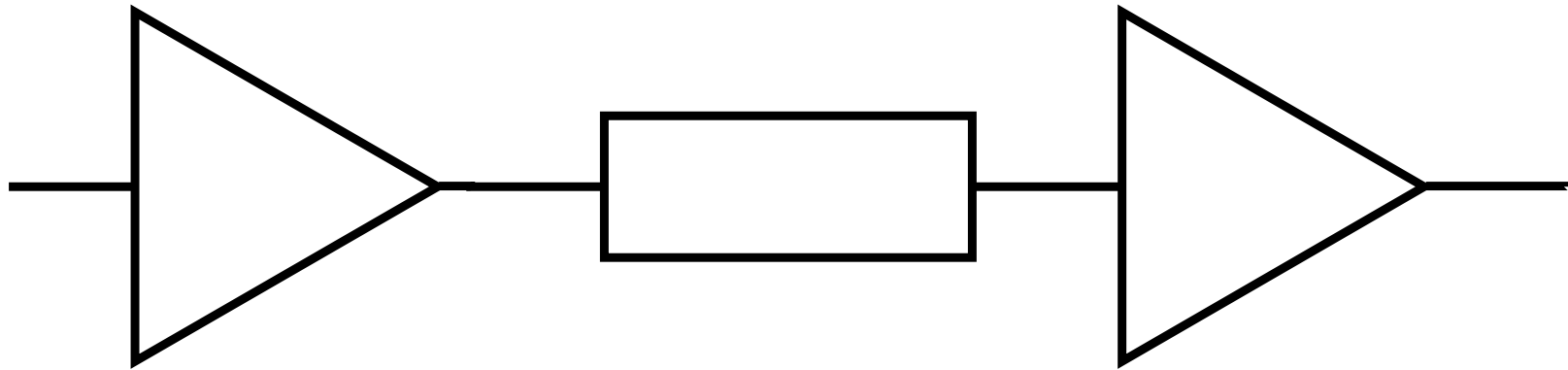
- Source and sink buffers are given



Central question:

What are the optimal buffer sizes  $b_i$  to minimize the overall delay? Is there a closed form solution?

# Equivalent circuit with buffers $i$ and $i+1$



Buffer  
size  $b_i$

$\pi$  model

Buffer  
size  $b_{i+1}$

# Calculate a recurrence relation

- Calculate Elmore delay  $ED_i$  between buffers  $i$  and  $i+1$

- Elmore delay from source to sink is  $ED = \sum_{i=0}^n ED_i$

- Find optimal buffer size  $b_i$  by setting  $\frac{\partial ED}{\partial b_i} = 0$

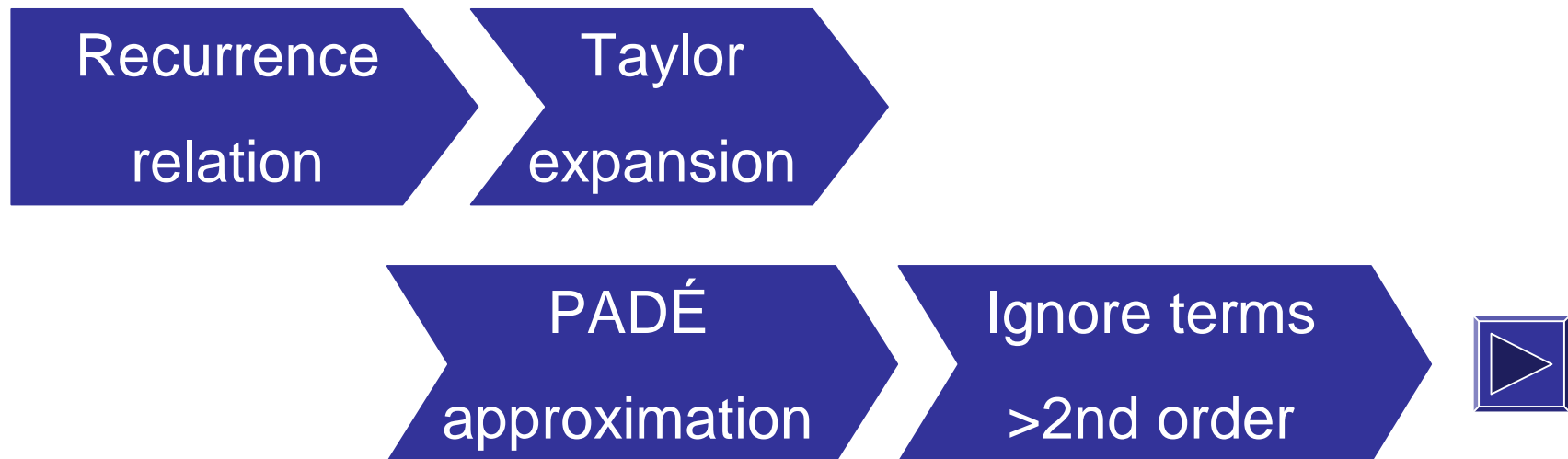
$$b_i = \sqrt{\frac{b_{i+1} + \frac{C}{C_B(1+n)}}{b_{i-1}^{-1} + \frac{R}{R_B(1+n)}}} = \sqrt{\frac{b_{i+1} + \frac{\alpha}{(1+n)}}{b_{i-1}^{-1} + \frac{\beta}{(1+n)}}}$$

$$b_0 = \lambda, b_{n+1} = \mu$$



# Closed form solution I

- Want a function  $f(x)$  that gives optimal buffer size when being evaluated at position  $x$ .
- Source is at  $x=0$  and sink is at  $x=1$
- Has to match boundary conditions  $f(0)=\bar{e}$  and  $f(1)=\mu$



# Closed form solution II

- Intermediate result

- First order ordinary differential equation with two boundary conditions
- c is integration constant, easy calculation

$$f'^2 = \left( \frac{df}{dx} \right)^2 = 2 \cdot n \cdot \beta \cdot f^3 + c \cdot f^2 + 2 \cdot n \cdot \alpha \cdot f$$

- ODE can be solved with Weierstrass elliptic function  $\wp$

- Important property:

$$\wp'(z)^2 = 4\wp^3 - g_2\wp(z) - g_3$$



## Closed form solution III

- Transformation of variables gives

$$f'(x)^2 = 4f^3 - g_2 f(x) - g_3$$



- $g_2 = g_2(\alpha, \beta, c, n)$   $g_3 = g_3(\alpha, \beta, c, n)$  are real constants
- ODE can be solved by setting

$$f(x) = \wp(x, g_2, g_3)$$

## Closed form solution IV

- According to Weierstrass, the solution  $f(x)$  can be given in terms of  $\wp$ ,  $\wp'$ .
- We denote  $h(x) = 2\beta \cdot n \cdot x^3 + c \cdot x^2 + 2\alpha \cdot n \cdot x$
- The closed form solution is given by

$$f(x) = \lambda + \frac{\sqrt{h(\lambda)}\wp'(x) + \frac{1}{2}h'(x)\left(\wp(x) - \frac{1}{24}h''(\lambda)\right) + \frac{1}{24}h(\lambda)h'''(\lambda)}{2\left(\wp(x) - \frac{1}{24}h''(\lambda)\right)^2}$$

# Closed form solution V

- Input:
  - Physical parameters of the wire (resistance, capacitance, geometrical dimensions)
  - Physical parameters of a unit size buffer (input capacitance, output resistance, intrinsic delay)
  - No. of buffers  $n$
- Evaluate  $f(x)$  at equally spaced points to get optimal buffer size:

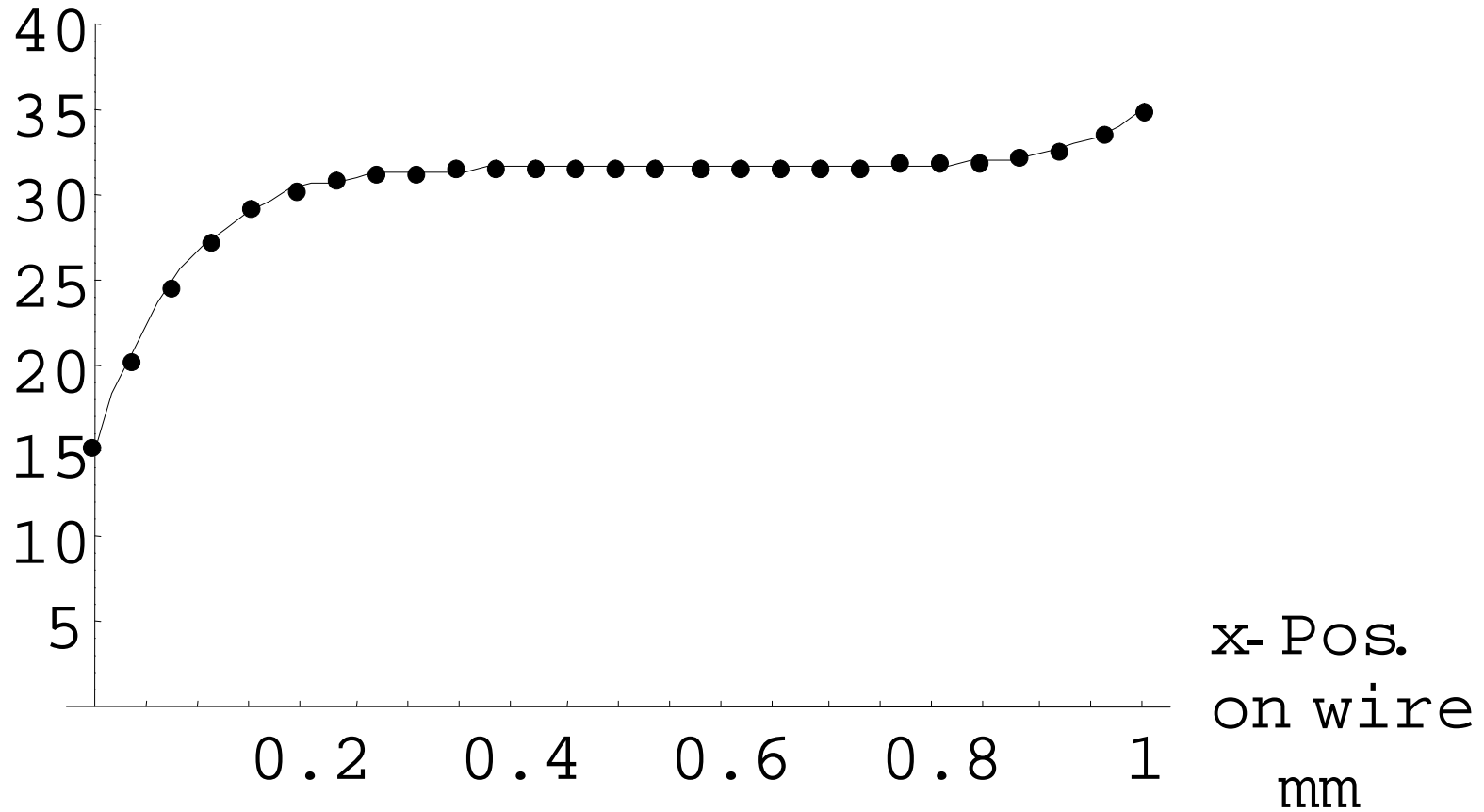
$$f(x) = f(x_i), x_i = \frac{i}{n+1}, 1 \leq i \leq n$$

# Testing procedure

- A) Iteratively
  - Fix every buffer size except buffer  $i$
  - Apply recurrence relation, giving buffer size  $b_i$  as a function of  $b_{i-1}$  and  $b_{i+1}$
  - Proceed to next buffer
  - Repeat until buffer sizes converge
- B) Use closed form solution
  - Weierstrass function available in Mathematica or Matlab
- Calculate the Elmore delay for both sets of buffer sizes
  - Compare the relative difference

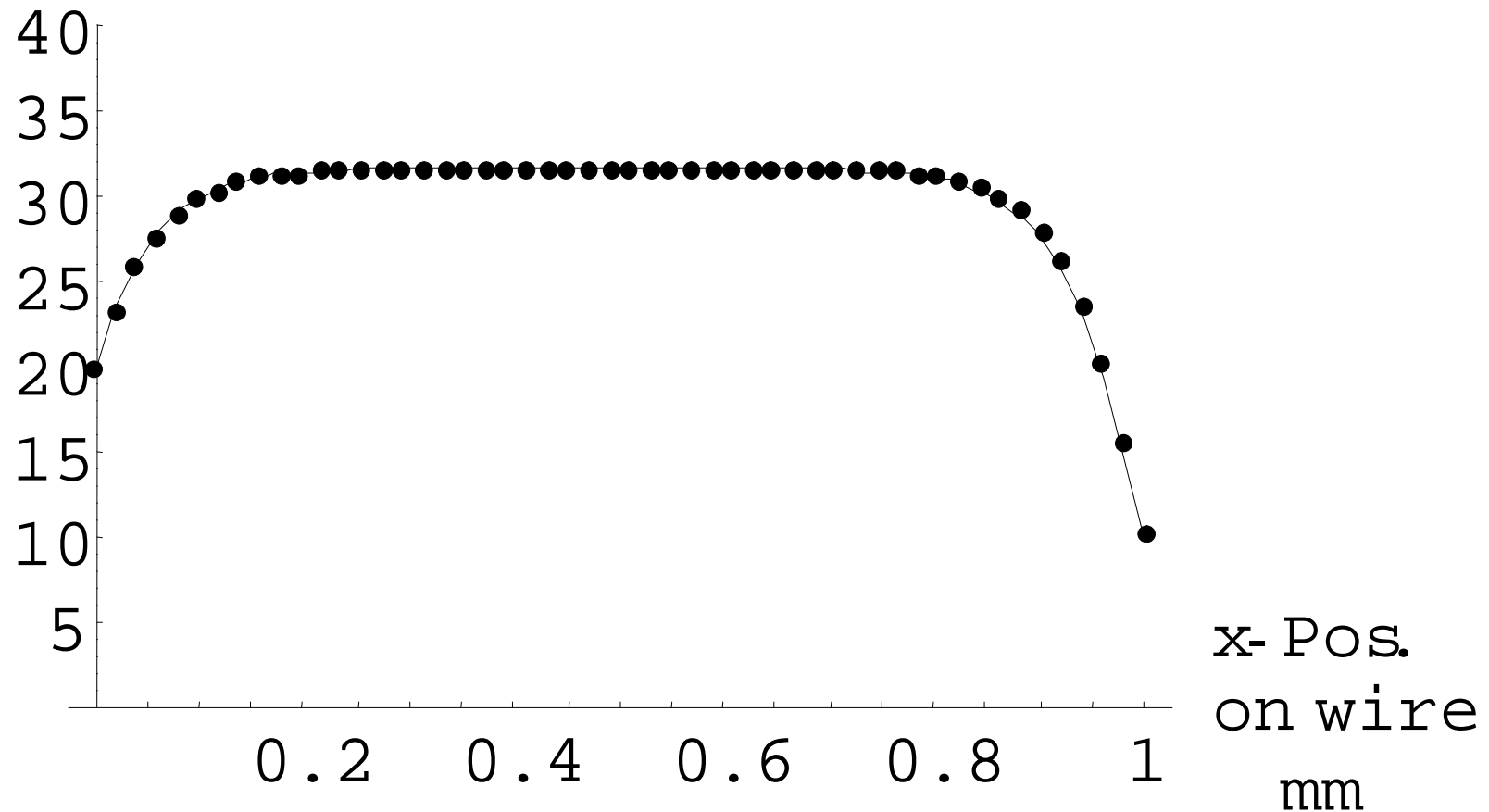
## Example 1: 25 buffers

Buffersizeb



## Example 2: 50 buffers

Buffersizeb





# Results

## ■ Achievements:

- New closed form solution
- Buffer sizes derived from closed form solution match those from the iterative method very well
- Discrepancy diminishes as  $n$  increases

## ■ Limitations:

- Only line topology allowed
- Equally spaced buffers - does not consider obstacles

# Conclusion

- Discussed the underlying model
- Derived recurrence relation
- Obtained a closed form solution in terms of Weierstrass elliptic function
- Discussed some results and consequences

Thank you for your attention!

# Substitute to apply Weierstrass function

- Start with

$$2\beta n f^3 + c f^2 + 2\alpha n t$$

- Perform substitution

$$\hat{f} = \frac{2}{n\beta} f - \frac{c}{6n\beta}$$

- Result:

$$4\hat{f}^3 - \underbrace{\left(\frac{c^2}{12} - \alpha\beta n^2\right)}_{g_2} \hat{f} - \underbrace{\left(\frac{1}{12}\alpha\beta n^2 c - \frac{c^3}{216}\right)}_{g_3}$$



# Elliptic Integrals

- Some apparently simple integrals have no simple solution.

$$I(x) = \int_c^x R(t, \sqrt{P(t)}) dt$$

- $R(.)$  is a rational function of two arguments
  - $P$  is a polynomial of third or fourth degree
  - $c$  is a constant
- Such integrals are called **elliptic integrals**, can be solved with elliptic functions
  - Jacobi's elliptic and theta functions
  - General theory: Weierstrass elliptic function

# Elliptic Functions

- $\omega_1$  and  $\omega_2$  are complex numbers (half-periods) generating a lattice  $\Gamma$  in the complex plain
- A function  $f$  is called doubly-periodic if it fulfills

$$f(z) = f(z + 2\omega_1) = f(z + 2\omega_2), \omega_1 \neq \omega_2, \forall z$$

- Elliptic functions are doubly-periodic functions with certain properties
- For each lattice  $\Gamma$  there is always a non-constant elliptic function  $\wp$ , the Weierstrass elliptic function, defined as

$$\wp(z) = \frac{1}{z^2} + \sum_{\gamma \in \Gamma \setminus \{0\}} \left( \frac{1}{(z - \gamma)^2} - \frac{1}{\gamma^2} \right)$$

# Weierstrass elliptic function I

- Consequently,  $\wp$  has three arguments:

$$\wp(z) = \wp(z; \omega_1, \omega_2)$$

- Most important property for us:

$$\wp'(z)^2 = 4\wp^3 - g_2\wp(z) - g_3$$

- $g_2$  and  $g_3$  are constants that depend on the half-periods  $\omega_1$  and  $\omega_2$ . They are called invariants.

- Can also write  $\wp(z) = \wp(z; g_2, g_3)$

- It is always possible to construct appropriate half-periods  $\omega_1$  and  $\omega_2$ , given the invariants  $g_2$  and  $g_3$ .

# Weierstrass elliptic function II

- For example, we can easily give the solution to this elliptic integral (here  $g_2$  and  $g_3$  are taken as constants)

$$u = \int_y^{\infty} \frac{ds}{\sqrt{4s^3 - g_2s - g_3}} \longrightarrow y = \wp(u; g_2, g_3)$$

- Proof: Integrate the Weierstrass differential equation
- There are (complicated) closed form solutions involving  $\wp$  and  $\wp'$  for other endpoints of the integral.





# Continuous version of the recurrence relation

- Replace  $b_i$  by  $f(x_i)$

$$f(x_i)^2 = \frac{f(x_{i+1}) + \frac{\alpha}{n+1}}{f(x_{i-1})^{-1} + \frac{\beta}{n+1}}$$

- Substitute  $x_i = x, \Delta x = \frac{1}{n+1}$   
 $x_{i+1} = x + \Delta x, x_{i-1} = x - \Delta x$   
 $\frac{1}{n} \cdot \frac{1}{n+1} \approx \Delta x^2$

- Result

$$f(x)^2 = \frac{f(x + \Delta x) + \Delta x^2 n \alpha}{f(x - \Delta x)^{-1} + \Delta x^2 n \beta}$$

# Taylor expansion + Pade approximation

- Substitute  $f(x) = e^{g(x)}$
- Perform Taylor expansion, result: rational function

$$f(x)^2 = e^{2g(x)} = \frac{P(\Delta x)}{Q(\Delta x)} = A(\Delta x)$$

- Compare coefficients to get

$$\begin{aligned} A(\Delta x) &= a_0 + a_1 \Delta x + a_2 \Delta x^2 + O(\Delta x^3) \\ &= e^{2g(x)} \underbrace{(1 + \Delta x^2 (g'' + \alpha n e^{-g} - \beta n e^g))}_{=0} + O(\Delta x^3) \end{aligned}$$

# A differential equation

- Drop terms of third order and higher
- Buffer sizing function must satisfy

- $f(x) = e^{g(x)}, x = \frac{1}{n+1} \forall i = 0, \dots, n+1$

- $g(x)$  is a solution to the second-order ODE

$$g'' = \beta n e^g - \alpha n e^{-g}$$

with the boundary conditions

$$g(0) = \ln(\lambda), g(1) = \ln(\mu)$$

# The ODE for buffer sizing function

- Apply chain rule of differentiation
- Undo previous substitution, e.g.

$$g(x) = \ln f(x)$$

- Buffer sizing function is the solution to the following first order ODE with given boundary conditions

$$f'^2 = 2n\beta f^3 + cf^2 + 2n\alpha f$$

- C can be determined numerically

