



Evaluation of Runtime Monitoring Methods for Real-Time Event Streams

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Overview

1 Motivation

2 Background

3 Comparison

- Comparison based on standard arrival curves
- Comparison based on complex arrival curves
- 4 FPGA evaluations

5 Summary





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Real-time system at runtime

Why real-time system needs to be monitored?





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What monitors do?

- System events are monitored
- Violation is reported by the monitor





State-of-art





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Dynamic counters (Kai et.al, 2011)





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- *I*-repetitive function (Neukirchner et.al, 2012)





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Contributions





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Analyzed the differences of dynamic counters and *l*-repetitive function





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- Analyzed the differences of dynamic counters and *l*-repetitive function
- Developed a new approach to apply dynamic counters to monitor periodic-burst events





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Contributions

- Analyzed the differences of dynamic counters and *l*-repetitive function
- Developed a new approach to apply dynamic counters to monitor periodic-burst events
- Prototyped hardware implementations on FPGA and presented FPGA resource usage





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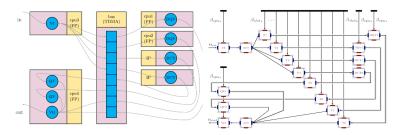
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Arrival Curve (Modular Performance Analysis)



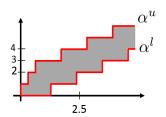
System view of the M-JPEG The corresponding RTC-MPA model

Biao Hu

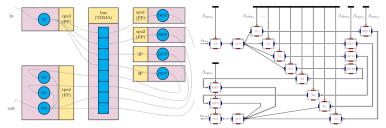




Arrival Curve (Modular Performance Analysis)



Arrival curve: Event stream model, $\alpha^{l}(t-s) \leq R[s,t) \leq$ $\alpha^{u}(t-s), \forall t \geq s \geq 0$, where R[s,t) denote the number of events in the time interval $s \leq \tau < t$.



System view of the M-JPEG The corresponding RTC-MPA model





Distance Function (Compositional Performance Analysis)

Minimum/maximum distance function

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An example: periodic events with jitter

Arrival Curve

• (a) Upper Arrival Curves :
$$\alpha^{u}(\Delta) = \lfloor \frac{\Delta+J}{\delta} \rfloor$$

(b) Lower Arrival Curves :
$$\alpha^{l}(\Delta) = max\{0, \lfloor \frac{\Delta-J}{\delta} \rfloor$$





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• (b) Lower Arrival Curves : $\alpha^{I}(\Delta) = max\{0, \lfloor \frac{\Delta-J}{\delta} \rfloor$ Distance Function

• (a) Minimum Distance Functions : $d^{min}(n) = \max\{0, n \cdot \delta - J\}$

• (b) Maximum Distance Functions : $d^{max}(n) = n \cdot \delta + J$





Assumption

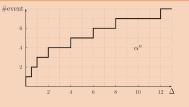
Any monotonous and time-invariant arrival curve can be conservatively approximated as the minimum on a set of staircase functions with the form $\alpha_i^u(\Delta) = N_i^u + \lfloor \frac{\Delta}{\delta_i} \rfloor$. $\forall \Delta \in \mathbb{R}_{\geq 0} : \alpha^u(\Delta) \geq \min_{i=1,n} (\alpha_i^u(\Delta))$





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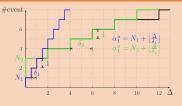




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- $\square DC_2 \leftrightarrow \alpha_2$



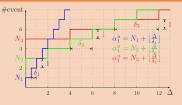




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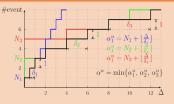




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- $\min(DC_1, DC_2, DC_3) \leftrightarrow \min(\alpha_1, \alpha_2, \alpha_3)$







Definition

A *l*-repetitive function is a special minimum distance function that satisfies: $d(n) = \begin{cases} d_n(given), & n \leq l, \\ \max_{\omega \in [1,l]} (d(\omega) + d(n-\omega)), & n > l. \end{cases}$

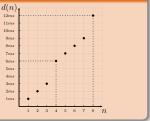




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Method: by constructing *I*-repetitive function, only the arrival time of most recent *I* events needs to be kept.





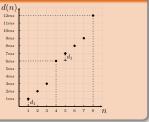


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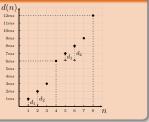


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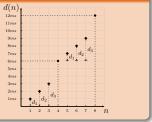


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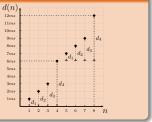


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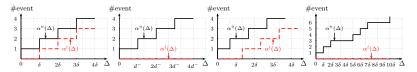
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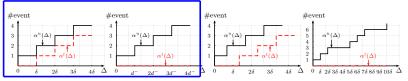
Periodic events, Sporadic events, Periodic events with jitter, Periodic burst events.







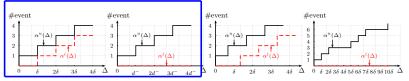
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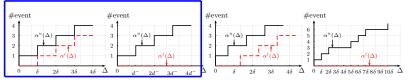
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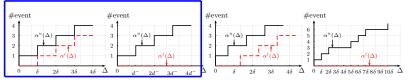
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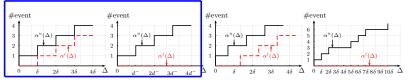
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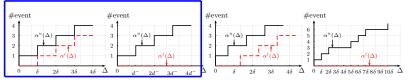
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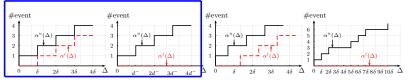
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When J is greater than δ , an initial burst N happens.

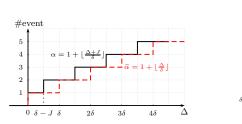


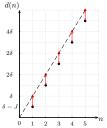


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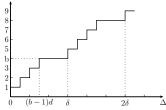
The jitter always exist when using dynamic counters or *I*-repetitive function

- Only one dynamic counter
- **2** l = 1 for *l*-repetitive function



Periodic burst events

#Events

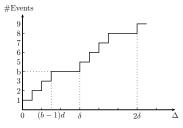


Definition

 $\begin{aligned} \alpha^{u}(\Delta) &= \lfloor \frac{\Delta}{\delta} \rfloor b + \min(\lceil \frac{\Delta - \lceil \frac{\Delta}{\delta} \rceil b}{d} \rceil, b) \\ \text{where } d \text{ is the minimum timing} \\ \text{separation, } b \text{ is the maximum events} \\ \text{within an interval } \delta. \end{aligned}$



Periodic burst events



Definition

$$\alpha^{u}(\Delta) = \lfloor \frac{\Delta}{\delta} \rfloor b + \min(\lceil \frac{\Delta - \lceil \frac{\Delta}{\delta} \rceil b}{d} \rceil, b)$$

where *d* is the minimum timing
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Monitoring method

- **1** *I*-repetitive function: I = b
- 2 dynamic counters: ?

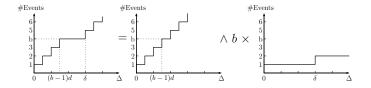




Dynamic counters



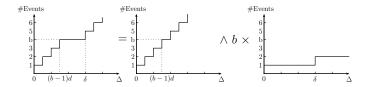
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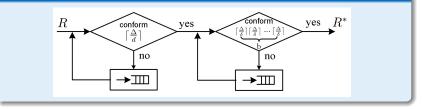
Periodic burst is equivalent to a staircase function with a period d, and the constraint of b staircase functions with a period of δ .



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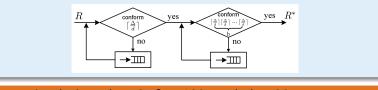






A case study of monitoring periodic burst by dynamic counters

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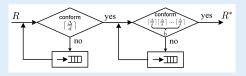
A case simulation: b = 6, $\delta = 180$, and d = 20



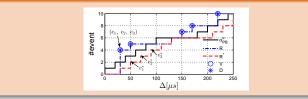


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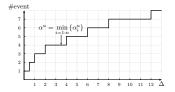






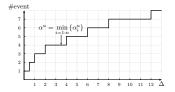






Arrival curve is the minimum of a set of staircase functions, i.e., $\alpha^{u} = \min_{i=1:n} \{\alpha^{u}_{i}\}, \text{ where } \alpha_{i} = N^{u}_{i} + \lfloor \frac{\Delta}{\delta_{i}} \rfloor (N^{u}_{i} < N^{u}_{i+1}, \delta_{i} < \delta_{i+1})$

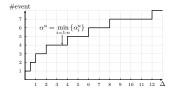




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Dynamic counters: n counters



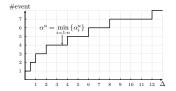


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2 /-repetitive function: ?

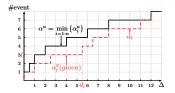






For *I*-repetitive function



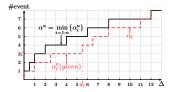


For *I*-repetitive function

1 Assume the repetitive segment is α_l^g that is given in the offline analysis. $\alpha_l(\Delta) = \begin{cases} \alpha_l^g(\Delta)(given), & \text{if } \Delta \leq d_l, \\ \alpha_l^g(\Delta - k \cdot d_l), & \text{if } \Delta > d_l. \end{cases}$



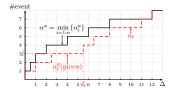




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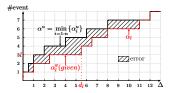




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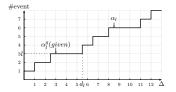
For *I*-repetitive function

1 Error percent (EP):
$$EP(\Delta) = \frac{\int_0^{\Delta} (\alpha^u(t) - \alpha^a(t)) dt}{\int_0^{\Delta} \alpha^u(t) dt}$$



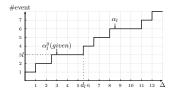






Arrival curve is segmentally repetitive, and each segment is sub-additive, i.e., $\alpha_l(\Delta) = \begin{cases} \alpha_l^g(\Delta)(given), & \text{if } \Delta \leq d_l, \\ \alpha_l^g(\Delta - k \cdot d_l), & \text{if } \Delta > d_l. \end{cases}$

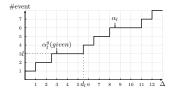




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Dynamic counters: ?



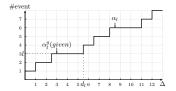


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- **1** Dynamic counters: ?
- 2 /-repetitive function: /

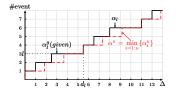






For dynamic counters

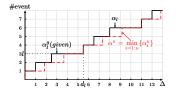




For dynamic counters 1 Assume a set of staircase functions $(\min_{i=1...n} \{\alpha_i\}, \alpha_i = N_i^u + \lfloor \frac{\Delta}{\delta_i} \rfloor)$ is used to approximate α_i , where $\delta_{max} = \max_{i=1...n} \{\delta_i\}.$



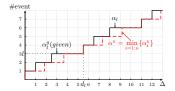
Complex arrival curve of Type 2



For dynamic counters 1 Assume a set of staircase functions (min_{i=1..n} {α_i}, α_i = N_i^u + ⌊Δ/δ_i⌋) is used to approximate α_l, where δ_{max} = max_{i=1..n} {δ_i}. 2 lim_{Δ=+∞} {min_{i=1..n} {α_l(Δ) - α_i}(Δ)} = Δ/d_l / - N_x^u - ⌊Δ/δ_{max} ⌋ ≥ 0 ⇒ d_l <= l · δ_{max}.



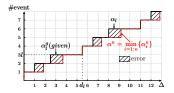
Complex arrival curve of Type 2



For dynamic counters



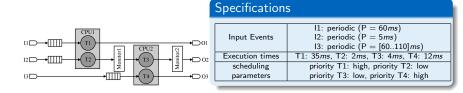
Complex arrival curve of Type 2



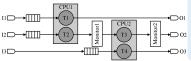
For dynamic counters

1 Error percent (EP):
$$EP(\Delta) = \frac{\int_0^{\Delta} (\alpha_l(t) - \alpha^a(t)) dt}{\int_0^{\Delta} \alpha_l(t) dt}$$

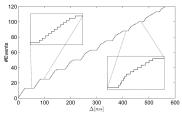




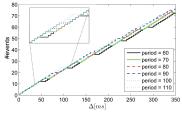




Specifications						
Input Events	11: periodic ($P = 60ms$) 12: periodic ($P = 5ms$) 13: periodic ($P = [60110]ms$)					
Execution times	T1: 35ms, T2: 2ms, T3: 4ms, T4: 12ms					
scheduling	priority T1: high, priority T2: low					
parameters	priority T3: low, priority T4: high					

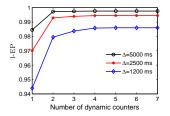


Arrival curve in Monitor 1

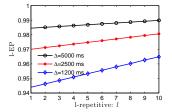


Arrival curve in Monitor 2



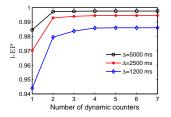


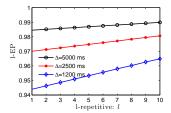
Result in monitor 1 by dynamic counters



Result in monitor 1 by *I*-repetitive function







Result in monitor 1 by dynamic counters

Result in monitor 1 by *I*-repetitive function

The needed number of dynamic counters or *I* in repetitive segment with respect to the period of input I3.

period [ms]	60	70	80	90	100	110
#DC	13	16	18	21	23	25
1	12	15	17	20	22	24





Outline

1 Motivation

2 Background

3 Comparison

- Comparison based on standard arrival curves
- Comparison based on complex arrival curves

4 FPGA evaluations



ПΠ

Setup

$Prototype \ both \ approaches \ in \ ALTERA \ Cyclone \ III \ EP3C120F780$

- Processor frequency: 50 MHz
- Timer bits: 16 bit

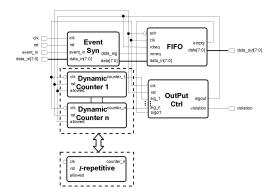


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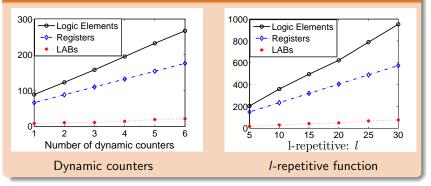






FPGA Resource

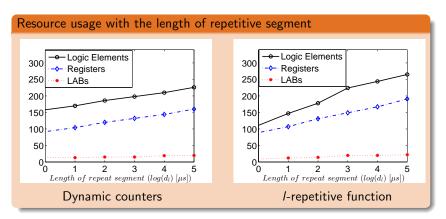
Resource usage with the number of dynamic counters or I in I-repetitive function







FPGA Resource







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Analyzed the difference of dynamic counters monitoring and *l*-repetitive function monitoring by using standard arrival curves, and two typical complex arrival curves.





- Analyzed the difference of dynamic counters monitoring and *l*-repetitive function monitoring by using standard arrival curves, and two typical complex arrival curves.
- 2 Developed a new monitoring scheme for dynamic counters to monitor periodic burst events





- Analyzed the difference of dynamic counters monitoring and *l*-repetitive function monitoring by using standard arrival curves, and two typical complex arrival curves.
- 2 Developed a new monitoring scheme for dynamic counters to monitor periodic burst events
- Prototyped the two monitoring approaches in FPGA, and found resource overhead is more sensitive to the number of dynamic counters, and not sensitive to the length of one repetitive segment.