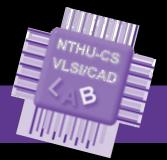
Communication Driven Remapping of Processing Element (PE) in Fault-tolerant NoC-based MPSoCs

Chia-Ling Chen, <u>Yen-Hao Chen</u> and TingTing Hwang Department of Computer Science National Tsing Hua University, Taiwan



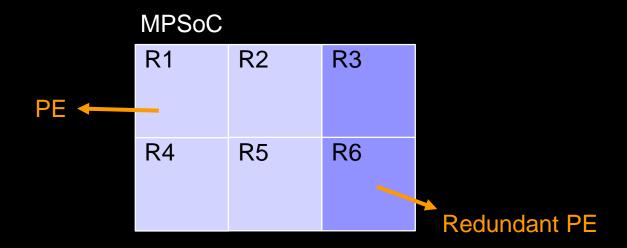


Outline

- Multi-Processor System-on-Chip (MPSoC)
 - Task remapping for fault-tolerant
- Communication driven remapping method
 - Model communication cost on edges
 - Allow tasks be moved to non-neighboring PEs
- Initial mapping improvement
- Experimental results
- Conclusions

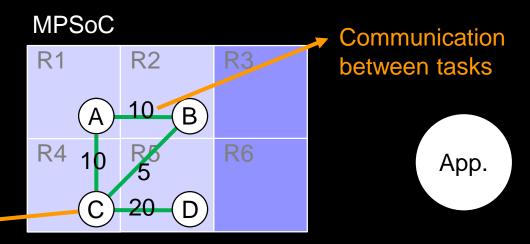
Multi-Processor System-on-Chip

- Processing elements (PEs) integrated on a chip
 - Each PE can execute a task independently
 - Redundant PEs are included for fault tolerance



Task Graph

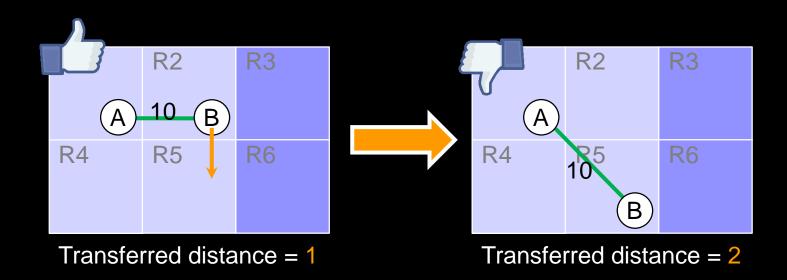
- Application is partitioned into several tasks
 - Task (node) basic executable unit
 - Must be mapped to a good PE for execution
 - Edge communication overhead
 - The amount of data (data size) transferred between tasks
 - A shorter edge (transferred distance) is preferred
 - For lower communication overhead



Task

Task Mapping

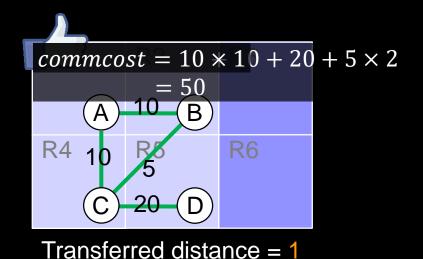
- Process of deciding PE of each task for execution
- Mapping with shorter transferred distance
 - Less communication overhead
 - Improve system performance
 - Decrease energy consumption

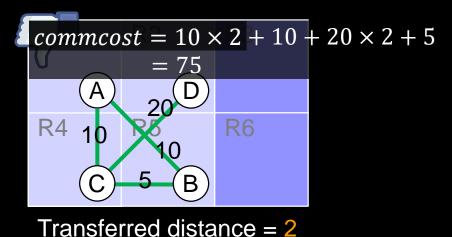


Communication Cost Metric

•
$$commcost = \sum_{i \in transferred_data} size_i \times distance_i$$

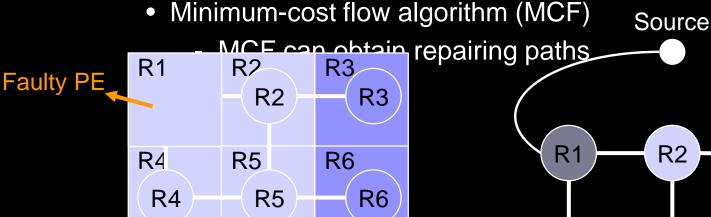
- Size transferred data size
- Distance transferred distance

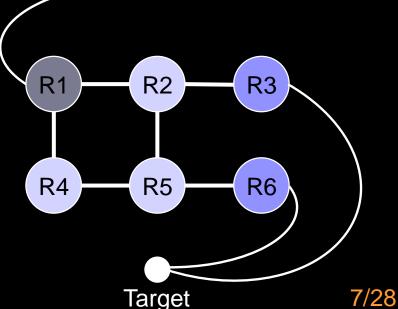




Task Remapping for Fault-tolerant

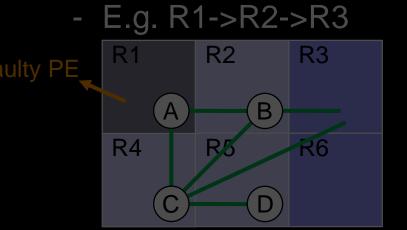
- Remap all tasks again when faulty PE detected
 - Topology graph construction
 - Repairing path
 - Each path from source to target is a remapping

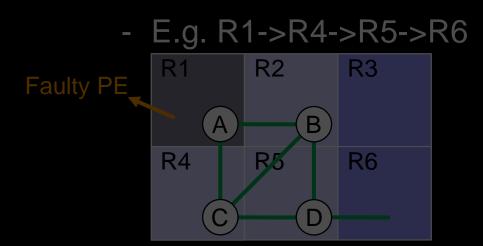


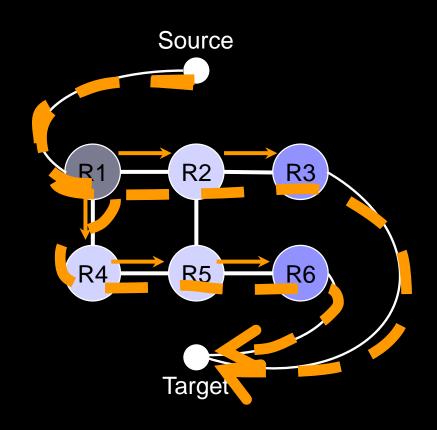


Repairing Paths

A repairing path represents a remapping solution

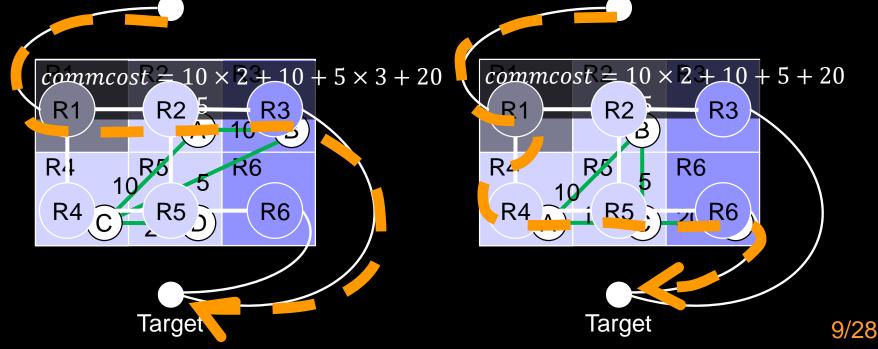






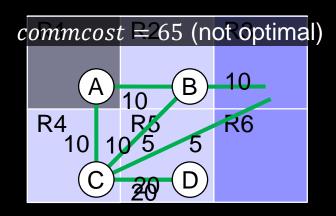
Problem Description

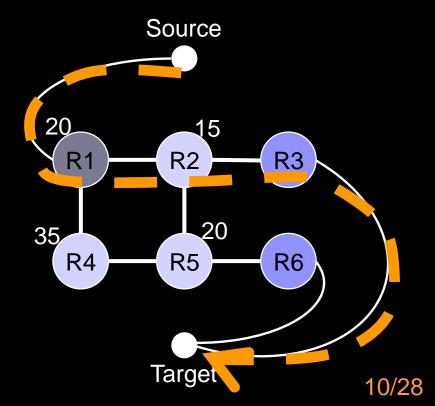
- Each repairing path indicates a valid remapping
 - All can successfully remap all tasks to good PEs
- But, the remappings have different commcost
- Thus, we would like to find the repairing path with minimized communication cost



Previous Work [1]

- The communication cost is modeled on nodes
 - Not precise enough to represent commcost
- Tasks can be only moved to neighboring PEs
 - Guarantee $\Delta distance = 1$
 - Reduce repair rate





Outline

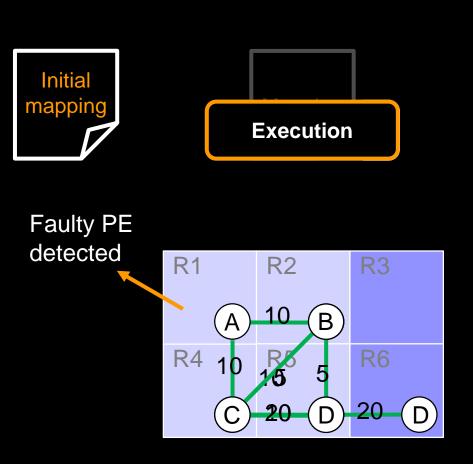
- Multi-Processor System-on-Chip (MPSoC)
 - Task remapping for fault-tolerant
- Communication driven remapping method
 - Model communication cost on edges
 - Allow tasks be moved to non-neighboring PEs
- Initial mapping improvement
- Experimental results
- Conclusions

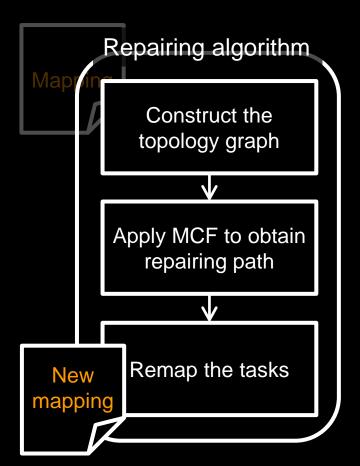
Features of Our Method

- Model communication cost on edges
 - Precisely model the communication overhead
- Allow task be moved to non-neighboring PEs
 - Greatly increase the repairing flexibility
 - Achieve 100% repair rate across all test cases

Flow Diagram of Task Remapping

Repairing is applied when detecting faulty PE





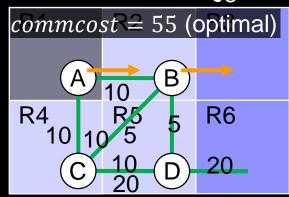
Communication Cost Model on Edges

Increment in commcost when task remapped to PE

-
$$\Delta commcost(R \rightarrow R') =$$

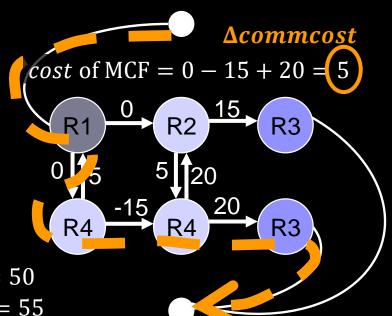
$$\left(\sum_{lenghten\ communication} size - \sum_{shorten\ communication} size \right) \times \Delta distance$$
Source

$$\Delta commcost(R2 \rightarrow R2) = (10 + 50) \times 1$$
$$= 05$$



$$commcost_{ori} = 1 * 10 + 1 * 10 + 1 * 20 + 2 * 5 = 50$$

 $commcost_{new} = 2 * 10 + 1 * 10 + 1 * 20 + 1 * 5 = 55$
Correct $\Delta commcost$ is $55 - 50 = 5$

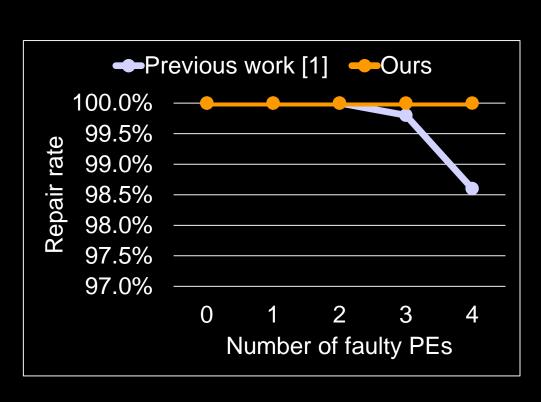


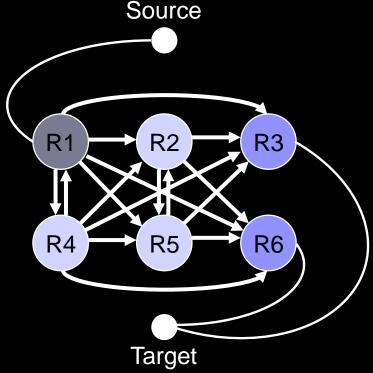
Target

14/28

Tasks moved to Non-neighboring PEs

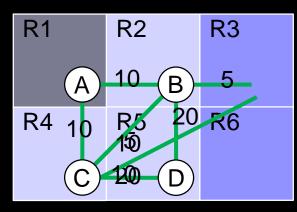
- Explore larger solution space
- Achieve higher repair rate
 - 100% repair rate for all experimental cases





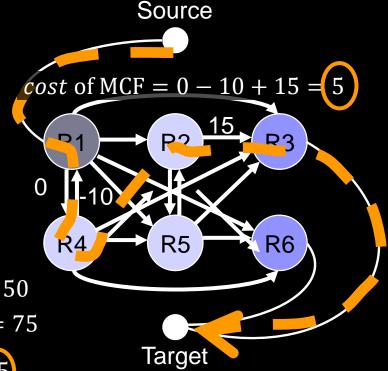
Problem on Down-and-up Paths

- The exact cost of some repairing paths can't be calculated by the MCF algorithm
 - With down-and-up (back-and-forth) movement



 $commcost_{ori} = 1 * 10 + 1 * 10 + 1 * 20 + 2 * 5 = 50$ $commcost_{new} = 2 * 10 + 3 * 10 + 1 * 20 + 1 * 5 = 75$

Correct $\triangle commcost$ should be 75 - 50 = 25

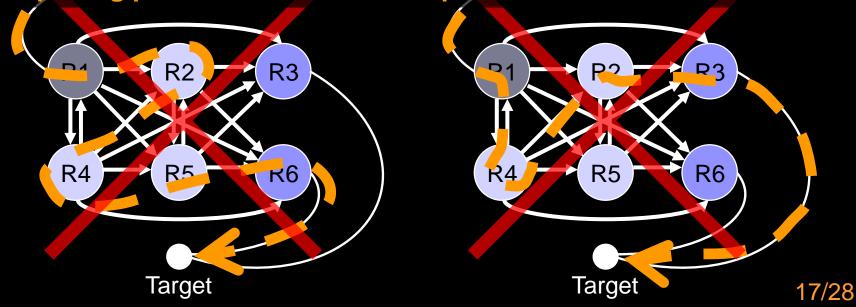


Conditions on Repairing Path

- Let $R_0 \to R_1 \to \cdots \to R_n$ be a given repairing path
- Let R_k be the k^{th} PE of the path located at (x_k, y_k)
- The repairing path must subjects to conditions:

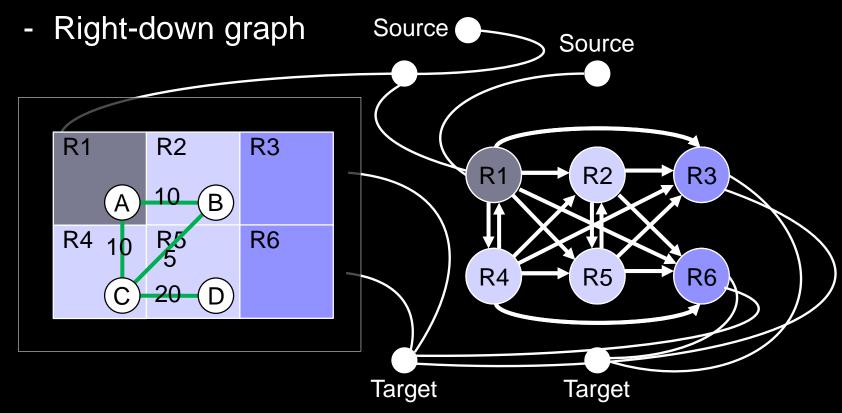
Repairing path cannot have down-and-up or back-and-forth movement

We can proof that the overall $\Delta commcost$ is exactly the cost of MCF, if the repairing path has no down-and-up or back-and-forth movement



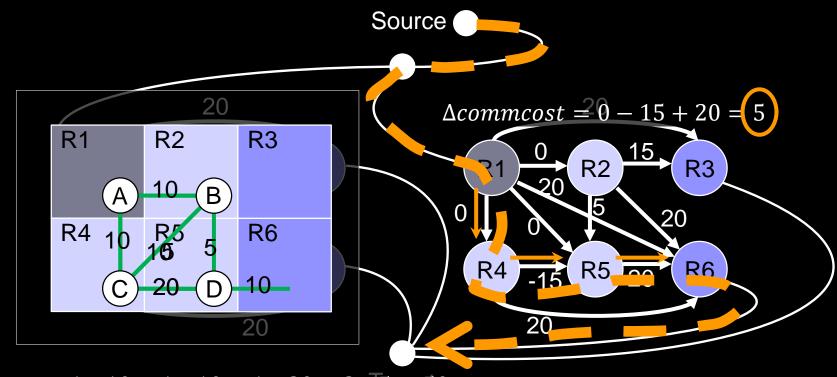
Topology Graph Construction

- Construct two topology graphs
 - Right-up graph



A Complete Example

- Obtain the repairing path with minimum overhead
 - Repairing path: R1->R4->R5->R6



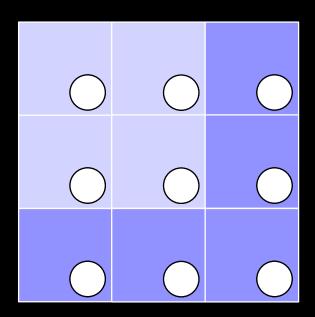
 $commcost_{ori} = 1 * 10 + 1 * 10 + 1 * 20 + 2 * 5 = 50$ $commcost_{new} = 2 * 10 + 1 * 10 + 1 * 20 + 1 * 5 = 55$ $Correct \Delta commcost is 55 - 50 = 5$

Outline

- Multi-Processor System-on-Chip (MPSoC)
 - Task remapping for fault-tolerant
- Communication driven remapping method
 - Model communication cost on edges
 - Allow tasks be moved to non-neighboring PEs
- Initial mapping improvement
- Experimental results
- Conclusions

Initial Mapping Improvement

 Our method can also be applied to improve the communication cost of a given initial mapping



Step1: Place a faulty PE

Step2: Apply our repairing method

Step3: Repeat Step1 and Step2 for iterations

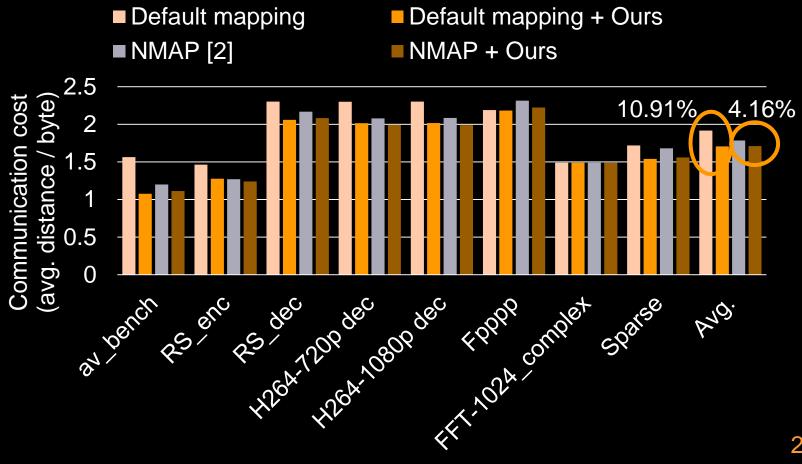
Step4: Select the mapping with lowest commcost

Outline

- Multi-Processor System-on-Chip (MPSoC)
 - Task remapping for fault-tolerant
- Communication driven remapping method
 - Model communication cost on edges
 - Allow tasks be moved to non-neighboring PEs
- Initial mapping improvement
- Experimental results
- Conclusions

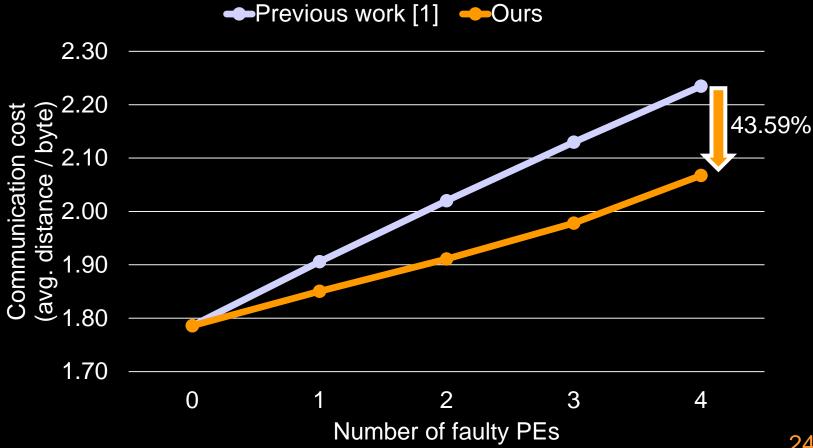
Initial Mapping Improvement

- 10.91% improvement on default mapping
- 4.16% improvement on NMAP [2]



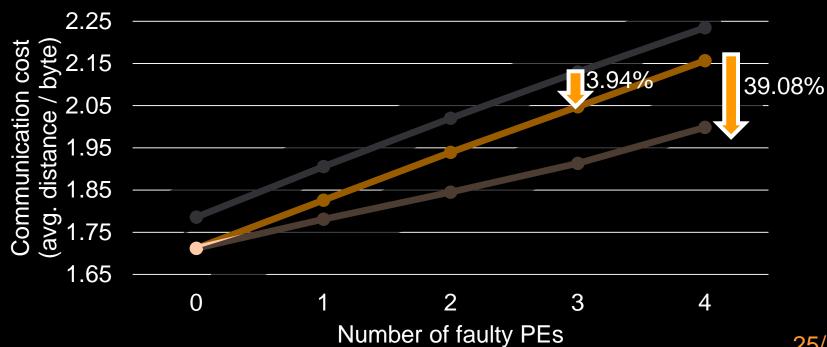
Repairing Communication Overhead

- 43.59% less communication overhead on average
 - Use a more precise communication cost model



Overall Effect

- 3.94% fewer overhead on initial mapping
- 39.08% less commcost for remapping
 - Initial mapping: NMAP [2]; Repairing: previous work [1]
 - Initial mapping: NMAP+Ours; Repairing: previous work [1]
 - Initial mapping: NMAP+Ours; Repairing: Ours



Outline

- Multi-Processor System-on-Chip (MPSoC)
 - Task remapping for fault-tolerant
- Communication driven remapping method
 - Model communication cost on edges
 - Allow tasks be moved to non-neighboring PEs
- Initial mapping improvement
- Experimental results
- Conclusions

Conclusions

- A communication driven remapping algorithm
 - Tolerate processing element failures
 - Use precise communication cost model
 - Allow flexible non-neighboring task movement
 - Improve the initial mapping
 - Be compared with previous work
 - Higher repair rate
 - Less communication overhead
 - Initial mapping (4.16%)
 - Remapping (43.59%)

Thank you for listening

Backup Slides

Environment & Benchmarks

- Transaction Generator (TG) [10]
- NoC traffic benchmark suite MCSL [12]
- Mesh 4 × 5 architecture

Application	# Tasks	# Communication Links	
		Among tasks	Among PEs
av_bench [11]	40	57	25
RS_enc	262	348	18
RS_dec	182	392	71
H264-720p_dec	2311	3461	65
H264-1080p_dec	5191	7781	65
Fpppp	334	1145	120
FFT-1024_complex	16384	25600	116
Sparse	96	67	34

Lemma of Topology Graph

Let
$$\begin{cases} \operatorname{task} A \text{ is moved from } R_i = (x_i, y_i) \text{ to } R_{i+1} = (x_{i+1}, y_{i+1}) \\ \operatorname{task} B \text{ is moved from } R_j = (x_j, y_j) \text{ to } R_{j+1} = (x_{j+1}, y_{j+1})' \end{cases}$$
 where $i < j$ and $i, j \in \{0, 1, ..., n\}$

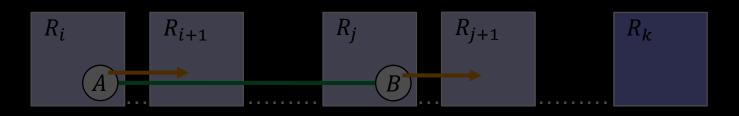
$$Suppose conditions \begin{cases} x_0 \le x_1 \le \cdots \le x_n \text{ or } x_0 \ge x_1 \ge \cdots \ge x_n \\ y_0 \le y_1 \le \cdots \le y_n \text{ or } y_0 \ge y_1 \ge \cdots \ge y_n \end{cases}$$
 are satisfied
$$\Delta commcost_{A,B} = size_{A,B} \times \left(distance(R_{i+1}, R_{j+1}) - distance(R_i, R_j) \right)$$

$$= size_{A,B} \times \left(\left(\left| x_{j+1} - x_{i+1} \right| + \left| y_{j+1} - y_{i+1} \right| \right) - \left(\left| x_j - x_i \right| + \left| y_j - y_i \right| \right) \right)$$

$$= size_{A,B} \times \left(distance(R_j, R_{j+1}) - distance(R_i, R_{i+1}) \right)$$

$$= size_{A,B} \times \left(distance(R_j, R_{j+1}) - distance(R_i, R_{i+1}) \right)$$

$$= size_{A,B} \times \Delta distance(B) - size_{A,B} \times \Delta distance(A)$$



Cost of MCF Algorithm ($\Delta commcost$)

$$\Delta commcost = \sum_{\text{Task graph edge } A,B} \Delta commcost_{A,B}$$

$$= \sum_{\text{Task graph edge } A,B} size_{A,B} \times \Delta distance(\mathbf{A}) - size_{A,B} \times \Delta distance(\mathbf{B})$$

$$= \sum_{\text{Moved task } A} \left(\sum_{\text{shorten edge } A,B} size_{A,B} - \sum_{\text{lenthen edge } A,B} size_{A,B} \right) \times \Delta distance(\mathbf{A})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_{i+1} \in \text{repairing path}} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_i \to R_i} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_i \to R_i} \Delta commcost(R_i \to R_{i+1})$$

$$= \sum_{R_i \to R_i} \Delta commcost(R_i \to R_i)$$

$$= \sum_{R_i \to R_i} \Delta commcost(R_i \to R_i)$$