# An Exact Schedulability Analysis for Global Fixed-Priority Scheduling of the AER Task Model

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#### Introduction

Safety-critical systems needs to be carefully designed to meet the specific requirements they have as catastrophic consequences can be resulted otherwise.

Examples for safety-critical systems,

- Avionics applications
- Automotive applications

Timing verification becomes crucial in these systems.





#### Introduction

Shifting to multi-core processors is considered as the only solution to the performance limitation challenge faced with single-core systems.



#### Phased-Execution Task Models

- Phased-execution models are a good solution when addressing these challenges.
- The PRedictable Execution Model (PREM) [5] and the AER Model [2] falls under this category.
- In the AER Model, a task is divided into 3 phases.
  - 1. Acquisition Read and copy all necessary data and instructions for the task from the main memory to the local memory
  - 2. Execution Execution of the task without having to access the main memory
  - 3. Restitution Write-back the results to the main memory after the execution



### Schedulability Analysis for the AER Model

- Schedulability analysis for the AER model under global fixed-priority scheduling has been proposed in [1] and [4].
- However, these schedulability tests are non-exact and only sufficient.
- Due to their pessimism they may classify a schedulable task set as unschedulable.
- ★ An *exact* schedulability test provides the result as unschedulable *only* when it is not actually schedulable.
- ★ Therefore, having an exact schedulability test is needed to correctly evaluate the schedulability of a task set.
- ★ Also, an exact test provides a reference for the comparison of the accuracy of the existing schedulability tests.

- Our work introduces the first exact schedulability test for the AER model under global fixed-priority scheduling.
- Our schedulability test utilizes timed automata (TA) where the schedulability problem is described as a reachability problem.
- Evaluation of the proposed analysis using synthetic task sets against the state-of-the-art, which shows that the proposed analysis provides up to 65% more schedulable task sets than the state-of-the-art, while providing acceptable solving times.

## System Model

- The multiprocessor system we consider has *m* identical cores.
- The task set is comprised of *n* independent tasks  $\tau = {\tau_1, \tau_2, ..., \tau_n}$ .
- Each task  $\tau_i$  can be represented by the tuple  $(T_i, D_i, O_i, P_i, C_i)$ .



A task set is considered schedulable when all of its tasks finish before their deadlines.

- The total execution time,  $C_i = C_i^r + C_i^e + C_i^w$
- The AER model executes each phase of the task without any interference from other tasks executing on other cores.
- However,  $C_i^{x} \in [C_i^{x_{min}}, C_i^{x_{max}}]$ , where  $x = \{r, e, w\}$
- Therefore, the total execution time,  $C_i \in [C_i^{min}, C_i^{max}]$

- We assume global non-preemptive fixed-priority scheduling.
- The tasks are scheduled from a memory perspective.
- The memory phases are added to a common global priority queue until they are scheduled.
- When the ready queue is not empty and the bus is available, 3 conditions are checked.
  - **C1**) If the task at the front wants its write phase to be scheduled.
  - $\mathbb{C}^2$  If the task at the front wants its read phase to be scheduled and if a core is available.
  - C3 If the task at the front wants its read phase to be scheduled and if a core is not available.

## Network of Timed-Automata in UPPAAL



- 1. The common period values reported by Kramer et al. [3] for automotive applications were used.
  - Period values for tasks  $(T_i)$ :  $\{1, 2, 5, 10, 20, 50, 100, 200, 1000\}$  ms
  - Respective probability percentages: {3, 2, 2, 25, 25, 3, 20, 1, 4}
- 2. For a given total utilization U, the Dirichlet-Rescale (DRS) algorithm is used to randomly generate  $u_i$  utilization values for each task in the task set.

## Schedulability Evaluation



- Basic configuration :  $U = 1, m = 4, \gamma = 0.1$  and n = 10
- Each plotted point represents 200 task sets.

#### **Runtime Evaluation**



- Basic configuration :  $U = 1, m = 4, \gamma = 0.1$
- Each plotted point represents 200 task sets.

## Log-uniform Period Distribution



- Basic configuration :  $m = 4, \gamma = 0.1, n = 10$
- Log-uniform period distribution between 10ms and 100ms
- Each plotted point represents 100 task sets.

- 1. We present an exact schedulability test using timed automata for the globally scheduled AER model.
- 2. The presented schedulability test reduces pessimism of the existing tests by a large margin.
- 3. It provides a baseline against which other tests can be compared.

### Reference

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