Optimizing Decision Diagrams for Measurements of Quantum Circuits

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Background

Variational quantum algorithm (VQA):

•Promising application of near-term quantum computers

•Require measurement methods that efficiently extract information from quantum state

Shallow-circuit quantum measurement:

•Estimate quantum state by measuring each qubit

•Measurement based on single-qubit gate such as Pauli operations *X*, *Y*, *Z*

•Robust to noise of quantum computers

Classical Shadow:

•Obtain the estimator by polynomial number of measurements

- •There are several methods for increasing accuracy of estimator
- •Decision Diagram (DD)-based method [Hillmich et al. 2021]: High accuracy

Shallow-circuit quantum measurements

Consider *n*-qubit Hamiltonian *H* represented by following equation

$$H = \sum_{P \in \{I, X, Y, Z\}^n} \alpha_P P$$

Shallow-circuit quantum measurement problem: Given quantum state $\rho \in \mathbb{C}^{2^n}$, we estimate $\text{Tr}(H\rho)$ with small variance by single qubit pauli-operation and measurement on computational basis ($|0\rangle$, $|1\rangle$)



Optimization of probability distribution β over measurement basis $\mathcal{P}^n(\mathcal{P} = \{X, Y, Z\})$

Variance of estimator

$$\operatorname{Var}(\nu) = \frac{1}{S} \left(\left(\sum_{P,Q} \alpha_P \alpha_Q g(P, Q, \beta) \operatorname{Tr}(PQ\rho) \right) - \left(\operatorname{Tr}(H\rho) \right)^2 \right)$$

This is small when # of measurement basis covering P, Q is small Reduces the variance

Explanation of terminology "cover": $Cover(P) \coloneqq \{B \in \mathcal{P}^n | B_i = P_i \text{ whenever } P \neq I\}$ *B* covers $P \Longrightarrow Tr(P\rho)$ can be estimated by measurement on *B*. e.g.) *XXXX* covers *XIXI*

Existing decision diagram (DD)-based method [Hillmich, et al. 2021]:
Extract measurement basis from DD representing Hamiltonian
Demonstrate high accuracy compared with other classical shadow methods

Extract measurement basis from DD

H = 0.25XXX + 0.25YYX + 0.5ZZZDD for quantum measurements: Convert Rooted directed acyclic multi-graph G = (V, E)Date given to edges: Root Single-qubit Pauli-operator •Weight $w(e) \in (0,1]$ Probability of selecting edge Requirements: Y(0.25)Z(0.5)• $v \in V$ has at least one out-going edge • For out-going edge of v ($e \in out(v)$), $\sum_{e \in out(v)} w(e) = 1$ Path: Corresponds measurement basis Terminal Estimate $Tr(H\rho)$ using measurement basis obtain by random walk

Factor of variance: Edge weight, Shape of graph

Building method for DD

Flow of existing building method for DD

- Reduce Hamiltonian
 Reduce # of Pauli-operators
 Tuning probability distribution
- Initialize DD
 Build DD from reduced Hamiltonian
- •Normalize DD Normalize edge weights Eliminate redundant nodes



Reduce Hamiltonian

H = -0.810IIII + 0.045XXXX + 0.045XXYY + 0.045YYXX + 0.045YYYY + 0.172ZIIII - 0.225IZII + 0.172IIZI - 0.225IIIZ + 0.120ZZII + 0.168ZIZI + 0.166ZIIZ + 0.166IZZI + 0.174IZIZ + 0.120IIZZ

Reduce (Pauli grouping)

Merge Pauli-operatorsTuning probability distribution

Reduced Hamiltonian $\mathcal{R}(H)$

= 0.045XXXX + 0.045XXYY + 0.045YYXX + 0.045YYYY + 1.714ZZZZ

Coefficients correspond to probability distribution over measurement basis

Algorithm: Check whether merging is possible for all combinations of terms

Initialize DD

 $\mathcal{R}(H) = 0.045X_1X_2X_3X_4 + 0.045X_1X_2Y_3Y_4 + 0.045Y_1Y_2X_3X_4 + 0.045Y_1Y_2Y_3Y_4 + 1.714Z_1Z_2Z_3Z_4$



Qubit order: Order in which bits are referenced

Create path for each pauli-operator
Give weight to bottom edge (Weight corresponds to coefficient)

Next, normalize this DD •Normalize edge weights •Merge some nodes

Normalize edge weight

Requirement for edge weight:

For out-going edge of v ($e \in out(v)$), $\sum_{e \in out(v)} w(e) = 1$

While propagating weight to root node,

normalize edge weight so that above requirement is satisfied.



Initial DD



Merge equivalent node

Merge node v, u, when the two nodes satisfy following conditions. (1) V_c (child-nodes of node v) = V_c' (child-nodes of node u) (2) For any node $a \in V_c$, Pauli-operation and weight of edge (v, a) and edge (u, a) are equivalent.



Node reduction by merging I-edges

When graph satisfies some conditions, following operation can be applied



Our contribution

Issue of DD: Size of Hamiltonian: Large \Rightarrow Size of DD: Large Can not be applied to large-size Hamitonian

Our contribution: •Propose a method for reducing size of DD •Our method does not compromise accuracy of estimator

Relaxing condition of equivalent nodes

Merge nodes v, u, when the two nodes satisfy the following conditions (1) V_c (child-nodes of node v) = V_c' (child-nodes of node u) (2) For any node $a \in V_c$, Pauli operations of edge (v, a) and edge (u, a) are equivalent

Difference from existing method: Ignore edge weight



Existing method: v, u are not equivalent

Proposed method: *v*, *u* are equivalent

of nodes merged is increased Reduce # of nodes in DD

Edge weighting scheme



Qubit ordering

Reduce # of nodes by qubit ordering

DD for quantum measurements# of nodes: Depends on qubit orderBDD (representative DD)# of nodes:

Depends on variable order (corresponds to qubit order)

Naïve algorithm for finding qubit order that minimizes # of nodes: •Examine all qubit orders

•Time complexity: O(n!)

Proposed method:

•Fast algorithm based on dynamic programming (technique utilized in BDD variable ordering)

•Obtain the optimal ordering when ignoring the reduction by merging I-edges 15



Divide DD

 $Q = \{q_1, q_2, ..., q_n\}$ (*q_i*: Qubits in a DD)



Problem of finding optimal qubit order: We can separate upper part and lower part

Optimal qubit order: Examine all combinations of *S*

of boundary nodes

Information for finding optimal qubit order:Set of qubits *S*# of boundary nodes



Master Profile Chart(MPC)



Finding qubit order that minimize # nodes = Finding minimum weighted path

Breadth fast search can solve this problem Time complexity: Exponential to # of qubits

Experimental results

Comparison of existing method with proposed method

			CS [4]	Existing DD-based [21]			Proposed		
Molecule	Encode	Terms	Vari.	Node	Edge	Vari.	Node	Edge	Vari.
LiH (12-qubit)	JW	751	266	166	271	11.3	74	137	8.28
	Parity	778	760	215	429	31.2	172	316	36.7
	BK	765	163	290	521	29.6	167	300	44.1
BeH ₂ (14-qubit)	JW	796	1792	319	484	123	177	297	314
	Parity	831	3524	375	702	635	266	468	697
	BK	833	2273	336	545	561	226	390	3597
H ₂ O (14-qubit)	JW	1302	4789	389	682	1444	220	397	1474
	Parity	1332	11209	540	1065	3552	229	492	5660
	BK	1333	26004	476	864	5325	265	487	9481

Proposed methods: Reduce DD size without compromising accuracy

[4] S. Aaronson, "Shadow tomography of quantum states," 2020.

[21] S. Hillmich, et. al, "Decision diagrams for quantum measurements with shallow circuits," 2021.

Conclusion

DD based quantum measurement:High accuracyScalability issue regarding DD size

Proposed methods:Relax condition of equivalent nodesEdge weighting schemeQubit ordering optimization

Experimental results:

•Demonstrate proposed methods can reduce DD size without compromising accuracy