

## SISCO: <u>Selective Invariant Sharing,</u> <u>Clustering and Ordering for Effective</u> <u>Multi-Property Formal Verification</u>

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#### Introduction

- Model Checking tools are not designed for multi-property verification
- Static Cone-of-Influence (COI) methods exists but they do not scale for practical or industrial designs
- Key observation from recent experience:
  - 1. Inductive Invariants discovered while proving one property can help another property
  - 2. This motivates ordering properties in a way that earlier ones helps the later ones
  - 3. Unfortunately there is no way to know a priori which one will help which
- Question that lead to this work:
  - 1. Are all invariants required for solving undecided goals?
  - 2. Can we improve effectiveness of falsified goals?

# Current State of the Art $Properties P \xrightarrow{P_1}{P_2} Properties P \xrightarrow{P_1}{P_2} P_2$



#### **Property Directed Reachability (PDR)** Frame = 0cube= Generalization **Create Solver** GetBadState(...) Procedure γ IF(cube IF(BlockState(...)) != NULL Ν Ν Push Clauses Increment Frame Up and Down SAT and Create Solver Υ IF(PushClause(...)) **UNSAT STOP**

## **Motivating Examples**



### Motivating Example (Falsified)



	races		With CEX Traces						
#Block	#PO	#Calls	#Clause	#Time	#Block	#Block #PO #Calls #C		#Clause	#Time
8	42	76	12	<1	8	29	63	9	<1



Witho	out CEX	traces a	ces and invariants With CE				CEX traces and all invariants				With CEX traces and selective invariants					
#Block	#PO	#Calls	#Clause	#Time	#Block	#PO	#Calls	#Clause	#Time	#Block	#PO	#Calls	#Clause	#Time		
7	31	135	23	<1	3	13	60	14	<1	4	14	65	15	<1		

## **Motivating Example**

Design Stats of H	WMCC Benchmark 6s421
Flops	951
Inputs	153
Properties	150
Gates	6294

Without CEX traces and invariants					With CEX traces and all invariants					With CEX traces and selective invariants					
#Block	#PO	#Calls	#Clause	#Time	#Block	#PO	#Calls	#Clause	#Time	#Block	#PO	#Calls	#Clause	#Time	
882	44681	272187	23399	87.91	884	32053	4336287	1066727	816.16	883	76745	286433	33625	82.92	

# SISCO

<u>Selective Invariant Sharing, Clustering and Ordering for</u> Effective Multi-Property Formal Verification



## **SISCO: Contributions**

#### **Clustering and Ordering**

- Clustering of goals based on clause and variable overlap
- Order goals within the cluster based on statistics dumped

#### **Selective Invariant Sharing**

- Selectively share invariants of already proven goals within a cluster
- Avoid sharing of invariants across different clusters

#### 03

01

02

#### **Storing CEX traces for falsified goals**

• CEX traces are stored within a cluster to avoid repetitive task of the PDR runtime engine

## **Clustering and Ordering**

Clustering of goals based on clause and variable overlap

- $\circ$  Jaccard Index (A,B) =  $|A \cap B| / |A \cup B|$
- JI (clauses(P<sub>1</sub>), clauses (P<sub>2</sub>)) > CL\_Upper\_Threshold
- $\circ~$  JI (clauses(P1), clauses (P2)) > CL\_Lower\_Threshold  $$\Lambda$$

JI (variables(P<sub>1</sub>), variables(P<sub>2</sub>)) > Literals\_Threshold

Order properties within the cluster based on statistics

 $\circ$  (Total calls to SAT solver / Frames Explored)\*

P<sub>2</sub> P₄

\* PURSE: Property Ordering Using Runtime Statistics for Efficient Multi-Property Verification

## **Selective Invariant Sharing**

When a property gets proven store invariants along with flop count.

- Property  $P_0$ : 100 and Property  $P_1$ : 111
- Invariant of Property P<sub>0</sub>: {001, 010, 100, 111}
- Flop count:  $\{x_0: 4, x_1: 4, x_2: 4\}$
- Maintain a map that has the count of each flops and sort it in descending order of flops count.

When the next property starts

- $\circ$  ~ Compute the flop count of the frame clauses and sort them similarly
- Pick the first 'k' variables from the invariant map and the current property frame clauses and check the overlap.
- If it crosses a certain threshold then add that invariant
- Check for other stored invariants as well
- Continue the process for each new frames added

![](_page_14_Picture_12.jpeg)

![](_page_14_Figure_13.jpeg)

## **Storing CEX Traces**

When a property gets falsified store CEX trace

- $\circ$  ~ The hash of each state in the CEX trace is stored also
- $\circ$   $\;$  Initially when a bad state is encountered its hash is matched
- For example the CEX trace for property 13 is:

 $\bullet \quad \{13 \rightarrow 11 \rightarrow 9 \rightarrow 6 \rightarrow 2 \rightarrow 1 \rightarrow 0\}$ 

- When property 15 starts it needs to block the following:
  - { (15,1) (15,2) (15,3) (15,4) (15,5) (14,4) (12,3) (11,3) }
- Once a state that needs to be blocked exists in the CEX trace then stop and report SAT
- CEX trace is recomputed by extracting the states from the current queue followed by the states from the CEX trace of the already falsified property
- $\circ$   $\;$  Here CEX trace for property 15 is:

$$\{15 \rightarrow 14 \rightarrow 12 \rightarrow 11 \rightarrow 9 \rightarrow 6 \rightarrow 2 \rightarrow 1 \rightarrow 0\}$$

![](_page_15_Figure_12.jpeg)

![](_page_16_Figure_0.jpeg)

![](_page_17_Figure_0.jpeg)

![](_page_18_Figure_1.jpeg)

![](_page_19_Figure_1.jpeg)

Designs		Des	ign Stats		Cluster In	formation	Undecided Goals			
	#Inputs	#Flops	#Gates	#Prop.	#Clusters	#Time	Baseline	State-of-Art	SISCO	
6s168	34	98	2722	16	1	0.01	2	0 (7200)	3	
pdtvsarmultip	17	130	2743	33	1	0.02	3	0 (772)	0 (822.10)	
pdtvsar8multip	23	195	6956	33	2	0.03	4	2	2	
bob12m18m	57	261	2140	163	1	0.14	158	6	6	
bob12m09	44	285	30250	85	4	0.15	0 (4037.13)	0 (4127)	0 (4631.15)	
sm98tcas16multi	279	310	5068	6	1	0.01	4	1	1	
sm98tcas16tmulti	279	310	5168	6	1	0.01	4	5	5	
6s395	21	463	3703	129	13	1.73	128	1	1	
s13207.1	62	638	2721	152	1	0.56	59	6	6	
s13207	31	669	2721	121	1	0.2	34	7	7	

Designs		Desigr	n Stats		Cluster In	formation	Undecided Goals			
Designs	#Inputs	#Flops	#Gates	#Prop.	#Clusters	#Time	Baseline	State-of-Art	SISCO	
6s421	153	951	6294	150	1	1.5	0 (87.91)	0 (3004.92)	0 (116.95)	
b17	37	1415	27549	97	1	3.63	78	0 (2935.16)	0 (3838.53)	
6s419	279	1653	26811	9	2	0.13	9	3	3	
6s258	286	1790	27501	80	9	9.6	80	77	77	
6s391	433	2686	13716	387	3	3.21	375	56	56	
mentorbm1	224	4376	31613	13	1	0.11	0 (553.90)	1	0 (1821.35)	
6s116	4912	4922	25104	17	1	0.18	1	3	3	
6s321	22	13126	66695	6	1	0.07	6	2	2	
6s118	452	13706	410804	100	2	2.06	100	83	83	
6s281	91099	177235	2179584	90	1	1.86	76	71	72	

## **Experimental Results (Observations)**

Selective Invariant Shared All Invariant Shared

![](_page_22_Figure_2.jpeg)

- SISCO outperforms all invariant sharing in 75.13 % of the cases and provides 2x improvements in 37.56 % cases
- SISCO provides an average improvement of 4.73x in the runtime of individual goals across all designs
- Frames explored by the last unsolved goal is more in SISCO than all invariant sharing

## Conclusion

#### 01

#### **Clustering and Ordering**

SISCO clusters and orders properties based on data dumps after the initial unrolling phase

#### 02

03

#### **Storing Invariants and CEX Traces**

Invariants are shared selectively and CEX traces are stored along with their hashes for faster comparison

#### **Flexibility in Framework**

One can use a completely different ordering, clustering or modified PDR and tests its effectiveness

# Thanks!

Do you have any questions?

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